THE DESIGN AND APPLICATION OF A MICROPROCESSOR DEVELOPMENT SYSTEM

A Thesis

Presented to

the Faculty of the Department of Computer Science University of Houston

In Partial Fulfillment

of the Requirements for the Degree Master of Science

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Βу

Jerry Burns Pace December 1978

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ABSTRACT

The material presented in this Thesis concerns two topics: the first is the design of a Microprocessor Development System and the second is the application of this system for developing a rather extensive programming example.

The Microprocessor Development System was designed around a Z-80 microprocessor. The system contains 8K of RAM, 12K of ROM, 4 serial I/O ports and room for 3 additional cards. A 2K monitor was implemented in ROM and a cross assembler was set up on a large mainframe HOST system. An I/O routine was written to allow the microprocessor system to converse directly with the HOST system. Programs could then be developed on the HOST system, assembled with the cross assembler and loaded directly into the microprocessor for debugging.

The programming example discussed is a program to emulate a multi-terminal network processor, a device which is used to multiplex several terminals on a timesharing system via a single modem line. Excellent results were obtained when using the HOST/Microprocessor combination for developing and testing programs for the microprocessor system.

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INTRODUCTION

The recent availibility of low-cost microprocessors has opened the door for many new and useful applications. This thesis will discuss the design of one of these microprocessor systems; specifically, a Z-80 microprocessor and the application of this unit as a network processor.

The work reported here divides naturally into two parts: Part I was the development of the hardware. This included purchasing and assembling the microprocessor, modifying the software monitor so that the microprocessor could communicate with a large HOST computer (HIS 66/60) and installing a cross assembler on the HOST to assemble programs for the microprocessor.

Part II involved choosing a development example which would illustrate the capabilities of the development system. It was decided to write a program which would use the development system to develop an emulator for a Remote Network Processor (HONEYWELL RCP 707). The network processor was chosen to demonstrate the ability of the microprocessor to do complex jobs with relatively inexpensive hardware.

PART I

MICROPROCESSOR DEVELOPMENT SYSTEM

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Chapter 1

HARDWARE

The microprocessor development system hardware is composed of two 7" by 9" printed circuit cards (the CPU card and the I/O card), a printed circuit CPU BUS Mother Board with provisions for 5 cards, and a multi-output power supply. Detailed wiring diagrams of the CPU cards can be found in Appendix A.

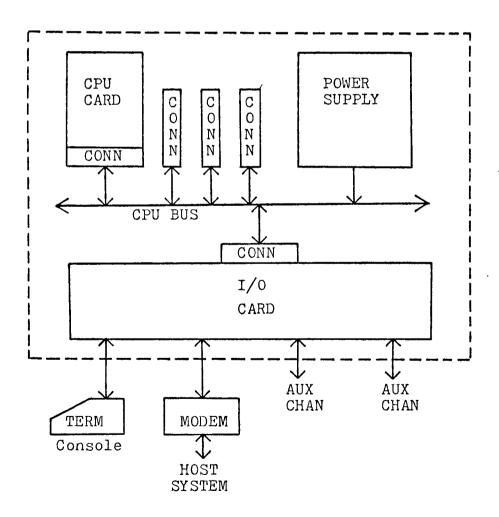


Fig. 1A

SYSTEM BLOCK DIAGRAM

The <u>CPU card</u> contains the Z-80 microprocessor, a 2 MHZ crystal clock, 8K of dynamic RAM, and 4K of programmable ROM, along with all the decoders, drivers, and receivers necessary to handle the CPU bus.

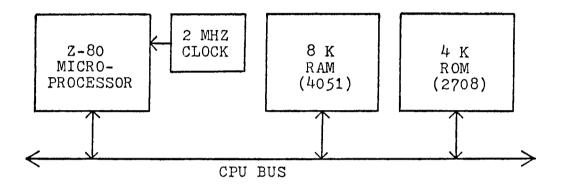


Fig. CPU CARD BLOCK DIAG.

The CPU chosen for this project was a Z-80 microprocessor. The reasons for choosing this particular unit were:

- It was one of the fastest and most powerful 8 bit microprocessors available.
- (2) It was very easy to design a system around.
- (3) Parts for this system were readily available and relatively inexpensive.

A crystal clock was used, instead of another type, due to its inherent stability and accuracy. The 2708 programable ROMs used to store the programs, both on the CPU card and the I/O card, combined a large storage capacity in a relatively small space, and were also very cost effective. The 4051 4K dynamic RAMs were used because, at this time, 4K dynamic RAMs were the least expensive type; and since the 2-80 had a built-in refresh counter, no extra hardware was required for refresh circurity.

The <u>I/O card</u> holds all the I/O interfaces consisting of 4 programmable I/O controllers (3 asynchronous and 1 synchronous unit), a 4 channel programmable real time clock/timer unit, an 8 input interrupt request register, and an additional 8K of programmable ROM.

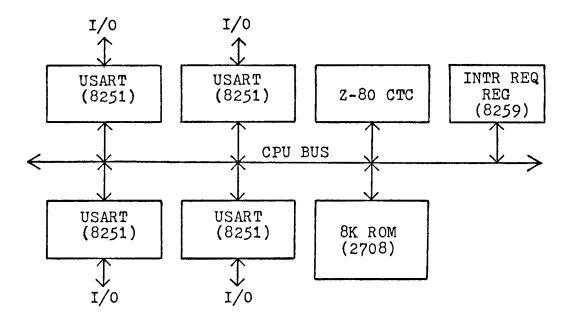


Fig. 1C I/O CARD BLOCK DIAG.

The programmable I/O controllers used on the I/O card were 8251 USARTS (<u>UNIVERSAL SYNCHRONOUS/ASYNCHRONOUS RECEIVER/</u> <u>TRANSMITTER</u>). These units were chosen for their ability to be programmed by the CPU, to operate in virtually any serial data transmission technique presently in use. They will operate in full duplex asynchronous mode to 9600 baud, and in full duplex synchronous mode to 50K baud. They also connect directly to the CPU bus and require no special interface circuitry.

The programmable clock/timer used on the I/O board was a Z80-CTC. It contains 4 independent programmable 8 bit counter/16 bit timer channels. Each channel can be programmed to operate either as a counter or a timer, which can generate interrupts and automatic interrupt vectoring with no external logic.

The 8259 interrupt controller is used here only as an interrupt request holding register. The software interrupt routine uses a polling technique to find the correct device to service.

The <u>CPU BUS Mother Board</u> serves simply as a mounting surface for five 100 pin card edge connectors which interconnect the signals and supply power to the cards. All CPU bus signals and all power lines are connected through this bus card.

The power supply is a three output regulated supply which produced 5 volts at 6 amps, +12 volts at 1.5 amps and -12 volts at 1.5 amps. Since a -5 volt supply was also required by the system, a -5 volt regulator driven by the -12 volt supply was used to supply -5 volts at 1 amp. The power supply delivers much more power than is required by the present system, which allows for the addition of other cards for future expansion of the basic system.

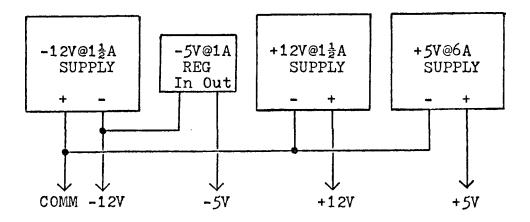


Fig. 1D POWER SUPPLY BLOCK DIAG.

Chapter 2

MONITOR PROGRAM

A system monitor program (ZAPPLE MONITOR by TDL), located on the CPU card in the upper 2K of ROM, provided all necessary functions for loading, displaying, modifing, and debugging assembly language programs. This monitor, however, had no provisions for connecting a HOST processor or for loading assembled code from a HOST processor. Therefore, an I/O routine had to be written to connect the Z-80 system to the HOST system (Appendix D).

The monitor contains, among others, routines for the following functions:

- (1) assign alternate peripherial devices for I/O or console,
- (2) display and/or change any single location in memory on the console,
- (3) display blocks of memory on the console,
- (4) fill blocks of memory with a single constant,
- (5) display and/or change registers from the console, and
- (6) set up one or two software break points.

Altogether, there are 23 separate functions in the standard monitor and provisions for 3 user defined functions. In addition, the monitor has many useful subroutines for I/O and data conversion which can be called by other programs. Appendix B gives a list of all the commands and a brief explaination of their use.

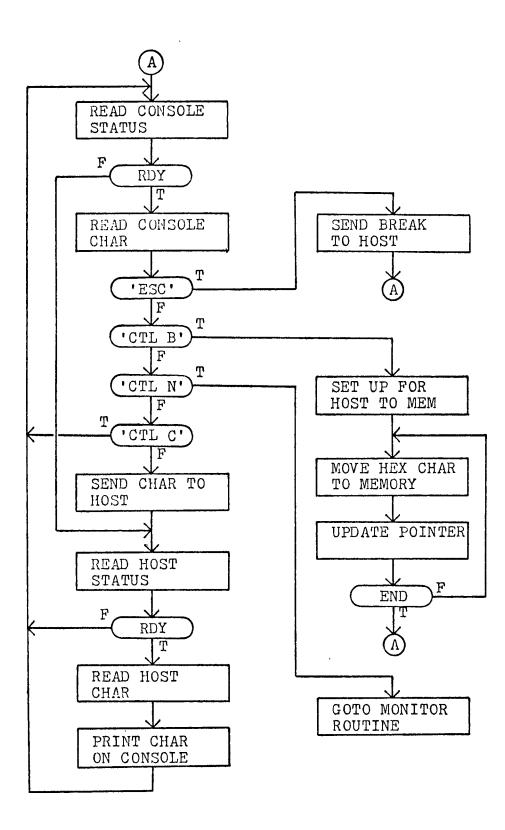


Fig. 2A Z-80 to HOST I/O Program

To make the task of conversing with the HOST processor as simple as possible, the Z-80 system was made to emulate a terminal and connected to the standard time sharing network (TSS) of the HOST processor. A special routine, which could be entered from this program, was written to load assembled code from the HOST processor to the memory of the Z-80 system. This made it possible to use the HOST system for writing, editing, assembling, and storing programs for the Z-80. The assembled code from these programs could then be down loaded to the Z-80 system for testing and debugging.

Figure 2A is a flow chart of the Z-80 to HOST I/O routine. The main loop of the program continously teststhe status of the console input and the modem input. When a character is ready to be read, the status flag will be set and the character will be read. If it is a console character, it will be tested and if it is also one of the command characters, a special routine will be entered to execute the command; otherwise, the character will be sent directly to the modem for transmission to the HOST. If, however, the character comes from the modem, it will immediately be printed on the console.

There are 4 command characters input from the console:

- (1) An 'ESC' character is used instead of the conventional break key because the I/O channel cannot detect a break. This input causes the program to go to the BREAK routine, which sends a break to the HOST for 250 MS and then returns to the main loop.
- (2) A 'CTL B' (CONTROL B) is used to enter the HOST to memory routine. This routine first asks for an offset value, next asks for the HOST file name, and then sends the command to the HOST to start input to memory. The input to this program, which must be a standard HEX FORMAT file — if not the programs aborts and returns to the main loop —, is then loaded and printed on the console at the same time. When the file is completely loaded, the routine returns to the main loop.
- (3) A 'CTL C' is normally used to cause an immediate disconnect. However, this was considered an undesirable feature. Therefore, this character is ignored and not sent to the HOST.
- (4) A 'CTL N' is used to cause a direct return to the monitor program.

Other than the above 4 characters, all keyboard characters are treated the same as in any standard TSS terminal and sent directly to the HOST.

Chapter 3

Z-80 CROSS ASSEMBLER

To allow Z-80 programs to be assembled on the HOST system, a cross assembler was acquired (XFOR-80 by MOSTEK). This cross assembler, although written in FORTRAN, was not written specifically for the HONEYWELL 66/60. Therefore, some slight modifications had to be made before it would work on this system. Some of the special characters had to be changed because they were not allowed on the TSS network and some special file instructions had to be added to the program.

The XFOR-80 is a 2 pass assembler which will assemble all standard Z80 source statement and also MACROS. As implemented on the HOST system (HIS 66/60), the input can be in the form of a TSS file, created on line, or a deck of punched cards or any other compatible file storage medium.

The output from the program is in the form of two separate disk files. One file is the line printer listing containing the assembled code along with the listing of the program instructions. It can be displayed on the console of the Z-80 system and/or printed on the line printer of the HOST system. The other file is the assembled code in standard HEX format which can be loaded into the Z-80 memory for execution or debugging.

PART II

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DEVELOPMENTAL EXAMPLE RNP EMULATION PROGRAM

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CHAPTER 4

RNP OVERVIEW

The <u>Remote Network Processor</u> is a device used for combining several terminals and/or several remote computers and/or remote batch facilities, in such a way that they can communicate with a host processor on a single high speed modem line. There are basically two protocols: one called RMC (REMOTE MESSAGE CONCENTRATION) for remote terminals and remote computers, and the other called RBS (REMOTE BATCH SYSTEM) for remote batch stations. It was decided to only implement the first, RMC, because it was simpler and would still serve well as an example. The following is a brief explanation of the RMC protocol. A more detailed explanation is available in the <u>HONEYWELL RNP/FNP</u> <u>INTERFACE</u> manual, number DB92.

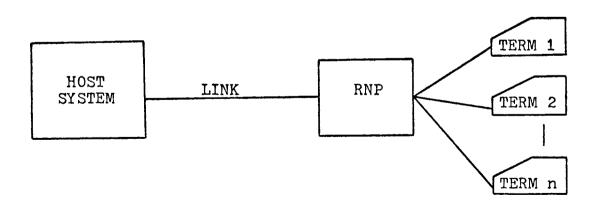


Fig. 4A RNP LOGICAL CONFIGURATION

A logical configuration of the RNP system is shown in Fig. 4A. The connection between the HOST and the RNP is called the link. This can be in the form of a modem or a direct wired connection.

The HOST/RNP link must be in one of the three following states:

(1) Physically disconnected.

(2) Logically disconnected (physically connected but idle).
(3) Logically connected (physically connected and active).
Control of the link is carried on through the exchange of Q-Frames during all periods in which the link is active and there is no link or logical messages to exchange.
This exchange is always initiated by the HOST, therefore, avoiding contention of the line.

All communication between the HOST and the RNP related only to the HOST/RNP link is carried in the <u>link message</u>, and communication between the HOST and each individual terminal is carried in the <u>logical message</u>. These messages, over the link, are carried in TRANSMISSION BLOCKS (XMIT BLOCK), which consists of a link message as the first, or only message. They may also contain one or more logical messages, each of which contains a unique address identifying its destination. The entire block is terminated by the 'EOT' character. Each XMIT BLOCK must be acknowledged (ACK) in the next received block or the same block is retransmitted (NAK). No new XMIT BLOCK (one having a new sequence code and different logical messages) may be sent until the previous one is acknowledged (ACK). The sequence code and the acknowledgement of the link message are used to insure against lost or duplicate XMIT BLOCKS. In addition, if no answer is received to a transmission within a specific amount of time, the same XMIT BLOCK is retransmitted.

There are two types of XMIT BLOCKS: the Service Message and the Data Message. The Service message differs from the Data message by the presents of only the link message and no logical messages in the XMIT BLOCK and the header of the link message contains a 1028 instead of a 1108 in the FC.

The Service message is used to control the link and conveys the following 4 messages:

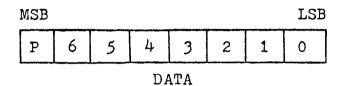
- (1) RFD Tells the receiver that the sender is going to disconnect the link. Must be acknowledged with an RFD.
- (2) DIS Tells the receiver that the sender is disconnecting the link. No reply is necessary and both processors disconnect.
- (3) A CALL Sent by the HOST to tell the RNP to accept all incoming calls.
- (4) N CALL Sent by the HOST to tell RNP to accept no new incoming calls.

The following table illustrates the error recovery rules. In these rules, the sequence code (SC) in the link message refers to the code in the header which alternates between 101₈ and 102₈. A changed SC indicates a new XMIT BLOCK. The ACK or NAK refers to whether or not a message is received in error. All retransmissions repeat the full XMIT BLOCK.

ERROR RECOVERY RULES

Same SC	New SC	ACK	NAK	ERR X	Transmit link message
	Х	x			with NAK and same SC. Process Logical messages. Change SC and transmit next XMIT BLOCK with ACK.
	Х		Х		Process Logical messages. Change SC and transmit next XMIT BLOCK with ACK.
х		x			Disregard Logical messages. Change SC and transmit next XMIT BLOCK with ACK.
х			х		Disregard Logical messages. Retransmit last XMIT BLOCK with same SC and ACK.

The XMIT BLOCKS are made up of messages which are composed of strings of characters. All characters used on the HOST/RNP link must be ASCII 8 bit (7 data bits + parity) characters. The bit notation is shown below:



A detailed description of the XMIT BLOCKS plus a description of the link and logical messages, along with a break down of their respective headers, is given in Appendix C.

Chapter 5

RNP EMULATION PROGRAM

The RNP Emulation Program has been written to simulate the actions of the HONEYWELL Remote Network Processor, configured to handle only the RMC protocol. It is written in modular form with individual subroutines for all major functions. This makes the program easily adaptable to many differing hardware configurations, without requiring major programming changes. Also, some of the functions are table driven to allow them to be changed or enlarged more easily. The program is also written such that it is 'ROMable' (i.e. written such that it can be stored in ROM and executed), therefore, no program variables are located within the program itself, but are all stored in RAM outside the program.

In addition to the normal RNP functions (HOST to terminal I/0), several additional functions were implemented. The program also allows device to device transmissions, without involving the HOST system; so this type of communication can go on even when the HOST is off line. Each device on the RNP has the ability to assign a destination device to itself with a keyboard command, and it can connect to or disconnect from the HOST with a similar command. This capability eliminates the need for a control console.

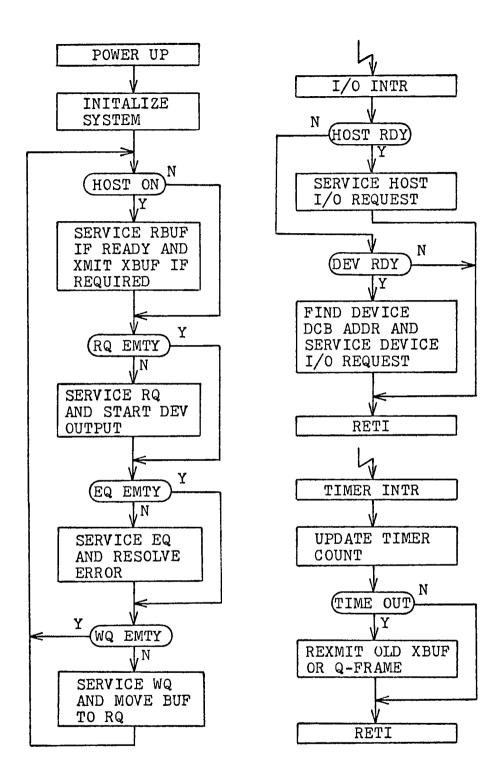


Fig. 5A RNP Basic Flow Diag.

The RNP program (Fig. 5A) accepts character input from terminals and/or remote computers (designated <u>devices</u>), and puts them into input buffers. Each completed buffer is then placed on a queue, to await output to the HOST or to another device. When the queue is services, each buffer is either output to another device or converted to a logical message. This logical message is then sent to the HOST in an XMIT BLOCK which normally contains logical messages from other devices.

Input from the HOST, in the form of an XMIT BLOCK, is converted from logical messages for several devices into individual output buffers, which are then placed on a queue. When this queue is services, the buffers are output to the respective devices by the I/O subroutines.

<u>Device I/O buffers</u> are composed of 64 byte buffer blocks. All available blocks are kept on the AQUE. When a block is needed by a process it is removed from the top of the AQUE, and when it is no longer needed it is put back on the bottom of the AQUE. The total number of blocks available is limited only by the amount of RAM available in the system.

All device I/O buffers are dynamic in size with a basic block size of 64 bytes, of which 58 are usable for character storage, 4 are used for the buffer header, and 2 are used for the linkage pointer. The basic buffer structure is shown in Figure 5B, along with the method of block linkage. Buffer size can vary from 1 to a maximum of 4 blocks for a total of 232 characters.

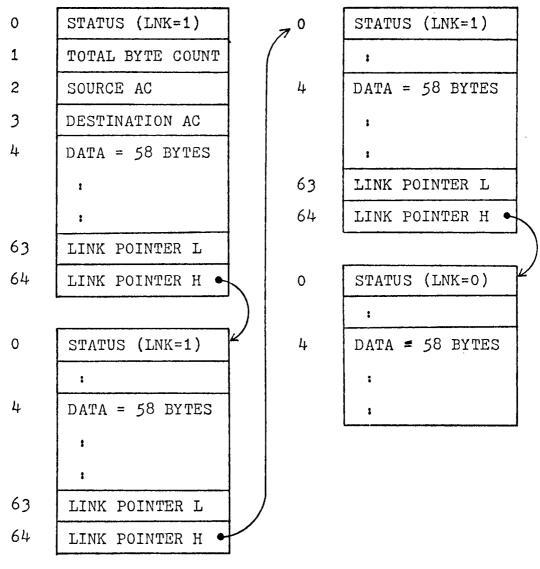


Fig. 5B MULTI-BLOCK BUFFER FORMAT AND LINKAGE

Blocks are linked up into a buffer by the use of a linkage pointer in the last 2 bytes of the block. There is also a bit in the status byte (link bit) which must be set to indicate that one block is linked to another. If the link bit is not set, it indicates this is the last block, or the only block, in the buffer

The buffer header (Fig. 5B) is only present in the first block of each buffer. It contains the buffer status byte as the lst byte, the total byte count of the buffer $(1 \le \text{count} \le 232)$ as the 2nd byte, the source address code of the buffer as the 3rd byte, and the destination address code as the 4th byte. The 5th through the 62nd byte is used for data storage and the 63rd and 64th hold the linkage pointer if necessary.

The status bytes also carries other information in addition to the link bit. The following is the bit arrangement of the status byte and the meaning of each bit.

	7	6	5	4	3	2	1	0
BUFs		LNK	ETB	PAR	HOST			

- LNK The Link bit indicates this is not the last block in the buffer and that a linkage address will be found as the end of the block.
- ETB The Extended Buffer bit indicates this is not the last buffer in this I/O operation (i.e. Input was more than 232 bytes).

- PAR The Parity bit indicates a character with bad parity occured somewhere within this buffer.
- HOST The HOST bit indicates this buffer's destination is the HOST system.

The last 3 bits above only occur in the lst status byte of each buffer (the buffer header), however, the link bit is in every status byte of every block in the buffer to indicate the presents or absence of another block.

The <u>HOST buffers</u> are a fixed size with 1024 bytes being allocated for each (Fig. 5C). There are two HOST buffers; one buffer (XBUF) for transmitting messages to the HOST, and one buffer (RBUF) for receiving HOST message transmissions. The buffer header for the HOST buffers consists of 4 bytes. The lst byte holds the status information, the 2nd and 3rd bytes hold the total byte count (6 = count = 1020) and the 4th is reserved for future designation.

0	STATUS
1	TOTAL BYTES L
2	TOTAL BYTES H
3	
4	DATA = 1020 BYTES
	:
	8

Fig. 5C

HOST BUFFER FORMAT

The status bytes for the HOST buffers are explained below:

		7	6	5	4	3	2	1	0	
RBUF	-	ERR	LOG	SVM	NSC	ACK	EMTY	FULL	BSY	
		••••••••••••••••••••••••••••••••••••••			•					
XBUF	-	ERR					EMTY	FULL	BSY	I
ERR	-	The Err has occ					or of s	ome ki	nd	
LOG	- The Logical message bit indicates there are logical messages present in the RBUF for processing.									
SVM	-	- The Service Message bit indicates this RBUF is a service message.								
NSC	-	The New SC bit indicates this RBUF has a different SC in the header than the last transmission.								
ACK	-	The Acknowledge bit indicates the HOST has acknow- ledged the reception of the last transmission.								
EMTY	Y - The Empty bit indicates there is no data in the buffer.									
FULL	ULL - The Full bit indicates no more data can be put into the buffer.									
BSY - The Busy bit indicates the buffer is now in the act of being changed by some process and cannot be accessed by another.								act		
Each	I/	0 device	e has a	issocia	ted wi	.th it	a <u>Devi</u>	<u>ce</u> Cor	trol	
Block	<u>c</u> (1	DCB) whi	.ch con	tains	all pe	ertiner	nt data	for t	hat	

device. All permanent data such as bus channel address and interrupt mask are stored here, and all temporary parameters such as byte count and buffer address used during I/Ooperations are also stored here. Figure 5D shows the format of the DCB used for device I/O.

0	STATUS						
1	INTR MASK						
2	DEVICE BUS ADDR						
3	AC (DEVICE LU#)						
4	DST AC (READ)						
5	BUF ADDR L						
6	BUF ADDR H						
7	REMAINING BYTES						
8	BLOCK STATUS						
9	TOTAL BYTES						
10	DATA POINTER L						
11	DATA POINTER H						
7 8 9 10	REMAINING BYTES BLOCK STATUS TOTAL BYTES DATA POINTER L						

Fig. 5D DEVICE CONTROL BLOCK FORMAT

The individual DCB bytes are as follows:

- (1) The device status holds the status bits for this device.
- (2) The <u>interrupt mask</u> is used for enabling/disabling the interrupt register.
- (3) The device bus address is the number of the I/O port.
- (4) The <u>address code</u> is the logical unit number of the device.
- (5) The <u>destination address code</u> is the logical unit number of the device.
- (6) The present <u>buffer address</u> is the address of the buffer assigned to this device during an I/O cycle.
- (7) The <u>remaining bytes</u> are the bytes left to be input to or output from this block.

- (8) The present <u>block status</u> is a save area for the status of the block in use.
- (9) The <u>total bytes</u> is the total remaining bytes left to be input to or output from the buffer.
- (10) The buffer <u>data pointer</u> is the pointer to the next byte of input or output.

The status byte of the DCB is explained below:

	7	6	5	4	3	2	1	0 ·
DCB	ERR	IDM	ETB	ONLN	ACK	XBF	RD	BSY

- ERR The Error bit indicates that an error has occured on an I/O transfer to this device.
- IDM The Identification Message bit indicates that an ID header should be attached to the start of each buffer sent by this device.
- ETB The Entended Buffer bit indicates that the present buffer is not the end of the message.
- ONLN The On Line bit indicates that this device is ready for I/O.
- ACK The Acknowledge bit indicates that the last transmission from this device to the host has been received.
- XBF The Xmit message buffer bit indicates that there is now a message waiting on the XQUE to be sent to the HOST (only 1 message is allowed to be on the XQUE from any single device at any particular time).
- RD The Read bit indicates that this is an input operation.
- BSY The Busy bit indicates that this device is performing an I/O operation and cannot start another until this one is complete.

The HOST has associated with it a <u>Host Control Block</u> (HCB) which contains all pertinent data for the HOST. Since the HOST I/O driver routines are not shared by any other devices, the HCB does not have to hold nearly as much information as the DCB. Figure 5E shows the format of the HCB.

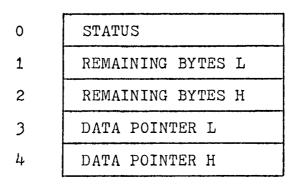


Fig. 5E HOST CONTROL BLOCK FORMAT

The individual bytes are as follows:

- (1) The Device Status byte holds the HOST status.
- (2) The Remaining Bytes indicates the total bytes left to be received or transmitted in this buffer.
- (3) The Data Pointer is the pointer to the next byte to be received or transmitted.

The status byte for the HCB is explained below:

	7	6	5	4	3	2	1	0
HCB	ERR	SYNC	ETB	ONLN	ACPT		RD	BSY

- ERR -The Error bit indicates that an error has occured on an I/O transfer to the HOST.
- SYNC-The Synchronized bit indicates that the HOST I/O unit has received the correct sync characters.

- ETB The Extended Buffer bit indicates the present buffer is not the end of the message.
- ONLN The On Line bit indicates the HOST is ready for I/O.
- ACPT The Accept all calls bit indicates the HOST will accept all new devices which sign on.
- RD The Read bit indicates this is a HOST receive operation.
- BSY The Busy bit indicates the HOST is performing I/O.

The DCB for any particular device is acquired through a table (DCBTAB). The AC of the device is all that is needed to calculate the offset for the table, which contains all of the DCBs for every device in the system. Figure 5F is an example of how this table is set up.

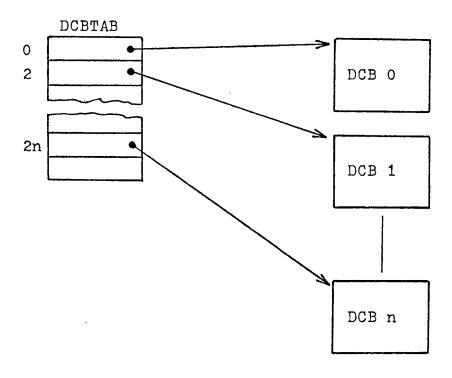


Fig. 5F DCB POINTER TABLE

All I/O within the program between two devices and between devices and the HOST is carried on through the I/O buffer. To prevent the possibility of interference between devices, all buffers are handled through queues on a first in first out basis. If a buffer taken from the top of a queue is destined for a device which is presently busy, it is put back on the bottom of the same queue to wait until the device is not busy. Figure 5G is a diagram of the queue buffer relationship.

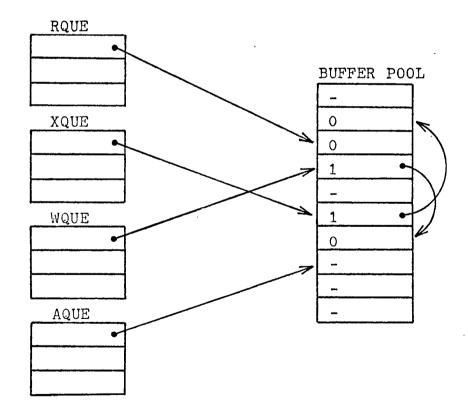


Fig. 5G QUEUE - BUFFER LINKAGE

The I/O is structured such that when a read or a write is started, for a particular device, the operation must run to completion before another can be started. This means that for a write operation, the complete buffer must be output, and for a read, either a line delimiter (CR) must be found, or the total length of input must exceed 232 characters. In the latter case, the present I/O buffer is terminated with an 'ETB' and put on a queue, and a new I/O buffer is started to receive the rest of the input.

There are 5 queues used by the program:

- (1) The Available blocks queue (AQUE) holds the addresses of all blocks not presently in use.
- (2) The Receive queue (RQUE) holds the address of all output buffers ready for output to a device.
- (3) The Transmit queue (XQUE) holds the address of all output buffers ready for output to the HOST.
- (4) The Write queue (WQUE) holds all overflow from the XQUE and the RQUE, and buffers which need to wait for output to a device or to the HOST.
- (5) The Error queue (EQUE) holds both buffer addresses and queue designators which are placed there when an error occurs in any other queue

Chapter 6

SUBROUTINE DESCRIPTIONS

MAIN_PROGRAM

The <u>Main</u> program loop (MAIN) is continously executed by the system until an interrupt occurs from one of the timers. If no device requires service at this time, then control is returned to this routine. During its execution the MAIN routine services each queue, if a buffer is on the RQUE it takes the buffer off the queue and starts output to the device indicated by the buffer, if a buffer is on the WQUE it moves this buffer to either the XQUE or the RQUE, and if a buffer is on the EQUE it moves the buffer to the required new queue.

MAIN also services the host receive buffer (RBUF) and the host transmit buffer (XBUF) when necessary. When the RBUF needs service, a flag causes the RBUF service routine to be entered which takes RBUF and converts all of its messages to output buffers and stores the addresses on the RQUE. And when the XBUF needs service, a flag causes the XBUF service routine to be entered, which takes all buffers off the XQUE and converts them to logical messages and puts them into the XBUF for transmission to the host processor.

The <u>Initalize</u> routine (INIT) is the routine which initalizes all necessary variables, initalizes all queues, sets up the stack pointer, sets up the I/O controller, starts an input cycle to all devices and enables the interrupts.

The <u>Clock and I/O</u> set up routine (CIO) is called by the INIT routine to set up the clock and timer interrupts, to set up the I/O device controllers, and to set up the priority encoder for device service requests.

COMMAND PROCESSOR

The <u>Command Processor</u> routine (CMDPRC) executes all system commands. It receives these commands from all devices in the form of buffers. The buffers are analyzed by the command processor and the appropriate action is taken. All commands start with a 'CTL C' followed by the command.

The commands are as follows:

- ASSIGN Consists of an 'A', followed by a one or two digit number. This command assigns the device, designated by the number, as the destination of the device issuing the command. All further inputs from the sorce device are routed to this destination device.
- ATTACH Consists of 'CTL A'. This command does two things: first it assigns the Host processor to the source device and second, it sends a select message to the Host to connect the source device to the Host.

DETACH - Consists of a 'D'. This command causes a detach message to be sent to the Host, which disconnects the source device from the Host, (Causes an immediate 'CP DISCONNECT')

The <u>Command Input</u> routine (CMDIN) is entered from DEVRD when a 'CTL C' is detected. It sets up a buffer to receive the command and pass it to the command processor routine.

DEVICE I/O ROUTINES

The <u>Device Input</u> routine (DEVIN) sets up the parameters in the device DCB for input from the device. It first acquires a block from the buffer pool and stores its address in the DCB. It then sets up the source and destination in the buffer header, initalizes the other parameters in the DCB for input from the device, and clears the read mask bit for this device in the I/O service mask.

The <u>Device Output</u> routine (DEVOUT) sets up the parameters in the device DCB for output to the device. On entry the buffer address is stored in the DCB. It then gets the byte count from the buffer and stores it in the DCB, initializes the other parameters in the DCB for output to the device, and clears the write mask bit for this device in the I/O service mask.

The <u>Interrupt</u> routine (INTR) is the device interrupt handler. This routine is entered once every millisecond from a timer interrupt. On entry it checks first for the Host needing service and then for any device needing service. If no service is needed, an exit is taken. However, if the Host needs service, the Host service routine is called. Additionally, if any device has a service request bit set and the device is not masked, the routine finds the correct DCB for that device, loads the registers with parameters from it, and then calls the service routine to service that device.

The <u>Device Read</u> routine (DEVRD), which is called by INTR, reads a character from the device designated by the parameters in the registers. It then stores this character in the designated buffer, updates the parameters, checks for the end of line character, checks for the last character in the present block or checks for a command character.

If the character read indicates the end of the line, then the buffer is closed and the device placed in idle mode. If the character read is the 1st in the present block, then the block is closed and a new one linked to the present one, and if the character read is a command character ('CTL C') then the present buffer is aborted and a command buffer is initiated by CMDIN.

The <u>Device Write</u> routine (DEVWR), which is called by INTR, writes a character to the device designated by the parameters in the registers. It then updates the parameters, checks for the end of the buffer, checks for the end of the present block, or checks for the extended buffer.

If this is the last character in the buffer, the routine restores the block to the buffer pool and places the device in the idle mode, if it is only the end of the present block it restores this block to the buffer pool and gets the address of the next block, and if it is an extended buffer, it restores this block to the buffer pool and then sets the device up to receive another output buffer.

HOST I/O

The <u>Start Receive Buffer</u> routine (STRBUF) is called by RBFSRV or GENXBF to start the next RBUF input cycle. It sets up the necessary parameters in the <u>Host Control Block</u> and unmasks the interrupt for input from the HOST.

The <u>Start Transmit Buffer</u> routine (STRXBF) is called by RBFSRV or GENXBF to start the next XBUF output cycle. It sets up the necessary parameters in the HCB and unmasks the interrupt for output to the HOST. The <u>Host Receive</u> routine (HOSTR) is entered from an interrupt, and if the Host I/O controller is in sync, a byte is read from the Host and placed in the RBUF. If it is an end of message character the full flag is set, the buffer is closed, and the interrupt mask set.

The <u>Host Transmit</u> routine (HOSTX) is entered fron an interrupt, and it transmits the next byte of XBUF to the Host. If it is the end of the buffer, the empty flag is set and the interrupt mask set.

HOST BUFFER CONVERSION

The <u>Receive Buffer Service</u> routine (RBFSRV) is called by MAIN to service the RBUF. It first strips the link message from the RBUF, analyzes the link message and uses it to set the RBUF status flags. Next, depending on the flag setting, this routine will retransmit the old XBUF, transmit a service message, generate and transmit a new XBUF or convert all logical messages in the RBUF to output buffers and put them on the RQUE. The routine then starts reception of the next RBUF.

The <u>Strip Link</u> routine (STRLNK) is called by RBFSRV to strip the link message off the RBUF. Depending on the data in the header, it will set or clear the flags in the RBUF status byte. The <u>Service Message</u> routine (SRVMSG) is called by RBFSRV to analyze RBUF service messages and either send back the correct acknowledge message or initiate the appropriate action.

The <u>Get Message</u> routine (GETMSG) is called by RBFSRV to get the next logical message off RBUF. It also calculates the length of the message and stores the address in a save area.

The <u>Get Buffer</u> routine (GETBUF) is called by RBFSRV to get a buffer to store the logical message. It gets enough blocks off the AQUE to hold the logical message and links them together as one buffer.

The <u>Convert Logical Messages</u> routine (CONMSG) is called by RBFSRV to convert the logical message to an output buffer and store the data in the buffer. It also puts the necessary header data into the buffer header.

The <u>Generate XBUF</u> routine (GENXBF) is called by RBFSRV to generate the next XBUF for output to the Host. It first generates a new link header for the XBUF using information from the present RBUF. It then gets buffers off the XQUE, converts them to logical messages and puts them into the XBUF. When the buffer is full or there are no more buffers on the XQUE, it starts transmission of the new XBUF.

The <u>Generate Link</u> routine (GENLNK) is called by GENXBF to generate a link message for the new XBUF. It analyzes the data in the present RBUF link header, generates the new link header and puts it in XBUF.

The <u>Convert To Logical Messages</u> routine (CONLOG) is called by GENXBF to convert input buffers to logical messages and put them into XBUF. It first generates a logical message header, puts it in the XBUF, and then moves the data from the input buffer to XBUF.

QUEUE HANDLERS

The queues are set up as simple linear lists with pointers to the top and bottom stored in the header. The header also holds the status byte, and the top and bottom buffer pointers

The <u>Put On Queue</u> routine (PUT) is called by many routines to put a value on a queue. It stores the two byte value, in the HL register, at the location pointed to by the top queue pointer. It then updates the pointer and, if the queue is full, sets the flag.

The <u>Get From Queue</u> routine (GET) is called by many routines to get a value from a queue. It gets the value pointed to by the bottom queue pointer and places that value in the HL register. It then updates the pointer and if the queue is empty, sets the flag.

CONCLUSION

This project has clearly shown the ability to produce a useful microprocessor development system with a bare minimum of hardware. The total cost of all hardware for this system came to less than \$1,000. The FORTRAN cross assembler, however, which normally sells for \$250.00, was donated by MOSTEK. And the cost of the console terminal is not considered because it was already owned by the university. But even including the prices of these items, this system still compares favorably with stand alone systems selling for up to \$10,000.

This system would be perfect for an application such as microprocessor training for a number of students. All programs could be written and assembled on normal timeshare terminals and then loaded into the development system for testing. The low price means several units could be accquired for the same price as one expensive stand alone system.

While the program discussed in the second part of the thesis was written and tested on the microprocessor development system, it was never completely tested with the HOST system. This was due to the lack of availability of an RNP line on the HOST system. And the effort necessary

to implement such a line on the HOST was too extensive to be completed in the allotted time. However, the software routines were all tested locally and worked well with test programs.

Although the emulator program should be very useful as it is, with a few changes to the harware and the software it could be made much more versatile and efficient. For example, the number of devices which it could service could be increased considerably by adding a different interrupt scheme and a DMA capability to the HOST I/O line, although as it stands, it could probably handle up to 16 low speed terminals (300 baud).

In addition, the emulator could be used with a different HOST machine just by changing the routines which determine the protocol. However, the basic framework of the emulator and most of the subroutines would remain the same as they are presently.

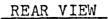
APPENDIX A

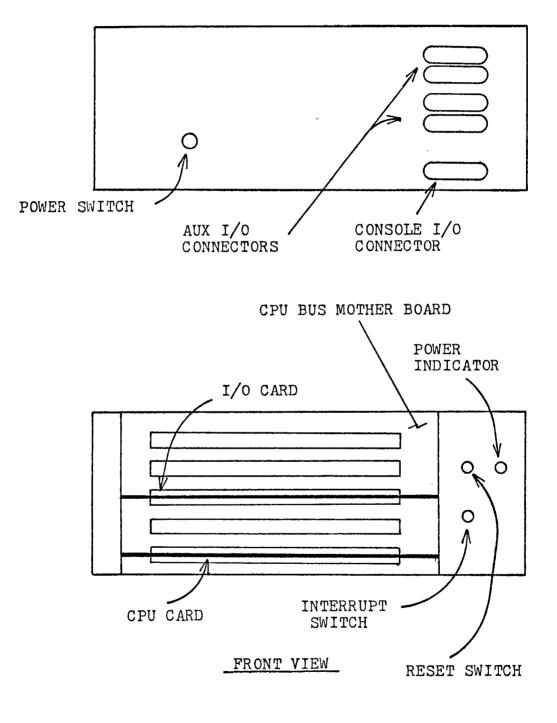
HARDWARE DIAGRAMS

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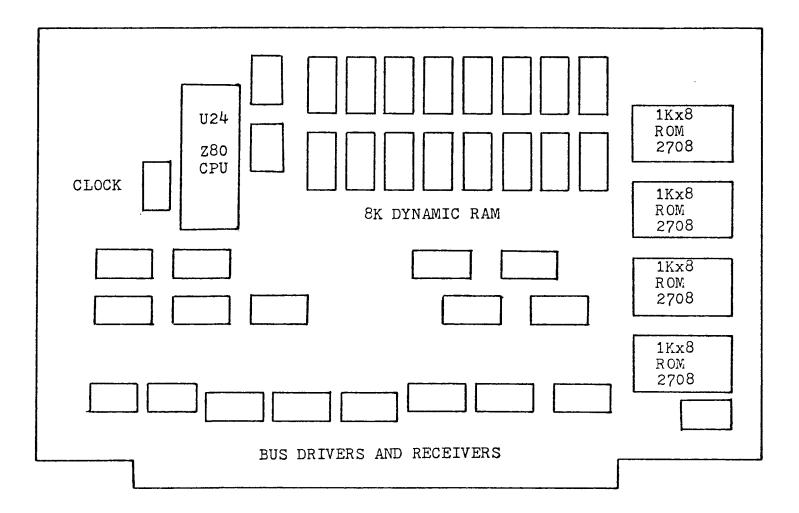
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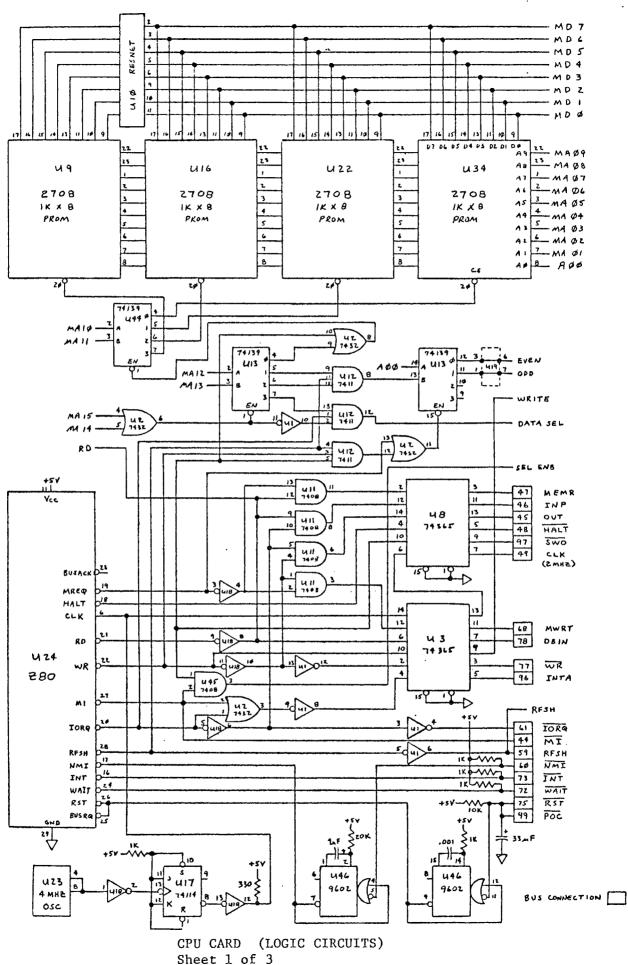
SYSTEM FRONT AND REAR VIEW



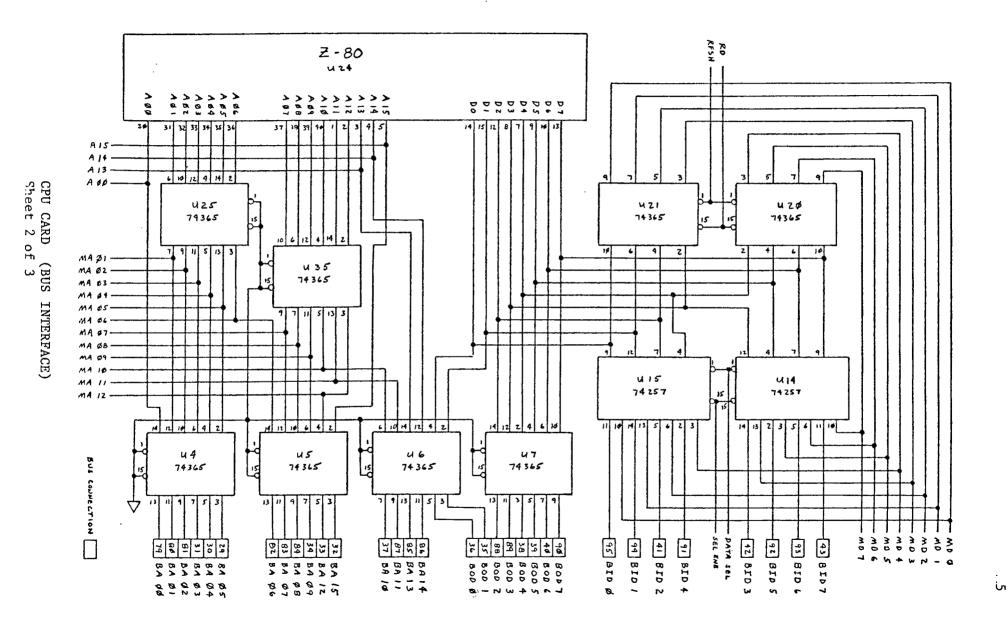
CPU CARD PHYSICAL LAYOUT

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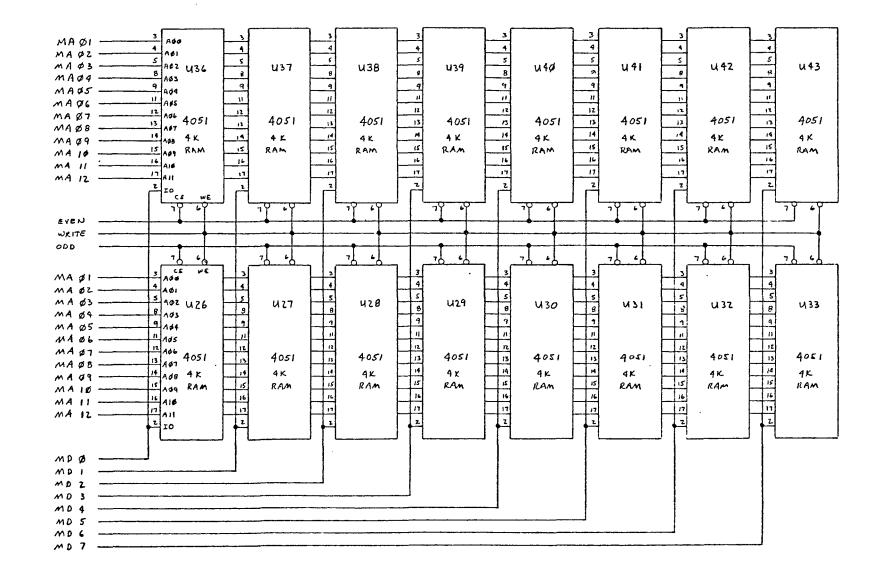
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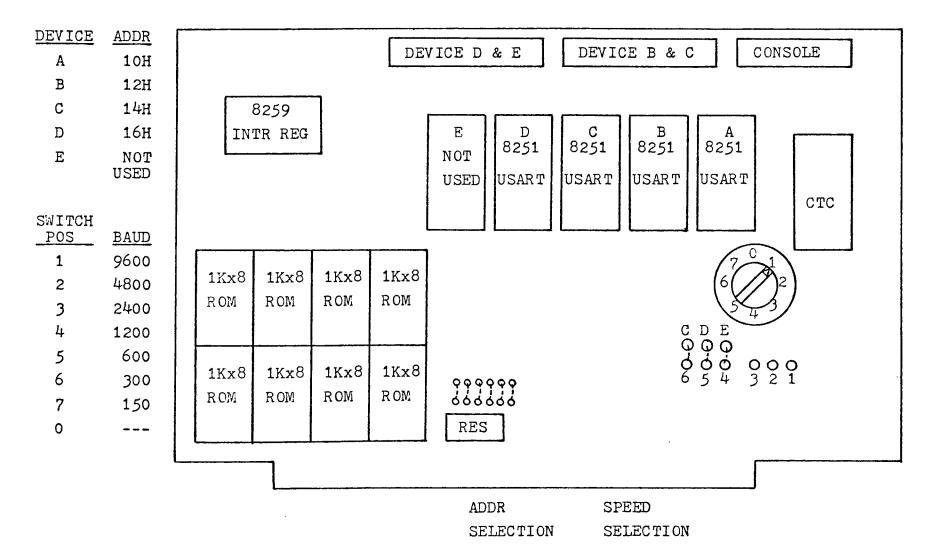


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CPU CARD (RAM STORAGE) Sheet 3 of 3

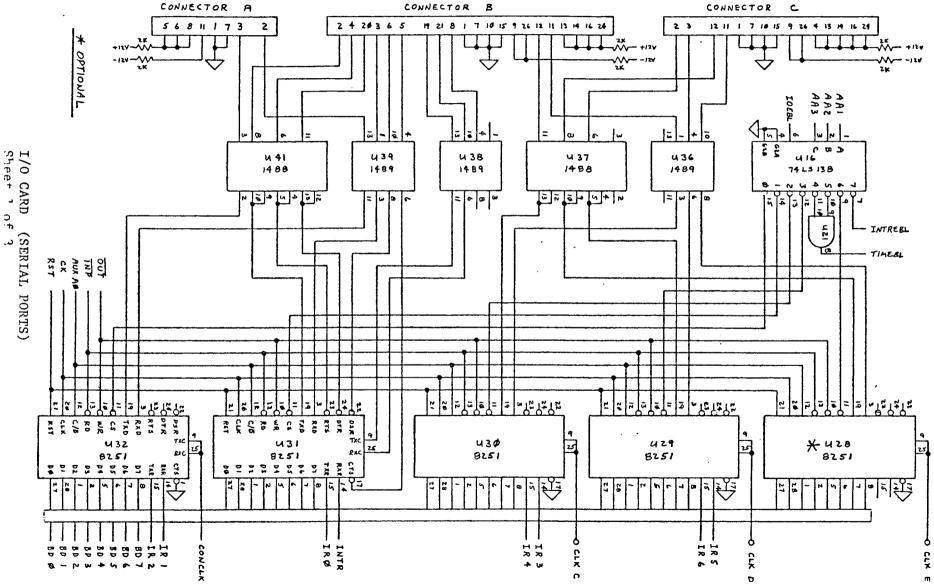
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I/O CARD PHYSICAL LAYOUT

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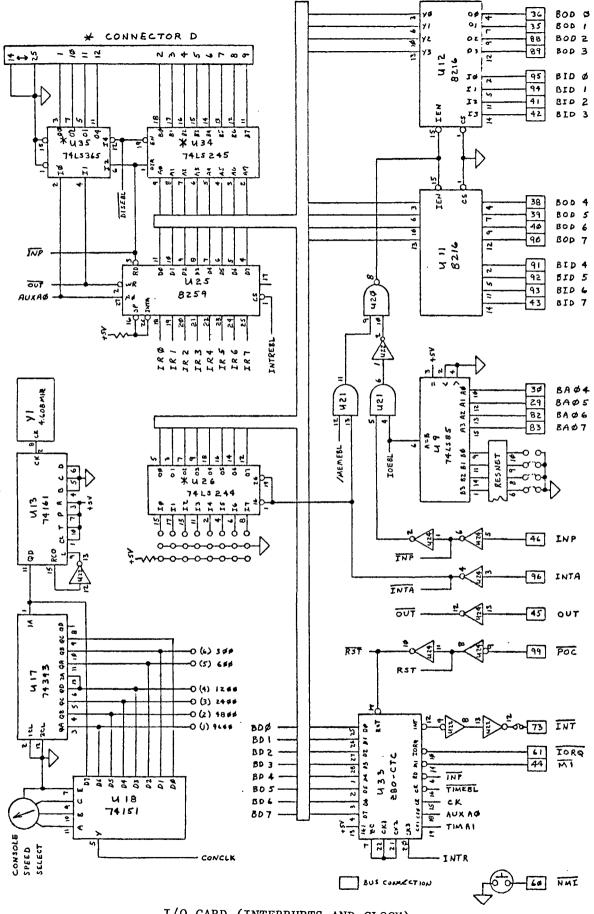


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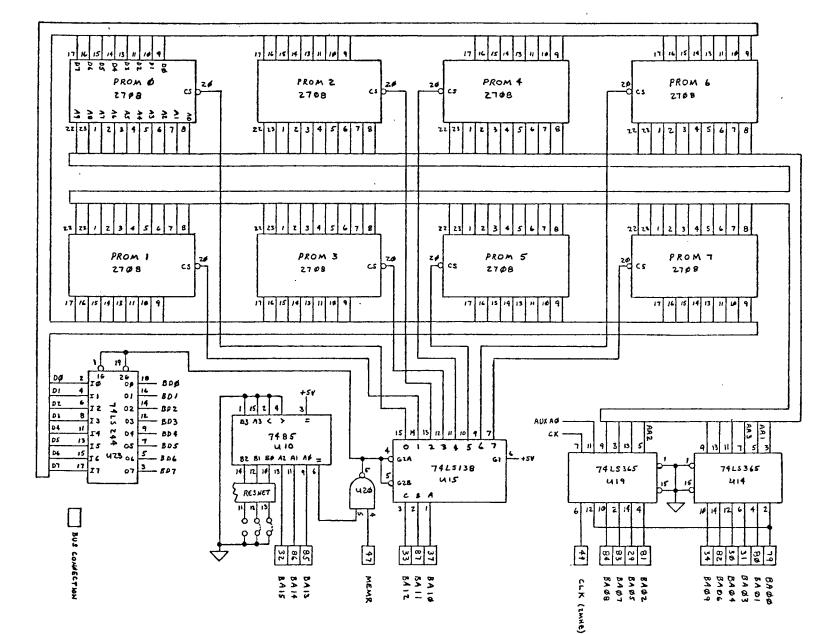
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I/O CARD (INTERRUPTS AND CLOCK) Sheet 2 of 3



I/O CARD (PROM STORAGE) Sheet 3 of 3

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APPENDIX B

ZAPPLE MONITOR COMMANDS

NOTE - The text in Appendix B was copied from the

ZAPPLE MONITOR OPERATIONS MANUAL

by

нн 199 Roger Amidon, Technical Design Labs

COMMANDS

The following is a list of commands for the Zapple Monitor. Precise definitions and usage notes are covered in the next section.

- A ASSIGN reader, punch, console or list device options from the console.
- B BYE (system shut down).
- C COMPARE the contents of memory with the reader input and display any differences.
- D DISPLAY the contents of any defined memory area in Hex.
- E END OF FILE statement generator.
- F FILL any define area of memory with a constant.
- G GOTO an address and execute. With breakpointing.
- H HEX MATH. Gives the sum and difference of two Hex numbers.
- I VERIFY ROM. Verifies contents of ROM against memory.
- J JUSTIFY MEMORY a non-destructive test for hard memory failures.
- K Jump to HOST I/O Routine.
- L LOAD a binary file.
- M MOVE a defined memory area to another starting address.
- N NULLS to the punch device.
- 0 PROGRAM ROM. Programs ROM from memory.
- P PUT ASCII characters into memory from the keyboard.
- Q QUERY I/O ports may output or input any value to or from any I/O port.
- R READ a Hex file. Performs checksum, relocating, offsetting, etc.
- S SUBSTITUTE and/or examine any value at any address (in hex).
- T TYPES the contents of a defined memory block in their ASCII equivalent.
- U UNLOAD a binary tape to the punch device.
- V VERIFY the contents of a defined memory block agains that of another block and display the differences.
- W WRITE a checksummed hex file to the punch device.
- X eXAMINE and/or modify any or all registers including the special Z-80 registers.
- Y "Yis there". Search memory for defined byte strings and display all addresses where they are found.
- Z "Z end". Locate and display the highest address in memory.

COMMAND SET USAGE

The following section lists the commands, and describes their format and their use. It should be noted that the Zapple Monitor recognizes both upper and lower case letters for its commands, and that in general, a command which is printing can be stopped with a CONTROL C, which is checked during a carrage return - line feed sequence. The following EXAMPLES show a comma (,) as a delimiter between parameters, however a space may also be used. If an error is made while inputting a command from the keyboard, it may be terminated by a rubout and the command re-typed. An asterisk is displayed indicating an ABORT of some kind.

COMMAND

DESCRIPTION

A ASSIGNMENT OF I/O DEVICE: The monitor system is capable of supporting up to 4 logical devices, these being: the CONSOLE, the READER, the PUNCH, and the LIST DEVICE. To these may be connected 4 different actual I/O devices, for a total of 16 direct combinations of I/O device and function. The specific permutations are:

LOGICAL DEVICE	ASSIGNED DEVICES
CONSOLE	TTY CRT BATCH USER (user defined)
READER	TTY CASSETTE PAPER (HIGH SPEED READER user written) USER (user defined)
PUNCH	TTY CASSETTE PAPER (HIGH SPEED PUNCH user written) USER (user defined)
LIST DEVICE	TTY CRT LINE PRINTER (user written) USER (user defined)

The default mode for each logical device is always the teleprinter.

Assignments are made using the following format;

EXAMPLE: AC =C(cr)

assigns the console equal to the Crt (video terminal) device. similarly:

EXAMPLE: AR = T(cr)

assigns the reader device to be the teleprinter.

While performing a command which requires a reader input (C,L,R), if the assigned reader is the Teleprinter, the software will look for a character from the TTY input. If a character is not recieved within a few seconds, it will ABORT, printing an asterisk (*) and return to the command mode. Similarly, if the assigned reader is the Cassette device, and you WISH to abort for some reason, changing the position of any of the SENSE switches will force an ABORT. On the external reader routines, returning with the carry set indicates an abort (or OUT OF DATA) conditinn.

When assigning a device, only the first letter initial of its name is required.

The Monitor itself is set-up to support the TTY, CRT, and Cassette routines. The other assignments require the addition of user's routines. These are addressed via the commands, which vector to starting addresses.

EXAMPLE; AL=L(cr)

assigns the list device to be the line printer. It vectors to (start address) +812H, or 12H above the end of the monitor. That would be the address for the line printer routine. For details of these arrangements, see the Source Documentation.

Within the above, the assign console equals batch "AC=B(cr)" deserves further mention. In BATCH mode, the READER is made the Keyboard input, and the LIST DEVICE is made the console output. This allows the running of a job directly from the reader input, with the result being output to the list device.

A typical use of this assignment would be the reconstruction of a lengthy text editing job where the text and your editing commands have all been saved on paper tape. With the BATCH MODE, you may assign the reader equals the TTY, the List device equals the TTY, and Console equals BATCH. Running the tape through the reader is the same as you redoing the entire text editing by hand, and the output will go to the TTY and be printed. On a very lengthy job, you could even start the process, and go away until it's done. Its usefulness is limited only by your imagination. B BYE. This command completely shuts down the system. It is useful where children might have access to the system, where a telephone communications link is established under remote control, or anytime when the operator wishes to make the system inaccessible to unauthorized use.

EX AMPLE:

R

completely kill the keyboard, Recovery from the shut-down is accomplished simply by inputting a CONTROL-SHIFT N from the keyboard. (ASCII equivalent is a Record Separator - "RS"; HEX character is a 1EH.) The monitor will sign on and print a greater-than sign (>), however, the register storage area will not be cleared.

C COMPARE the reader input with memory. This command is useful for verifying correct loads, verifying that a dumped tape matches with its source, etc.

EXAMPLE: C1000,2000(cr, start reader)

compares the memory block 1000H to 2000H with the input from the reader device.

For those with automatic readers, the operation is very simple. Assign the Reader equal to the device you wish to enter the data against, type C (starting address), (ending address)(cr), and the reader will start. The first character read by the reader will be the one matched with the starting address. If any discrepencies are encountered, the reader will stop, and the address (in hex) of the error will be printed on the display. The reader will restart, and continue in this fashion until the entire tape is compared.

If your reader cannot operate automatically, start the reader manually. If an error is encountered, however, while the incorrect address is being printed, the reader will continue, and get "out of sync" with the compare action. Therefore, it is necessary to manually stop the reader if an error is encountered, and manually reposition the tape to the byte following the error. (An excellent article on how to convert ASR33 type readers to automatic operation was recently presented in INTERFACE magazine.)

D DISPLAY memory contents. This command displays the contents of memory in Hex. Memory is displayed 16 bytes per line, with the starting address of the line given as the first piece of data on the line.

EXAMPLE: D100,1FF (cr)

will display in hex the values contained in the memory block 100H to 1FFH.

E END OF FILE. This command generates the end of file pattern for the checksum loader. It is used after punching a block of memory to the punch device using the "W" command. An address parameter for the end of file may be given if so desired.

EXAMPLE: E(cr)

will generate an "end of file marker".

EXAMPLE: E100(cr)

generates the EOF marker with the address parameter "100H". When loading such a file, upon completion, the address contained in the End of File will be placed in the "P" register. Execution of the program may then be initiated by typing "G(cr)".

F FILL command. This command fills a block of memory with a specific value. It is quite handy for initializing a block to a specific value (such as for tests, zeroing memory when starting up, etc.) *NOTE: Avoid doing this over the monitor's stack area. This area may be determined as being between the value you get when typing the Z command, and the value in the S register upon sign-on. It is approximately 60H bytes below the "Top of memory"(Z).

The format for the command is:

EXAMPLE: F100,1FF,FF

fills memory block 100H to 1FFH with the value FFH.

G

GOTO command. This command allows the user to cause the processor to GOTO an address and execute the program from that address. In the actural performing of the G command, a program, which has been placed in the stack area during the sign-on of the monitor, is executed. This program will first take all of the values in the register storage area (displayed with the X command), and stuff them in their correct registers in the CPU, and finally JMP to the program address being requested by the operator. If this short program up in the stack has been destroyed (as a result of a blow-up", or the F or M commands, etc.) the monitor will not be able to GO anywhere, and a manual restart of the monitor will be required. Whenever the monitor is restarted at the initialization point (first address I.E. OFOOOH), the contents of the registers are set to ZERO with the exception of the S (stack), which contains a valid stack address. This actual value depends on the amount of memory in the system, etc. In its simplest form, the letter "G" accompanied by a parameter causes the processor to go to that address and start execution.

EXAMPLE: G1000

would cause the processor to goto address 1000(H) and execute from that address.

Additionally, one or two breakpoints may be set.

EXAMPLE: G1000,1005,1010

would cause the program to start execution at address 1000H, and IN THE EVENT that the program gets to address 1005, OR 1010, the program will stop execution and return to the monitor, printing an "at" sign, and the address of the breakpoint that was executed. (ie. @1010). It then prints the ">"prompt, awaiting further instructions. This action also cancels any breakpoints previously set.

Breakpoints must be set at locations containing an instruction byte. This is a SOFTWARE breakpoint system, and requires either RAM at RST 7 (restart 7, addr. 0038H), or if using ROM, a permanent JMP to the monitor TRAP address (OFOIEH) at 0038H. Remember, this is a SOFTWARE breakpoint system, and the program being debugged must be in non-protected Read/Write memory.

EXAMPLE:	*C2 34	JNZ	1234H
	12		
	* 3 E 0D	MVI	A,CR
	*21	LXI	н,1000Н
	00 10		
	*77	MOV	M,A
	*23 *CD	INX	H
	78	CALL	5678H
	56		

The asterisks (*) mark the bytes that may be used as breakpoints.

H HEX MATH. This command allows the execution of hexidecimal arithmetic directly from the console. It will give the sum and difference of any two hex numbers entered.

EXAMPLE: H1000,1010(cr) 2010 FFF0

2010H being the sum, and FFFO being the difference of the two hex values.

J The J command is a non-destructive memory test. The command reads any given byte, complements it, writes into the location the complement, compares the complement with the accumulator, and rewrites the original byte into the location. The command is used with two parameters, delineating the block of memory to be checked.

EXAMPLE: J1000,1FFF

would perform the above test on the block 1000H to 1FFFH.

If errors are detected, the address at which the error is found and the error are displayed on the console before the test is continued.

EXAMPLE: J1000,1FFF(cr) 1F00 00001000

would indicate that the 4th bit (D3) at location 1F00H did not correctly complement itself.

This test is useful for the discovery of hard memory failures, and also serves as a quick check for accidentally protected memory. A fully protected memory block would print out as entirely "ls". (lllllll)

L LOAD BINARY FILE. This command loads a binary file from either a cassette or paper tape.

EXAMPLE: L1000(cr)

would load the tape at address 1000H. This would require that the program be an absolute program, designed for address 1000H. The start-of-file mark (automatically generated by the "U" command) is a series of 8 OFFH's (rubouts). When this is detected at the start of file, the bell will ring on the TTY to indicate the start of the load process. When the end-of-file is detected (again, a series of 8 rubouts) the load is terminated, and the address of the NEXT location that would have been loaded is printed on the console. There are two constraints on this type of file system. The middle of the program can not contain more than 6 OFFs (11111111) in a row (an unusual occurence), and if OFFH is the LAST data byte in the file, it will be ignored. This too is unusual, and only a minor inconvienience.

Binary programs loaded at other than their design address will not run. The "L" command does not perform checksum functions, and cannot handle relocatable files. This is a pure and simple byte-forbyte binary loader (see "U" command.)

MOVE COMMAND. This command is used to move a block of memory from one location to another. The original block is NOT affected by the move, remaining intact so long as the block moved into does not overlap with the block currently occupied. This command, like the "F" command should be used with some caution as moving a block into an area occupied by the stack, or the program or the monitor will cause unpredictable results.

EXAMPLE: M1000,1FFF,2000(cr)

moves the contents of memory contained in the block 1000H to 1FFFH to a starting address of 2000H. The new block has the limits 2000H to 2FFFH.

This command is very useful for working on programs without destroying the original, verifying blocks of memory loaded with existing memory, etc.

N NULL. This command punches nulls to the punch device. 72 nulls are punched whenever the command is used. It may be used repetitively for any desired leader length.

EXAMPLE: (N) *Note: the "N" or "n" will NOT echo, so as to not spoil the paper tape.

It will punch 72 nulls to the punch device.

P PUT ASCII characters into memory. This command allows ASCII characters to be written directly into memory. It is useful for placing labels in files, etc.

EXAMPLE: P1000(cr)

activates the command, and any further inputs via the keyboard would be placed into memory in their ASCII equivalent. The command is terminated by a CONTROL D character, w th the address of the location following

Μ

the last entry printed on the console (the Control-D is NOT stored). Recovery of the input data is affected by use of the "T" or "U" command.

Q QUERY INPUT/OUTPUT PORTS. This command allows any value to be output to any I/O port, and allows the value in binary on any I/O port to be read on the console.

EXAMPLE: Q01,7(cr)

would output an ASCII "7" to I/O PORT 1. (ASCII seven is a "bell" so on a TTY, the bell would ring.)

EXAMPLE: QI1(cr) 00001101

inputs the value at port 1, in the illustration above, we see that bits 0,2 and 3 are high, the others low. This is useful for observing the condition of status bits and other diagnostic activities.

R READ A CHECKSUMMED HEX FILE. This command reads checksummed hex files in the INTEL format, as well as being capable of loading the relocatable TDL files at any selected address and bias offset. When reading an ABSOLUTE file (INTEL format), there may be only a BIAS added. These files cannot be relocated. The format is: R(bias),(relocation)(cr).

If a checksum error or a failure to write the data to memory occurs, the loading process is stopped, an asterisk is printed (indicating some error condition), and the address that was attempting to be written will be displayed on the console device. This is to assist in determining the failure.

EXAMPLE: R(cr, start reader)

will load a hex file at its absolute address.

EXAMPLE: R,1000(cr, start reader)

will load a TDL relocatable hex file at address 1000H and modify the program to run at address 1000H.

EXAMPLE: R1000,100 (cr, start reader)

loads the file set up to run at 100H, but with a positive BIAS of 1000H added to it. Thus, the file, set up to run at 100H will be loaded at 1100H.

EXAMPLE: R1000(cr)

will load the file, set up to run at address 0000H, at address 1000. In other words, using the TDL relocating format, you may load any program, to execute anywhere in memory, anywhere in memory. (Think about it.....)

S SUBSTITUTE and examine. This command allows any address in memory to be examined directly, and allows substitution of one value for another at that address if desired.

EXAMPLE: SF810(sp)00-(sp)1A-(sp)C3-(sp)(cr)

In this case the "S" command examines address F810H. The hitting of the space bar (sp) displays the value of that address. (assuming value 00H at that address.) Hitting the space bar again displays the NEXT location in memory (F811H), and so forth. Simply typing S(sp) starts display from address 0000H. By repetitive typing of (sp), all of memory could be displayed one address at a time.

EXAMPLE: SF810(sp)00-(kb)FF(cr)

This command examines address F810H, showing the value 00H at that address. Immediately typing in FFH from the keyboard SUBSTITUTES FFH for 00H at that address. Repeating the example above would show:

EXAMPLE: SF810(sp)FF-

When an address is being examined, the address being examined may be moved BACKWORD by entering a backarrow (ba) or SHIFT-O, or underline, depending on the terminal used.

EXAMPLE: SF810(sp)00-(ba)AA-

shows that at address F80FH, the value AA exists. Typing a space bar will examine F810H again.

T TYPE ASCII characters from memory. This command allows the contents of memory to be displayed in their ASCII equivalents. All non-printing characters will be displayed as periods (.). It may be used to display the results of the 'P' command which allows keyboard entry of ASCII characters directly into memory. Also useful for finding text strings and messages in software. The initial address if first displayed, then the first 64 characters, the next address, etc. until the upper limit has been reached. EXAMPLE: T1000,2000(cr)

displays the ASCII equivalents of memory locations 1000H to 2000H. If the 'P' command had been used to place a 'message' into memory somewhere in that memory block, it would soon be apparent on the console display.

U UNLOAD BINARY. This command simply dumps core to the punch device. It may be used with a cassette system as well, with no start-up problems. It does not generate a checksum. The format which is generated will be a leader, eight OFFHs, binary data, eight OFFHs, and a trailer. The OFFHs are 'rubouts' and are called files ques. These are detected and counted to determine the start and the end of files.

EXAMPLE: UOO,FF (cr, start reader)

will generate a binary tape, formated as discribed above, of the values contained in memory locations OOH to FFH.

V VERIFY. This command allows the user to verify the contents of one memory block agains the contents of another memory block. This is very useful for functions such as verifying that a file generated from a program is a duplicate of the actual program, etc.

EXAMPLE: V1000,2000,3000

will compare the contents of the memory block 1000H to 2000H against the contents of the memory block commencing at 3000H and extending to 4000H. Any differences will be displayed.

EXAMPLE: V1000,2000,3000 100F 00 FF

indicated that the contents of address 100FH is a 00 while that at 300FH is an FF.

W WRITE Hex file. This command dumps memory to the punch device in the standard 'Intel-style' hex file format. Both start and end of file parameters are required. The proper 'end of file' (EOF) is generated by the 'E' command.

EXAMPLE: W00,FF(cr,start punch) (after punching) E(cr) will generate a checksummed hex file of the values in the memory block OOH to FFH. If the assigned punch and console are the same, the program will pause and wait for the operator to turn on the punch (ASR33, etc.). Use of the 'N' command at either the beginning and/or end of the file is optional, but recommended.

X eXAMINE REGISTERS. The "X" command allows the user to examine and/or modify all of the 280 registers.

A - Accumulator B,C,D,E,H,L - CPU REGISTERS M - Memory (pointed to by H & L) P - Program Counter (PC) S - Stack Pointer (SP) I - Interrupt Register X - Index (IX) Y - Index (IY) R - Refresh Register EXAMPLE: X(cr) displays the contents of MAIN registers A, B, C, D, E, F, H, L, M, P, S and I, in hex. EXAMPLE X'(cr) displays the contents of PRIME registers A, B, C, D, E, F, H, L, M, X, Y and R. Typing the letter "X" (or X'), followed by a specific register letter will display the contents of that register. Entering a new value via the keyboard (kb) will substitute the new value in the specific register. Hitting the space bar will display the next register in which you may then perform substitutions, etc. In the unique case of the "M" register, you may modify the 16 bit pointer (H&L) to that memory location.

EXAMPLE; XA 00-(kb)FF(cr) XA FF-(sp)00-(kb)FF(cr) XA FF-(sp)FF-(cr)

first examines the contents of register "A" (OOH), then substitutes an FF. In the next line, the FF is displayed, a space character displays the next register (again a OOH), and substitutes an FF for this value. The last line displays both registers as containing FFHs. Y SEARCH. This command allows unique byte strings, from one up to 255 bytes to be searched for in memory, and the addresses where they are found to be displayed. It is advisable to search for unique patterns rather than single bytes. The search operation may be stopped with a control-C.

EXAMPLE: YC3,21,F3,01(cr) 0081 00B2 0F08 >

indicates that the byte string (in hex) C3, 21, F3, 01, is found in memory at locations 0081H, 00B2H and 0F08H. This routine will search all 65-K of memory for a unique sequence of bytes in less that one second.

Z Z TOP OF MEMORY. This command locates and gives the highest address of available memory in your system.

EXAMPLE: Z 7FFF >

indicates that the highest available memory is at address 7FFFH. Note that NO carriage return is required. Also, if only one 1K board were in the system, and it was addressed to have its top byte at address 7FFFH, the Z command would so indicate regardless of the absence of lower memory.

Additional Functions

The following functions are not part of the original Zapple Monitor. One of the functions,(K), cause a jump to the HOST I/O routine. The other two, (I and O), are used to program and verify PROMS using an auxillary PROM programming card.

- K JUMP TO HOST I/O. EXAMPLE: K will go to the HOST I/O routine and start execution.
- VERIFY ROM. This command will verify that a 1K block of data has been correctly written into ROM by the '0' function.
 EXAMPLE: I1000 will compare the program in the ROM with the data in the block from 1000H to 1400H. Any differences will be displayed.

EXAMPLE: 11000

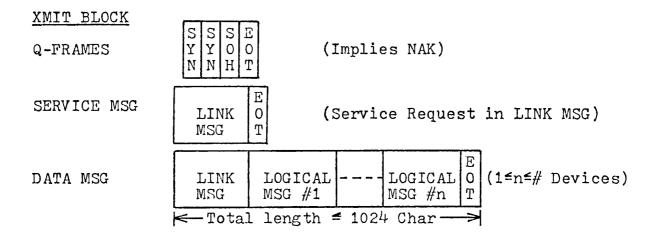
10F0 00FF

indicates that the contents of location FOH in the ROM is OOH, while that at 10FOH in memory is FFH.

- O PROGRAM ROM. This command will write data into a PROM from a 1K block of memory.
 EXAMPLE: 01000
 writes the data from 1000H to 1400H into the PROM.
- NOTE The above two commands (I and O) require a special card which contains a ROM programmer for 1K by 8 programmable ROMs.

Appendix C

RMC MESSAGE FORMATS



LINK MSC

S Y N	S Y N	S Y N	S Y N	S O H	F C	SC	A C	0 C	I C	S T X	E T X	B C C	(BCC = Block) (Check) (Char)	
	- Header ->													

LINK MSG INFORMATION

- SYN All XMIT BLOCKS must start with 2 or more SYN characters. SYN characters may also appear anywhere within the XMIT BLOCK but are disregarded.
- SOH Start Of Header character.

FC	- Format Code	1108	Transmission ACK/NAK.
		1028	Service message (RFD, DIS).
SC	- Sequence Code	101 ₈ ,102 ₈	Alternates on each new XMIT BLOCK.
AC	- Address Code	1008	Alway this value in single
			RNP networks

OC - Operation Code 1xy8	x = 0 ACK 1 NAK
	<pre>y= 0 No Instruction 3 A Call (accept all calls) 4 Ready for Disconnect (RFD) 5 N Call (accept no calls) 6 Disconnect (DIS) 7 Reserved</pre>
IC - Identification Code 1xx8	xx= # of messages in XMIT BLOCK. xx≤ 63
STX - <u>S</u> tart <u>Of T</u> ext character	Text may follow this character, but not normally found in LINK.
ETX - End Of Text character	Follows text, if present.
BCC - <u>B</u> lock <u>C</u> heck <u>C</u> haracter	X-OR of all characters from SOH to ETX not including SOH or any SYN characters in msg.

LOGICAL MSG

S O H	F C	S C	A C	C 1	C 2	I C	S T X	TEXT 0-232 Char	E T X	B C C
Ĭ	- 1	lea	ade	er		>				

LOGICAL MSG INFORMATION

SOH ·	- Start of Header		
FC ·	- Format Code	Bits 5-6=1	Bits 0-4 Indicate Mode
SC ·	- Sequence Code	101 ₈ ,102 ₈	Set but not checked.
AC ·	- Address Code	1xx8	xx= Destination Device #
001 ·	- Operation Code 1	1018	Indicates OC2 to be used.

002 -	Operation Code 2	1xy ₈	x= 0 1 2 7	Break ACK DIS ACK
			y= 0 1 2 3 4 5	Break DIS SELECT (Connect Term) Bad Parity
IC -	Identification Cod	1018 1108 1118 1128	Remo TTY VIP	te Computer 110 baud 150 baud 300 baud
STX -	Start Of Text			of logical message ows.
TEXT-	the following spec	ial char	racte	CII characters excluding rs: NAK, ENQ, US, DLE, EOT)
ETX -	End Of Text		ETB	ows text, can also be character for messages h are longer than 232 char.
BCC -	Block Check Charac	eter		as for LINK MSG including text characters.

:1/O PROGRAM FOR MODEM TO Z80 SYSTEM **3300H** ZEF: EQU ;MAIN MONITOR PROGRAM TRAP: EQU **381EH** ;TRAP RETURN ADDR L FADR: EQU **JC70H** PRINT CRLF & HL **;**GETS 2 BYTE PARAMETER 0033H EXPR1: EQU 101BH EQU ; I/O STATUS BYTE IOSTS: JA28H **;**TAPE READ ROUTINE LODO: FQU 0C78H C O : EQU CONSOLE OUT ROUTINE CONSOLE IN ROUTINE CI: EQU **JE13H** FILF: EQU 1020H FILE NAME BUFFFR TOM1: EQU 0C40H #MESSAGE PRINT ROUTINE CR & LF TO CONSOLE EQU 0004H CRLF: ; CALL INIT ;RESET I/O ENTER: LD A, ODH ; C R (FILE+DAH), A ; PUT AT END OF FILE BUFF LD CALL CRLF CR & LF LD 9,05H ; CHAR COUNT HL,MSGO **;**PROGRAM NAME LD CALL TOM1 PRINT IT ; C, '+' START: LD **;**PROMPT CHAR ; PRINT IT CALL CO INCO: IN A. (11+) FREAD CONSOLE STATUS AND 72 ; INPUT DATA JR Z, INM-S INO, CHECK MODEM CALL CI ;YES, READ CHAR ; CP 1 3 H ; = ESC;YES, SEND BREAK JR Z, CONBRK-\$ (P = 'CNTL-B'02H JR Z,HOSTLD-\$;YES, LOAD FROM MODEM ;= 'CNTL-C' CP 03H JR Z, INM-S ; IGNORE IT ;= 'CNTL-N' CР 0 E H RET Ζ ; GO TO MONITOR ; OUTM: CALL DUTMOD ;OUTPUT TO MODEM ; INM: IN A (15H) FREAD MODEM STATUS AND <u> 2</u> FINPUT DATA RDY JR Z, INCO-\$;NO, CHECK CONSOLE CALL INMOD FYES, READ CHAR LD C A ; CALL CO **;**PRINT ON CONSOLE INCO-\$ J R ; INMOD: IN A, (15H) FREAD MODEM STATUS AND 25 JINPUT DATA RDY JR Z, INMOD-\$;NO, WAIT ΪN A, (14H) ;YES, READ CHAR AND 7 F H **;**MASK PARITY RET

; OUTMOD: PUSH AF SAVE A A, (15H) ;READ MODEM STATUS IN XMIT RDY 01 AND Z.OUTMOD+1-\$;NO. WAIT JR POP AF ;YES, GET DATA (14H),A ;SEND DATA OUT RET ; SEND BREAK TO MODEM CONBRK: CALL BREAK START-\$ JR ; BRFAK: ;LOAD BREAK CODE LD A . 3DH (15H)/A ;OUTPUT TO MODEM USART (8251) OUT LD 005 A A ;WAIT COUNT (200 MS) CALL JAIT ;WAIT... ;LOAD NORMAL OPERATION CODE A . 35H LD (15H)/A **;OUTPUT TO USART** OUT LD C ; PRINT < ON CONSOLE СЭ JP. ; ; WAIT: B,090H **;**DELAY COUNT LD DJNZ WAIT+2-5 ;WAIT 1 MS DEC ; DEC WAIT COUNT, = 0A JR NZ .WAIT-S ;NO, DELAY AGAIN RET **;**YES, RETURN ; FREAD CHAR R D M O D : CALL INMOD PUSH AF SAVE AF ;SAVE BC PUSH BC LD C / A **;**PRINT ON CONSOLE CALL CO **;**RESTORE BC POP 9 C POP AF FRESTORE AF 0 R **JCLEAR CARRY** A RET ; HOSTLD: C,'>' **;**PROMPT CHAR LD CALL CO **;**PRINT ON CONSOLE ;GET BIAS, IF ANY CALL EXPR1 A , B ;LOOK AT DELIMITER LD SUB ; IF = CR) D H LD 8 . A ; RELOCATION = 0 LD C 🖌 A ; POP DE ;DE=BIAS Z, R0-5 ;CR ENTERED JR CALL EXPRI JGET RELOCATION POP ; BC=RELOCATION 30 ΕX R () : DE,HL ;HL'=BIAS, BC'=RELOCATION EXX CALL CRLF ; LD B,OAH CHAR COUNT HL,MSG2 LD LOAD POINTER

н1:	LD LD CALL LD CP JR CP	B, OAH CI (HL), A ODH Z, H2-\$ 13H	<pre>; PRINT MSG2 ; FILE NAME BUFFER ; MAX LENGTH ; READ NAME ; PUT IN BUF ; CARRIAGE RET ; YES, LAST CHAR ; = ESC</pre>
Н2:	JR INC DJNZ CALL LD CALL LD LD	Z, CONBRK-\$ HL H1-\$ CRLF B,5 HL,MSG1 MOUT B,08H	;YES, SEND BREAK ;INC POINTER ;CHECK MAX LNGTH ;CRLF TO CONSOLE ;CHAR COUNT ;LOAD POINTER ;OUTPUT MESSAGE TO MODEM ;MAX FILE LENGTH ;FILE NAME POINTER ;OUTPUT TO MODEM
;	0 R L D L D L D L D C A L L' P O P	AF OF 3H O4H (IOSTS) A A OC 3H (1006H) A HL & RDMOD (1007H) & HL L ODO AF	;GET I/O STATUS ;SAVE OLD STATUS ;CLEAR READER STATUS ;SET NEW STATUS ;PUT IN IOBYT ;JUMP CODE ;PLACE IN USER ROUTINE AREA ;MODEM I/O DRIVER ;PLACE IN USER ROUTINE AREA ;GO READ TAPE ;RECALL I/O STATUS ;RESTORE OLD STATUS
MOUT:	CALL LD CALL INC CP JR	OUTMOD C . A	;GET CHAR FROM BUFFER ;OUTPUT TO MODEM ; ;PRINT IT ;INC POINTER ;= CARRIAGE RET ;NO, CHECK COUNT ;= MAX COUNT
M01:;	RFT		
M S G O :		"ZIO" DADDH	CR & LF
M S G 1 : M S G 7 :	D E F M D E F M	'LIST ' 'FILE' ' NAM'	MESSAGE TO HIS
;;		IP I/O CHANNEL	. S
; INIT:	X O R L D	А ВС,4011н	;CLEAR ACC ;RESET USART, I/O CHAN(10,11)

	LD	DE, OFA35H	SET MODE & FUNCTIONS
	CALL	USET	;OUTPUT COMMANDS
	LD	C,15H	; I/O CHAN(14,15)
	CALL	USET	;OUTPUT COMMANDS
	LD	С,17н	;I/O CHAN(16,17)
	CALL	USET	;OUTPUT COMMANDS
	RET		;
;			
USET:	OUT	(C)•A	CLEAR COMMAND REG
	OUT	(C),B	FRESET USART
	OUT	(C),D	SEND MODE COMMAND
	OUT	(C),E	SEND FUNCTION COMMAND
	RET		;
;			
	END	000H	

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APPENDIX E RNP PROGRAM LISTINGS

	PROGRAM	F UNC T I ON	PAGE
001	C I O	CLOCK & I/O INIT	E 2 9
002	CMDIN	COMMAND INPUT	E 2 5
003	CMDPRC	COMMAND PROCESSOR	E 2 5
004	CONLOG	CONVERT BUFS TO LOG MSG	E12
005	CONMSG	CONVERT LOG MSGS TO BUFS	E 8
006	DEVIN	DEVICE INPUT SET UP	E16
007	DEVOUT	DEVICE OUTPUT SET UP	E14
008	DEVRD	DEVICE READ ROUTINE	E 2 1
009	DEVWR	DEVICE WRITE ROUTINE	E19
010	ERROUT	ERROR PRINT OUT	E27
011	GENLNK	GENERATE LINK MSG	E11
012	GENXBF	GENERATE NEW XBUF	E11
013	GET	GET VALUE OFF QUEUE	E33
014	GETBUF	GET BUFS FOR LOG MSG STORAGE	E 7
015	GETMSG	GET NEXT MSG FROM RBUF	E 7
016	HBUFS	HOST BUFFER STORAGE	E34
017	HOSTR	HOST READ ROUTINE	E 2 3
018	HOSTW	HOST WRITE ROUTINE	E 2 4
019	INIT	INITALIZE ROUTINE	E29
020	INTR	DEVICE INTERRUPT SERVICE	E17
021	MAIN	MAIN PROGRAM LOOP	E 2
022	PUT	PUT A VALUE ON A QUEUE	E32
023	RBFSRV	SERVICE RBUF	E 2
024	SRVMSG	ANALYZE SERVICE MSG	Ε6
025	STRBUF	START INPUT TO RBUF	E23
026	STRG	PROGRAM VARIABLES STORAGE	E35
027	STRLNK	STRIP OFF RBUF LINK MSG	E 4
028	SUBS	GENERAL USE SUBROUTINES	E31
029	QUES	QUEUE STORAGE AREA	E35

; FILE MAIN MAIN RNP SERVICE ROUTINE CALL INIT ;INITALIZE SYSTEM START: ; MAIN: LD A, (HCB) GET RNP STATUS BIT ONLN.A RNP ON LINE JR Z.CHKRQ-\$ **;NO, CHECK RQ FOR OUTPUT** ; CALL RBFSRV **;YES**, SERVICE RBUF ; ; IX = RQUE CHKRQ: LD IX, RQUE CALL GET GET ADDR OFF RQ C,CHKEQ-\$ JIF EMTY, CHECK EQ JR CALL DSTDCB ;HL=DST DCB, IY=BUF ADDR CALL DEVOUT START BUF OUTPUT JR NC, CHKEQ-\$; DCB(BSY)=1 NO, CHECK EQ ; PUSH IY YES, MOVE BUF ADDR POP HL ;TO HL ; PUT ON RQ CALL PUT ; LD IX, EQUE CHKEQ: ;IX=EQUE CALL GET GET ADDR OFF EQ ;IF EMTY, CHECK WQ C,CHKWQ-S JR ;IX=QUEUE ADDR, HL=BUF ADDR POP IX CALL PUT ; PUT BUF ON Q JR NC, CHKWQ-S JIF OK, CHECK WQ JELSE PUT BACK ON EQ CALL FULERR ; LD IX,WQUE CHKWQ: ;IX=WQUE CALL GET ;GET ADDR OFF WQ C,MAIN-S FIF WQ EMTY, GO TO MAIN JR ;HL=DST DCB, IY=BUF ADDR CALL DSTDCB BIT ACK (HL) ; DEV(ACK) = 1 PUSH IY MOVE BUF ADDR POP HL ;TO HL J R Z,RSWQ-\$;NO, PUTBACK ON WQ LD IX, RQUE YES, PUT ON RQ CALL PUT ; PUT ON Q RSWQ: NC / MAIN-\$;NOT FULL, BACK TO MAIN JR LD IX,WQUE FULL, PUT BACK ON WQ JR R S R Q - \$; ; :; FILE RBFSRV ; HOST RECEIVE BUFFER SERVICE ROUTINE ; ENTER-; EXIT- RBUF SERVICED, IF ERROR-CARRY SET RBFSRV: LD HL, RBUF ; CHECK NEXT RBUF

;	RET		;RBUF(BSY)=1 ;YES, RETURN ;NO, RBUF(FULL)=1 ;NO, CHECK RBUF(EMTY)
	BIT	STRLNK ERR/(HL) Z/CKSV-\$;YES, STRIP OFF LINK MSG ;RBUF(ERR)=1 ;NO, CHECK SERV MSG
; R E R R :	BIT JR LD BIT	NZ,CKNAK-\$ IX,SBUF ACTV,(IX+0)	;XBUF(ACTV)=1 ;YES, GO CHECK NAK
;	JK		INUT SEND NAK M30
	JR SET LD		;OC(NAK)=1 ;YES, GO START BUF ;NO, SET OC(NAK)=1 ;RECALC BCC ;
; STRBF:	PUSH Pop	IX HL STRXBF	;SAVE RBUF ADDR ;MOVE BUF ADDR TO HL ; ;REXMIT OLD XBUF ;RESTORE RBUF ADDR
	JR	SRBF - \$	START NEXT RBUF
;			
NKMG:	C A L L J R	N A K M S G S R B F - \$;SEND NAK MSG ;START NEXT RBUF
скsv:	JR BIT JR BIT JR PUSH	LOG,(HL) NZ,CKSC-\$ ACK,(HL) Z,RXMT-\$ HL GENXBF	;YES, SEND SERVICE MSG ;RBUF(LOG)=1 ;YES, CHECK SC
;	U K		JIANI NEXT ROOT
RBMT:	BIT JR JR		;RBUF(EMTY)=1 ;YES, START NEXT RBUF ;NO, CONTINUE PROCESSING RBUF
; Srmg:			XMIT SERVICE MSG IF CARRY, ERROR GO START NEXT RBUF
; CKSC:			;RBUF(NSC)=1 ;YES, PROCESS RBUF ;RBUF(ACK)=1

;YES, SEND ACK MSG JR NZ / AKMG - S RXMT: CALL STRXBF **FREXMIT OLD XBUF** SRBF-\$ GO START NEXT RBUF JR PUSH HL ;(SP)=RBUF NXMSG: NXM: GET NEXT MSG CALL GETMSG JR NZ RERR-\$ **;IF NZ, ERROR** C.PTRQ-S **;BUF SAVED**, IF CARRY J R CALL GETBUF ;GET NEXT AVAILABLE BUF **;IF EMPTY, SAVE MSG** J.R C.SVMG-\$ PTRQ: LD IX, RQUE ;IX=RQUE CALL PUT FUT ADDR ON RQ ; IF FULL, SAVE MSG JR C.SVMG-\$ CONVERT TO OBUF CALL CONMSG ; IF NOT 'EOT', NEXT MSG JR NC NXM-S ; AKMG: CALL ACKMSG SEND ACK MSG POP HL ;HL=RBUF ; RBUF(FULL)=0 RES FULL, (HL) SET EMTY, (HL) ;RBUF(EMTY)=1 ; SRBF: CALL STRBUF START NEXT RBUF RES ACK (HL) ; RBUF(ACK) = 0BSY, (HL) SET ;RBUF(BSY)=1 ; RET ; SVMG: INC С ;SET SAVE FLAG (SAVMSG) BC ;SAVE LENGTH & FLAG LD (SAVMSG+2), DE ; SAVE MSG ADDR LD LD (SAVMSG+4), HL ; SAVE BUF ADDR POP HL ;HL=RBUF RES FULL, (HL) ;RBUF(FULL)=0 RES EMTY, (HL) ; RBUF(EMTY) = 0RET ; ; **;;;** FILE STRLNK ROUTINE TO STRIP LINK MSG AND SET RBUF FLAGS ; ENTRY- HL=RBUF ADDR ; EXIT- STATUS FLAGS IN RBUF SET AS INDICATED ; IN LINK MSG PUSH HL STRLNK: ;MOVE RBUF TO IX POP IΧ ; (IX+0),00H LD CLEAR ALL FLAGS **;**SET UP BUFFER POINTER LD HL, RBUF+04H LD A, (HL) ;GET CHAR CP 01H ;=SOH JP NZ .LNKERR ;NO, ERROR INC **;YES**, CHECK FC HL A/(HL) LD ;NEXT CHAR CP 48H J=ACK\NAK MSG

	10	7.04886-4	;YES, CHECK SC
			NO, =SRVMSG
		42H NZ→LNKERR-\$	
<u></u>			YES, RBUF(SVM)=1
CHKSC:		нс	CHECK SC
	LD		START BCC
			INEXT CHAR
		41H	;=41H
			YES, CHECK LAST SC
	СР		;NO, =42H
		NZ, LNKERR-\$	
CHKLSC:	LD		; B = S C
	LD		;GET LAST SC
	CP		;=LAST SC
	JR	Z / CHKAC – S	;YES, CHECK AC
	LD	A / B	GET NEW SC
	LD	(SAVSC)/A	;SAVE NEW SC
	SET	NSC (IX+0)	;RBUF(NSC)=1
CHKAC:	INC	НL	;CHECK AC
	XOR	C	;CALC BCC
	LD	C . A	;SAVE BCC
	LD	A,(HL)	;NEXT CHAR
	СP	40H	;=40H
	JR	NZ / LNKERR-\$	NO, ERROR
	INC	HL	;YES, CHECK OC
	XOR	С	CALC BCC
			;SAVE BCC
	LD	A,(HL)	INEXT CHAR
	AND	70H	CLEAR LOWER BYTE
	СР	40H	;=40H
	JR	NZ, LNKERR-\$	
	LD	A, (HL)	;YES, GET OC
	BIT		;NAK BIT SET
		NZ . SVOC-\$	
			;NO, RBUF(ACK)=1
svoc:			SAVE PRESENT OC
		HL	
	XOR		CALC BCC
			SAVE BCC
		A,(HL)	
	BIT		;BIT 6 SET
		Z.LNKERR-\$	
	AND		CLEAR MSB
		(SAVIC),A	SAVE IC
	INC		CHECK STX
	XOR		CALC BCC
	LD		SAVE BCC
		A.(HL)	
	CP		;=STX
		NZ .LNKERR-\$	
			CLEAR B
СНКЕТХ:	DEC		;B=B-1
			B>256, NO ETX FOUND
	INC		FIND ETX
	146		

;CALC BCC XOR C LD C / A SAVE BCC ;NEXT CHAR LD A, (HL) CР ;=ETX 03H JR NZ, CHKETX-S ;NO, CHECK NEXT CHAR INC HL ;YES, CHECK BCC ;CALC BCC XOR С LD SAVE BCC C / A LD A, (HL) ;NEXT CHAR CP С ; C = 8 C C NZ .LNKERR-\$;NO, ERROR JR **;YES, CHECK FOR EOT** INC HL LD A, (HL) **;**NEXT CHAR CP 04 H ;=EOT JR NZ .LNKRET-\$;YES, RETURN ;NO, RBUF(LOG)=1 SET $LOG_{\prime}(IX+0)$ (SAVMSG+2), HL ; SAVE LOCATION POINTER LD LD HL,0000H ; LD (SAVMSG),HL ;COUNT=00,SAVE FLAG=0 LNKRET-\$ **;**RETURN JR LNKERR: SET ERR (IX+O) ;SET ERR BIT PUSH IX ; MOVE RBUF TO HL LNKRET: POP HL RET **FRETURN** ;;; FILE SRVMSG ; ; ROUTINE TO ANALYZE SERVICE MSG ENTRY- SAVOC=PRESENT VALUE OF OC ; EXIT- DECODE OC SERVICE MSG, ON ERROR, SET CARRY & RETURN, ELSE, TAKE INDICATED ACTION ; ; ;GET PRESENT OC A, (SAVOC) SRVMSG: LD **;CLEAR UPPER 5 BITS** AND 07H JR Z,NOINST-S ;=O, NO INSTRUCTION SUB 03H ;=3, ACCEPT ALL CALLS JR Z, ACALL-\$ DEC Δ Z,RFD-\$;=4, READY FOR DISCONNECT JR DEC ; Α ;=5, ACCEPT NO CALLS JR Z.NCALL-S DEC 2 Α JR Z.DIS-S ;=6, DISCONNECT 7 SCF ;ERROR, SET CARRY RET ; RETURN ; NOINST: JP ACKMSG SEND ACK MSG ; ACALL: HL,HCB GET RNP STATUS ADDR LD SET ACPT, (HL) ;HCB(ACPT)=1 1 JP ACKMSG SEND ACK MSG ; RFD: LD A,44H ;A=OC(ACK,RFD) **;**PUT IN SRVMSG BUFFER LD (SBOC) A

;GET BCC LD A, (SBCC) XOR D4H ;CALC NEW BCC (SBCC)/A **;**PUT IN SBUF LD SEND SERVICE MSG JP STRSBF ; GET RNP STATUS ADDR-LD HL,HCB NCALL: RES ACPT, (HL) ;HCB(ACPT)=0 SEND ACK MSG JP ACKMSG ; GET RNP STATUS ADDR DIS: LD HL,HCB RES ONLN.(HL) ;HCB(ONLN)=0 RET ; ;;; FILE GETMSG ROUTINE TO GET NEXT MSG OFF RBUF ; ********** ENTRY- SAVMSG+2=MSG ADDR, SAVMSG+4=SAVED BUF EXIT- B=BYTE COUNT OF MSG, DE=MSG ADDR, ; HL=BUF ADDR, CARRY IF BUFFER SAVED, NZ IF ERROR ; GETMSG: LD A, (SAVMSG) ;GET STATUS CLEAR FLAGS 0 R Α NZ,SAVED-\$;IF A>O, DATA SAVED JR **;NO, GET BYTE COUNT** LD ArO3H HL, (SAVMSG+2) ;HL=START OF MSG LD BC, OOFFH ;NO, SET BC=255 LD ;'ETX' FOUND CPIR ;NO, ERROR-RET WITH NZ RET NZ ; ; CALC BYTE COUNT HL,OOFFH ;HL=255 LD ;HL=BYTE CNT SBC HL,BC LD B/L ;B=BYTE CNT DE, (SAVMSG+2) ; DE=MSG ADDR LD SET ZERO XOR A SCF **;**SET CARRY RET ; ; SAVED: SCF SET CARRY DEC ;STATUS>1 Α Z,NOBUF-\$;NO, BUF NOT SAVED JR **;**CLEAR CARRY & SET Z XOR Α HL, (SAVMSG+4) ; YES, GET BUF ADDR LD NOBUF: LD DE, (SAVMSG+2) ;GET MSG ADDR RET ; ; ; ; ROUTINE TO GET 1-4 LINKED BUFFERS OFF AQUE ENTRY- B=BYTE COUNT ; ; EXIT- HL=BUF ADDR OF FIRST BUF IN LINKED SET LINK BIT SET IN ALL BUT LAST BUFFER AND ; POINTER TO NEXT BUF IN LAST 2 BYTES ;

; PUSH BC ;SAVE COUNT ;SAVE MSG ADDR GETBUF: PUSH DE LD A.B A=BYTE COUNT LD 8,00H ;B=0 IX / AQUE JIX=AQUE LD INC B ;B=B+1 NUMBUF: GET BUF OFF AQ CALL GET JR C.RSAQ-\$; IF EMPTY, RESTORE OTHER BUFS PUSH HL SAVE BUF ADDR SUB 3AH ;A=A-58, A<=0 NC, NUMBUF-\$ INO, NEXT BUF JR ; POP HL GET BUF ADDR LD (HL),00H CLEAR STATUS DEC B ;B=B-1, B=O Z.LSTBUF-\$ FYES, ONLY BUF JR ; POP IX **;NO**, LINK BUFFERS LNKBUF: (IX+62)/L LOWER BYTE NEXT BUF ADDR LD JUPPER BYTE NEXT BUF ADDR LD (IX+63)/H PUSH IX ;MOVE IX TO HL POP HL ; LINKED TO NEXT BUF SET LNK, (HL) DJNZ LNKBUF-\$ LAST BUF ; LSTBUF: XOR **CLEAR CARRY** Α RESTORE MSG ADDR POP DE POP BC **FRESTORE COUNT** RET : RSAQ: LD IX, AQUE LOAD QUEUE ADDR SCF SET CARRY RSQ: DEC B ;B=B-1, B=0 Z.LSTBUF+1-\$;YES, RETURN JR POP HL GET BUF ADDR ;PUT BUF BACK CALL PUT RSQ-\$ INEXT BUF JR ;;; FILE CONMSG ROUTINE TO CONVERT LOG MSG'S FROM RBUF TO OBUF'S ; ENTER- HL=OBUF, DE=RBUF MSG ADDR, ; (SAVMSG+1)=BYTE COUNT ; EXIT- CARRY SET IF 'EOT' ; PUSH HL CONMSG: SAVE OBUF POINTER EX DE.HL ; DE=OBUF, HL=MSG A,(HL) ;GET CHAR LD CP 01H ;="SOH" JP NZ.LOGERR2 **INO** ERROR INC HL ; LD A, (HL) ;GET FC

			C 7
60 (0.0	1-50	
		;=FC	
		INO. ERROR	
		START BCC	
		;GET SC	
		INEXT BYTE	
XOR C		;CALC BCC	
LD C	- A	;	
INC H	IL .	;GET AC	
LD A	(HL)	;NEXT BYTE	
PUSH A	F	;SAVE AC	
XOR C		CALC BCC	
	A	;	
		JGET OC1	
		INEXT BYTE	
		;=001	
	Z.LOGERR1-\$		
		CALC BCC	
		GET OC2	
		INEXT BYTE	
PUSH A		SAVE OC2	
		;CALC BCC	
LD C	A A	•	
INC H	IL	;GET IC	
LD A	(HL)	;NEXT BYTE	
CP 4	8H	; = I C	
JR N	Z.LOGERR-\$;NO, ERROR	•
XOR C		;CALC BCC	
LD C	• A	;	
		GET STX	
		INEXT BYTE	
		;=STX	
		NO, ERROR	
XOR C		CALC BCC	
•	,		
TNC U	۱L	;HL=POINTER TO TEXT	
		SAVE RBUF POINTER	
	SAVESCE A		
	SAVELLIPA	SAVE BUU	
;			
		;HL=0C2	
		; A=AC (LU#)	
PUSH H		SAVE OC2	
		GET DCB ADDR (A=DEV#, HL=	D C B)
		; I X = D C B	
		;GET OC2	
		;PUT OC2 IN DCB	
EX D)E,HL	;HL=OBUF, DE=DCB	
;			
LD A	<pre>(SAVMSG+1)</pre>	A=BYTE COUNT	
INC H	IL .	;	
LD ((HL) / A	;PUT IN OBUF	
	۱L	;	
		PUT DCB IN OBUF	
		;	

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;	INC EX	DE,HL	; ;HL=OBUF DATA POINTER ;DE=OBUF DATA POINTER ;HL=RBUF MSG ADDR
PUTMSG:	LD SUB JP	M,LTBF	; A=BYTE COUNT ; ; A=A-58 ; IF A>O, MOVE 58 BYTES ; IF A=O NO MORE DATA
;	LD	BC • 003AH	SAVE REMAINING BYTE COUNT SET BC=58 BYTES MOVE BLOCK OF DATA
	LD INC LD INC INC INC INC	E . (HL) HL D . (HL) HL . (SAVMSG+2) DE DE DE DE	<pre>;HL=OBUF POINTER ;DE=LINKED BUF ADDR ; ; ;HL=RBUF POINTER ; ; ;DE=OBUF DATA POINTER ;PUT DATA IN OBUF</pre>
NOBF:	CALL LD LD CP	MOVMSG A.(SAVBCC) C.(HL)	; A = C
;	INC	(SAVMSG)♪A HL	<pre>;A=0 ;CLEAR STATUS ; ;GET NEXT BYTE ;='EOT' ;SET CARRY ;YES, RETURN WITH CARRY SET ;CLEAR CARRY ;</pre>
LOGERR: LOGERR1: LOGERR2:	POP	HL HL ERR,(HL) A	; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ;
; MOVMSG: MVMG:	L D X O R L D I	A,(SAVBCC) (HL)	;GET BCC ;CALC BCC ;MOVE BYTE

PE,MVMG ;IF BC NOT O, NEXT BYTE (SAVBCC),A ;SAVE BCC JP LD (SAVMSG+2), HL ; SAVE RBUF MSG POINTER LD ; RET ; ;;; FILE GENXBF ROUTINE TO GENERATE NEXT XBUF HL,XBUF ;HL=XBUF GENXBF: LD BSY, (HL) ;XBUF(BSY)=1 BIT **;YES**, RETURN RET NZ ;NO, XBUF(FULL)=1 BIT FULL (HL) NZ,STXB-\$;YES, START XMIT OF XBUF JR ;NO, XBUF(EMTY)=1 BIT EMTY, (HL) ;NO, NEXT LOG MSG Z,GNMG-\$ JR ; CALL GENLNK GENERATE LINK MSG RES EMTY (HL) ;XBUF(EMTY)=0 ; GNMG: PUSH HL ;(SP)=XBUF LD IX, XQUE ; IX = XQUE GNI: GET NEXT IBUF OFF XQ CALL GET JR C.RTXB-\$ **;IF XQ EMPTY, RETURN** CONVERT IBUF TO LOG MSG CALL CONLOG JR NC GNI-\$; IF XBUF NOT FULL, GET NEXT IBUF ; CALL PUT ;PUT BACK ON XQ POP HL ;HL=XBUF SET FULL (HL) ;XBUF(FULL)=1 ; STXB: A/(RBUF) ;A=RBUF STATUS LD BIT BSY A ;RBUF(BSY)=1 ;YES, RETURN RET NZ CALL STRXBF ;NO, START XMIT XBUF ;XBUF(BSY)=1 SET BSY (HL) RET . : RTXB: POP ;RESTORE HL HL RET ; ; ; ROUTINE TO GENERATE LINK MSG IN XBUF ; ENTRY - HL = XBUF; EXIT- LINK SET UP IN XBUF, (SAVLOC)=END OF LINK GENLNK: PUSH HL SAVE BUF ADDR HL,XBUF LD CHECK XBUF STATUS BIT ACTV, (HL) ;xbuf(Actv)=1 JR NZ,GXSC-\$ **;YES, GENERATE NSC** ; LD HL/SBUF **;NO, CHECK SBUF STATUS**

E11

BIT ACTV (HL) ;sbuf(actv)=1 INO, SEND ACK MSG JR Z,NACTV-S LD A. (SBSC) ;YES, GET OLD SC RES ACTV (HL) ;SET SBUF(ACTV)=0 GEN NSC JR GNSC-\$; NACTV: LD A,42H START NSC ; JR G N S C - \$; GXSC: LD A, (XBSC) ;GET OLD SC ;set xbuf(actv)=0 RES ACTV, (HL) ; GNSC: XOR 03H GENERATE NEW SC POP IX ;IX=BUF ADDR LD (IX+OAH)/A **;**PUT NSC IN BUF LD C (I X + 09 H) JGET FC XOR C START NEW BCC LD $C_{\prime}(IX+OCH)$;GET OC ;CALC BCC XOR С (IX+10H) • A **;**PUT NEW BCC IN XBUF LD PUSH IX ;MOVE BUF ADDR TO HL POP HL SET ACTV, (HL) ;set buf(ACTV)=1 LD BC,0011H ; ADD HL, BC ;HL=END OF LINK LD A 🖉 0 4 H ;A='EOT' LD (HL),A ;PUT IN XBUF **;**SAVE END POINTER (SAVLOC) HL LD RET ; FILE CONLOG ; ROUTINE TO CONVERT IBUFS TO LOG MSGS AND ; PLACE THEM IN XBUF ; ; ENTRY- HL=IBUF ; EXIT- RETURN WITH CARRY IF ERROR, ; ; (SAVLOC) = END OF MSG SAVE IBUF CONLOG: PUSH HL SET CARRY SCF ;IBUF(BSY)=1 BIT - BSY, (HL) RET ΝZ ;YES, RETURN BIT ;IBUF(ERR)=1 ERR (HL) RET ΝZ **;YES**, RETURN SET BSY, (HL) ;SET IBUF(BSY)=1 ; INC GET BYTE COUNT HL LD C/(HL) ;PUT IN C LD B/00H ;B=0 LD HL/(SAVLOC) GET XBUF POINTER PUSH HL SAVE XBUF POINTER ADD HL, BC ADD BYTE COUNT DE, XBFN-OCH ;GET XBUF END - HEADER COUNT LD CALL HLDE ;RETURN CARRY IF HL>DE

•	JR	C,FULRET-\$	XBUF FULL YES, RETURN
;	LD LD LDIR	DE BC,DDDBH HL,LOGHDR	;GET XBUF ;SAVE FOR LATER ;BC=BYTE COUNT ;HL=LOG MSG HEADER BUF ;MOVE HEADER TO XBUF ;SAVE NEW XBUF LOC
;	POP POP CALL PUSH PUSH POP	HL DSTDCB IY HL	;IX=START OF HEADER ;HL=IBUF ;HL=DST DCB, IY=IBUF ;MOVE DCB TO IY ; ;IY=DCB
;	L D O R	A, (IY+3) 40H (IX+5), A C	;START BCC ;GET AC FROM DCB ; ;PUT IN XBUF ;CALC BCC ;
;	0 R	40H (IX+7),A C	;GET OC2 FROM DCB ; ;put in XBUF ;calc bcc ;
	OR LD XOR	40H (IX+8).A C C.A	;GET IC FROM DCB ; ;PUT IN XBUF ;CALC BCC ;SAVE BCC ;HL=IBUF
; mvbf: ;	LD BIT JR	A,(HL) LNK,A Z,LSBF-\$;GET IBUF STATUS ;IBUF(LNK)=1 ;NO, LAST BUFFER
	INC LD CALL PUSH LD INC LD EX POP JR	MOVBLK DE E / (HL) HL D / (HL) DE / HL	; ;B=BYTE COUNT ;MOVE BLOCK OF DATA ;SAVE XBUF LOC ;GET NEXT BUF ADDR ; ;HL=NEXT BUF ADDR ;DE=XBUF LOC ;MOVE NEXT BLOCK
LSBF:	INC LD LD	B,(HL)	;HL≃BYTE COUNT POINTER ;B=BYTE COUNT ;PUT BCC IN XBUF

E13

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A = 04 H LD ;A='EOT' INC DE ; (DE),A ;PUT IN XBUF (SAVLOC),DE ;SAVE XBUF LOC LD LD ; XOR CLEAR CARRY A RET ; FULRET: SCF SET CARRY POP HL **;**ADJUST STACK POP HL **FRESTORE HL** RET ; ; ; MOVBLK: INC HL ; INC HL INC HL ;HL=DATA LOC ; MVBT: LD A, (HL) GET DATA LD (DE)/A **;**PUT IN XBUF XOR C ;CALC BCC LD C . A SAVE BCC ΗL ;INC POINTERS INC INC DE ; ;CONTINUE TILL B=0 DJNZ MVBT-\$ RET ; ; ; SYN, SYN LOGHDR: DEFW 1616H DEFB 01H ; S О Н DEFB 60H ; F C ; S C DEFB 41H DEFB 40H ;AC (40H+DEV#) DEFB 41H ;001 ;0C2 (40H+ACK+REQ) DEFB 40H DEFB 48H ;IC (48H=TTY) DEFB 02H ; S T X DEFB 03H ; E T X ; ;;; FILE DEVOUT ROUTINE TO SET UP OUTPUT TO A DEVICE ; ; ; ENTER- HL=DST DCB, IY=OBUF ADDR ; EXIT- IF CARRY, DEVICE IS BUSY OR RD IS SET ; IF BAD COUNT OR DEV NOT ON LINE, OBUF ; PUT ON AQ ; ELSE, PARAMETERS STORED IN DCB (IX=DCB) ; LIST X DEVOUT: SCF SET CARRY BIT BSY, (HL) ; DCB(BSY) = 1RET NZ **FYES**, RETURN WITH CARRY ;NO, DCB(RD)=1 BIT RD/(HL) RET NZ YES, RETURN WITH CARRY

	J R S E T R E S R E S L D	BSY,(HL) ERR,(HL) ETB,(HL) A,(IY+3)	<pre>;NO, DEV NOT ON LINE ;SET DCB(BSY)=1 ;SET DCB(ERR)=0 ;SET DCB(ETB)=0</pre>
;	EX POP LD LD BIT JR SET LD	(SP),HL IX (IX+8),H (IX+7),L ETB,(HL) Z,6 ETB,(IX+0) E,(HL) C,(IX+2)	; ;HL=OBUF ;IX=DCB ;PUT OBUF ADDR IN DCB ; ;OBUF(ETB)=1 ;NO, SKIP NEXT INST ;YES, SET DCB(ETB)=1 ;E=PRESENT BUF STATUS ;C=DEVICE BUS ADDR ;
;	JR LD LD SUB JR	Z,NLNK-\$ B,3AH A,(HL) B C,IOERR-\$ D,A	;OBUF(LNK)=1 ;NO, GET BYTE COUNT ;YES, SET B=58 BYTES ;A=TOTAL BYTE COUNT ;A=REMAINING BYTE COUNT ;JF A <d, byte="" count<br="" error="" in="">;D=REMAINING BYTES ;SKIP NEXT INST</d,>
; NLNK: ;	INC INC	HL HL	;B=TOTAL BYTE COUNT ; ; ;HL=OBUF DATA POINTER
	CALL	SAVPR4	;SAVE ALL PARAMETERS IN DCB
;	CALL XOR OUT RET	SETMSK A (C) • A	SET I/O MASK BIT A=00
IOERR:	RES PUSH POP		;SET DCB(BSY)=0 ; ;HL=0BUF
PUTAQ:	LD CALL XOR BIT RET LD LD PUSH POP JR	IX > AQUE PUT A LNK > (HL) Z L > (IY + 62) H > (IY + 63)	; ; ; ; ; ; ; ; ; ; ; ; ; ;

;;; FILE DEVIN : ROUTINE TO SET UP INPUT FROM A DEVICE ; ENTER- HL=SRC DCB EXIT- IF CARRY, DEVICE IS BUSY OR NOT ON LINE ; OR RD NOT SET OR NO BUFFERS LEFT IN AQ ELSE, PARAMETERS STORED IN DCB (IX=DCB) SCE SET CARRY DEVIN: BIT BSY, (HL) ; DCB(BSY) = 1;YES, RETURN WITH CARRY RET ΝZ ;DCB(ONLN)=1 BIT ONLN, (HL) **;NO**, RETURN WITH CARRY RET Z ; DCB(RD) = 1BIT RD/(HL) **;NO**, RETURN WITH CARRY RET 7 SET BSY (HL) ;SET DCB(BSY)=1 ;SET DCB(ERR)=0 RES ERR (HL) ETB, (HL) RES ;SET DCB(ETB)=0 ; ;(SP)=DCB PUSH HL ;IX=AQUE LD IX,AQUE GET IBUF OFF AQ CALL GET POP ; IX = DCB IΧ ; IF EMPTY, RETURN WITH CARRY RET C PUSH HL ; ;IY=IBUF POP IY ; XOR ;A=00 Α ;CLEAR E LD E / A CLEAR IBUF STATUS LD (HL)/A ;PUT IBUF ADDR IN DCB LD (IX+7),L LD (IX+8),H INC HL **;**CLEAR BYTE COUNT LD (HL)/A ;D=58 BYTES LD D.JAH LD $C_{(1X+2)}$ C=DEVICE BUS ADDR INC HL ;GET SRC AC $A_{\prime}(IX+3)$ LD ;PUT IN IBUF LD (HL),A INC HL ; ;GET DST AC $A_{\prime}(IX+4)$ LD **;**PUT IN IBUF LD (HL)/A INC HL ;HL=OBUF DATA POINTER ; 0 R ;DST AC=0 Α Z,SETHST-\$;YES, SKIP ID HEADER JR. . ; BIT IDM, (IX+0) ;DCB(IDM)=1 JR Z,SETCNT-\$ **;NO**, SKIP ID HEADER ; LD A, (IX+3) GET SRC AC

LD (HDRID),A ;SAVE CALL IDHDR **;**PUT ID HEADER ON IBUF JR FININ-\$; ; SET HOST, (IY+O) ;SET IBUF(HOST)=1 SETHST: **;INITAL BYTE COUNT** SETCNT: LD B,D ; CALL SAVPR4 SAVE PARAMETERS IN DCB FININ: CALL SETMSK SET I/O MASK BIT **CLEAR DEVICE STATUS** ΙN A, (C) ΕI **;**ENABLE INTR RET ; ;; ID HEADER ADDED TO BUFFER- [XX] , XX=LU# ; START OF ID HEADER, "[" IDHDR: LD A,5BH LD (HL)/A **;**PUT IN IBUF INC HL ; JGET AC VALUE LD A, (HDRID) PUSH AF SAVE SRL A GET UPPER 4 BITS SRL A. ; ; SRL Α ; SRL Α CALL HEXASC CONVERT TO ASCII LD (HL)/A **;**PUT IN IBUF INC HL ; POP AF GET LOWER 4 BITS CONVERT TO ASCII CALL HEXASC ;PUT IN IBUF LD (HL)/A INC HL A,5DH ;END OF ID HEADER, ']' LD ;PUT IN IBUF LD (HL)/A INC HL ; A, ' ' ; SPACE LD LD (HL)/A PUT IN IBUF INC HL ; ;B=53 B,35H LD RET ; ;;; FILE INTR INTERRUPT HANDLER ; ; ENTER- INTERRUPT FROM LEVEL 1 EXIT- INTERRUPT DEVICE SERVICED ; ; INTR: PUSH AF SAVE REG'S IN A,(13H) GET HOST STATUS BIT 01H/A ; READ=1 JP YES, GO READ CHAR NZ•HOSTR PUSH BC LD A, OAH ; ; OUT (1EH)/A IN A, (1EH) **;INTR REQUEST REG**

LD ;SAVE IN B B / A LD A. (MASK) ; I/O MASK AND **;CLEAR MASKED REQUESTS** R JR 7.FIN-\$ FRETURN IF NO REQ'S SRL ;BIT 1 SET Α JP C.SYNWR **;YES**, SYNC WRITE ; PUSH DE IND, SAVE REMAINING REG'S PUSH HL ; ; PUSH IX PUSH IY : LD HL, INTTAB+2 ;HL=INTR TABLE ADDR+2 LSB: SRL ;LSB=1 Α ;YES, GET DCB ADDR JR C.DCBA-S INC HL ;NO, HL=HL+2 INC HL ; **;**NEXT BIT JR LSB-\$; DCBA: LD E, (HL) ;E=LOWER BYTE DCB ADDR INC HL : **;**D=UPPER BYTE DCB ADDR LD D/(HL) ;; ; ROUTINE TO HANDLE IO INTERRUPTS ; PUSH DE ; POP IX ;IX=DCB **;**DEVICE STATUS BYTE LD A, (DE) BIT BSYA **;**DEVICE BSY JR Z, INTERR-\$ **;NO, GOTO ERROR ROUTINE** ;GET BYTE COUNT LD B,(IX+9) ;GET DEV BUS ADDR LD C (I X + 2) H,(IX+13) ;GET DATA POINTER LD $L_{(IX+12)}$: LD BIT ; DCB(RD) = 1RD A Z,WRITE-S ;NO, GO WRITE BYTE J R ; ; READ DATA BYTE FREAD BYTE CALL DEVRD JR NC,SVPR-\$ **;SAVE PARAMETERS, IF NC** BIT $ETB_{\prime}(IX+0)$;DCB(ETB)=1 JR Z,FINIO-\$ **;NO**, END OF READ BIT ;YES, DCB(ERR)=1 ERR (IX+0)JR Z,FINIO-\$;NO, END OF READ PUSH IX **;YES**, INPUT NOT FINISHED POP HL ;HL=DCB ADDR RES BSY (IX+0);SET DCB(BSY)=0 RES IDM (IX+O) ;SET DCB(IDM)=0 CALL DEVIN START NEW READ CYCLE JR SVPR+3-\$ JEND OF THIS READ ; WRITE: CALL DEVWR WRITE BYTE JR NC SVPR-\$ SAVE PARAMETERS, IF NC ; RES BSY, (IX+0) ;SET DCB(BSY)=0 FINIO:

JR 5 SKIP NEXT INST ; SVPR: CALL SAVPR2 SAVE PARAMETERS FRESTORE REG'S POP ΙY POP IΧ ; JP RRET **FRETURN FROM INTR** ; ;SET DCB(ERR)=1 INTERR: SET ERR (IX+O) JR FINIO-\$; ;; SAVPR4: SAVE 4 PARAMETERS IN DCB LD (IX+11),D LD (IX+10),E SAVPR2: **;**SAVE 2 PARAMETERS IN DCB LD (IX+9),B LD (IX+13),H ; LD (IX+12)/L ; RET ; ;; ; SETMSK: LD $B_{\ell}(IX+1)$ JGET DEV MASK BIT RD (IX+O) ; DEV(RD) = 1JR NZ/4 ;YES, SKIP NEXT INST . SLA В **;**ADJUST FOR WRITE LD A (MASK) ;GET I/O MASK 0 R ;SET BIT В LD (MASK)/A **FRESTORE MASK** RET ; ;; ; JGET DEV MASK CLRMSK: LD $A_{\ell}(IX+1)$ BIT RD (IX+O) ; DCB(RD) = 1JR NZ·4 YES, SKIP NEXT INST SLA A **;**ADJUST FOR WRITE CPL **;INVERT** LD B / A ;MOVE TO B A. (MASK) GET PRESENT I/O MASK LD AND R CLEAR BIT (MASK) A FRESTORE MASK LD RET ; **;;;** FILE DEVWR ; ROUTINE TO WRITE A BYTE TO A DEVICE ; ; ENTER- HL=OBUF DATA POINTER, IX=DST DCB ; **B=PRESENT REMAINING BYTES** ; C=DEVICE BUS ADDR ; EXIT- CHAR MOVED TO DEVICE, PARAMETERS UPDATED DEVWR: LD A/(HL) GET CHAR FROM OBUF OUT (C),A WRITE TO DEVICE INC HL JUPDATE POINTER DJNZ SVWR-S CONTINUE TILL B=0

;			
;	LD LD CALL CALL POP	IX H (IX+8) L (IX+7) IX AQUE PUT C FULERR IX	
;	JR CALL BIT JR	NZ,NXLK-\$ CLRMSK ETB,(IX+0) NZ,6 RD,(IX+0)	;OLD OBUF(LNK)=1 ;YES, NEXT BUF ;CLEAR I/O MASK BIT ;DCB(ETB)=1 ;YES, SKIP NEXT INST ;SET DCB(RD)=1 ;SET CARRY ;END OF OUTPUT
, NXLK:	L D E X L D L D	D,(HL) DE,HL (IX+8),H (IX+7),L	;GET NEXT OBUF ADDR ; ; ;HL=NEW OBUF ;SAVE NEW OBUF ADDR IN DCB ;
	BIT JR LD LD SUB JR	LNK,E Z,NOLNK-\$ B,3AH A,D B C,BDLK-\$;GET TOTAL REMAINING BYTES ;E=NEW BUF STATUS ;OBUF(LNK)=1 ;NO, LAST BLOCK ;B=58 BYTES ;A=TOTAL REMAINING BYTES ;A=TOTAL-58 ;IF A<0, BAD LNK BIT ;SAVE IN DCB ;SKIP NEXT INST
; BDLK: ;	RES	LNK.E	SET OBUF(LNK)=0
NOLNK:	INC INC INC	(IX+10)≠E HL HL	<pre>>B=TOTAL REMAINING BYTES >SAVE PRESENT OBUF STATUS ; ; ; ; ; ; ; ; ; ; ; ; ;</pre>
; SVWR:	X O R R E T	A	CLEAR CARRY
; fulerr:		(SP),HL IX,EQUE	;SAVE QUEUE ADDR ;HL=QUEUE ADDR, (SP)=BUF ADDR ;IX=EQUE ;PUT ON EQ

FRETURN IF FULL RET C POP HL HL=BUF ADDR CALL PUT ;PUT ON EQ CALL C.GET ; IF FULL, REMOVE QUEUE ADDR RET ; ; ;;; FILE DEVRD ROUTINE TO READ A BYTE FROM A DEVICE ; *********************** ENTER- HL=IBUF DATA POINTER, IX=SRC DCB ; **B=PRESENT REMAINING BYTES** ; ; C=DEVICE BUS ADDR ; EXIT- CHAR MOVED TO IBUF, PARAMETERS UPDATED ; FREAD CHAR A,(C) DEVRD: IN CP 03H ;A='CNTL C' JP Z,CMDIN ;YES, SET CMD INPUT BUF LD (HL)/A ;PUT IN IBUF INC JUPDATE POINTER HL ;MOVE CHAR TO D LD D / A GET BUFFER TERMINATOR LD A, (ENDCHR) CP 0 **;**D=END CHARACTER JR Z,FINRD-\$ YES, END OF LINE DJNZ SVRD-5 ;NO, CONTINUE TILL B=0 ; LD A, OAEH ;A=3*58, 3 BLOCKS (IX+11) ;TOTAL COUNT <= A CР ;YES, NEXT BLOCK JR NC NX BK - \$ SET ERR (IX+O) ;NO, SET DCB(ERR)=1 EMTAQ+3-\$ **;**CLOSE BUFFER JR ; NXBK: PUSH IX SAVE SRC DCB ADDR PUSH HL SAVE POINTER LD IX,AQUE ; CALL GET JGET NEW BUF OFF AQUE **;IF CARRY, AQ EMPTY** JR C, EMTAQ-\$ POP DE ; ΕX DE,HL ;DE=NEW BUF > HL=LINK LOC PUT LINK ADDR IN OLD BUF LD (HL),E INC HL ; LD (HL),D ; PUSH HL POP ΙY ;IY=OLD BUF POINTER SET LNK, (IY-63) ;SET OLD BUF(LNK)=1 ; POP IX FRESTORE SRC DCB ADDR LD A, JAH ;A=58 LD B/A ;8=58 ADD A. (IX+11) FINCREASE TOTAL COUNT BY 58 LD (IX+11),A STORE IN SRC DCB ΕX DE,HL HL=NEW BUF ADDR XOR A ;A=0

E21

	INC LD INC LD INC	(HL).A HL	CLEAR STATUS CLEAR COUNT CLEAR SRC AC CLEAR DST AC HL=NEW BUF DATA POINTER
SVRD:	X O R R E T	Α	CLEAR CARRY
		IX	;HL=BUF POINTER ;IX=DCB ;SET SRC DCB(ETB)=1 ;SKIP NEXT INST
	SUB LD LD DEC PUSH INC IN AND JR SET	A, (IX+11) B L, (IX+7) H, (IX+8) HL (HL), A HL IX C A, (C) O8H Z, CKHOST-\$ PAR, (HL) HOST, (HL)	<pre>;ADJUST B ;GET TOTAL BYTE COUNT ;SUBTRACT UNUSED BYTES ;GET MAIN IBUF ADDR ;HL=IBUF ;HL=BYTE CNT ADDR ;PUT FINAL BYTE COUNT IN IBUF ;HL=IBUF ADDR ;SAVE SRC DCB ADDR ;DEVICE STATUS ADDR ;GET STATUS BYTE ;BAD PARITY BIT ;NO, CHECK HOST BIT ;YES, SET IBUF(PAR)=1 ;IBUF(HOST)=1 ;NO, CHECK FOR COMMAND</pre>
;			;IX=XQUE, HL=IBUF
CKCMD:	LD SCF OR CALL	A, (IX+6) A Z, CMDPRC	;PUT ON QUEUE ;GET DST AC ;SET CARRY ;A=O ;YES, PROCESS CMD ;IF CARRY, SAVE BUF
IBPT:	CALL		;GET WQUE ADDR ;PUT IBUF ON QUEUE ;QUEUE FULL PUT ON EQ
; .	BIT JR SET JR	CLRMSK ETB#(IX+O)	<pre>;RESTORE SRC DCB ADDR ;CLEAR I/O MASK BIT ;DCB(ETB)=1 ;NO, SKIP NEXT 2 INST ;YES, SET IBUF(ETB)=1 ;SKIP NEXT INST ;SET DCB(RD)=0 ;SET CARRY</pre>

E 2 2

; RET ; ;;; FILE HOSTIO ROUTINE TO SET UP HOST READ ; ; ENTER- ;; ; EXIT- ;; ; GET RBUF ADDR STRBUF: LD HL, RBUF BIT EMTY (HL) ; RBUF(EMTY)=1 RET ;NO, RETURN Ζ ; PUSH HL SAVE RBUF LD HL/HCB ;GET HCB ADDR BIT BSY, (HL) ;HCB(BSY)=1 ;YES, RETURN RET ΝZ ;HCB(RD)=1BIT RD, (HL) RET Z ;NO, RETURN ; A, (13H) GET STATUS IN DSRA ;HOST(DSR)=1 BIT ; Z,NODSR-\$ JR ; SET BSY/(HL) ;SET HCB(BSY)=1 ; LD BC 400H-4 ;BC=BYTE COUNT LD DE, RBUF+4 ;DE=DATA POINTER (HCB+2),BC SAVE IN HCB LD (HCB+4), DE ; LD XOR Α **;**CLEAR CARRY RET ; ;;; ROUTINE TO READ A CHAR FROM HOST CHAN ******* ; ENTER- ALL INTR'S DISABLED ; EXIT- PUT CHAR FROM HOST IN RBUF ; ; HOSTR: PUSH BC SAVE REG'S PUSH DE ; PUSH HL ; ; BC, (HCB+2) LD GET COUNT LD DE / (HCB+4) GET POINTER LD HL/HCB ;HL=HCB ADDR ; BIT SYNC, (HL) ;MODEM IN SYNC JR NZ, RDAT-S **YES**, READ CHAR BIT SYNC,A :NO, MODEM SYNC BIT SET Z • R R E T - \$ NO, RET JR SET SYNC, (HL) ;YES, SET HCB(SYNC)=1

;

A,(12H) FREAD CHAR RDAT: IN J=SYNC CHAR CР SYN ;YES, SKIP CHAR JR Z, RRET-\$ (DE)/A **IND, SAVE IN RBUF** LD JUPDATE POINTER INC DE · JUPDATE COUNT DEC BC EOT CP ;=EOT CHAR JR ;YES, END OF RBUF Z, FINSR-\$; LD A . B ; ADD A.C ;BC=0 Z, FINSR-\$;YES, END OF BUFFER JR ; (HCB+2), BC ; SAVE COUNT SRET: LD (HCB+4), DE LD SAVE POINTER ; POP HL ;RESTORE REG'S RRET: POP DE ; ; FIN: POP BC ARET: POP AF ;RESTORE AF ΕI ; RETI ; ; FINSR: RES RD,(HL) ;SET HCB(RD)=0 ;SET HCB(BSY)=0 RES BSY, (HL) ;set hcb(sync)=0 RES SYNC, (HL) LD HL, RBUF ;SET RBUF(FULL)=1 SET FULL/(HL) LD HL,400H-4 CALCULATE BYTE COUNT SBC HL,BC (RBUF+1),BC ;PUT IN RBUF LD JR RETURN RRET-\$; ;;; ROUTINE TO SEND CHAR TO HOST ; ENTER- ALL INTR'S DISABLED ; EXIT - CHAR SENT TO HOST ; ; HOSTW: PUSH DE SAVE DE LD BC (HCB+2) JGET COUNT GET POINTER DE (HCB+4) LD HL·HCB LD ;HL=HCB LD A. (DE) ;GET CHAR OUT (12H)•A SEND TO HOST INC DE ;INC POINTER DEC BC **;**DEC COUNT EOT ;='EOT' CР JR Z, FINSW-S YES, END OF BUFFER LD A,B ; ; A = 0 ADD A.C JR Z,FINSW-5 ;YES, END OF BUFFER (HCB+2),BC LD SAVE COUNT (HCB+4),DE LD SAVE POINTER

```
RRET-$
        JR
                         RETURN
;
                         ;set HCB(RD)=1
FINSW:
        SET
            RD, (HL)
                         ;SET HCB(BSY)=0
        RES
            BSY, (HL)
            HL / XBUF
                         ;HL=XBUF
        LD
                         ;set xbuf(emty)=1
        SET
            EMTY, (HL)
                         RETURN
             RRET-$
        JR
;
;;; FILE CMDPRC
;
;
        ROUTINE TO SET UP COMMAND BUFFERS
******
        ENTER - IX = DCB
;
        EXIT- BUFFER SET UP TO RECEIVE COMMAND
;
;
                         GET BUF ADDR
CMDIN:
        LD
            L_{\prime}(IX+7)
            H,(IX+8)
                          ;
        LD
        PUSH IX
                          SAVE DCB ADDR
                          FRESTORE PRESENT BUF TO AQ
        CALL IDERR
        POP IX
                          ;
                         GET PRESENT AC
        LD
            A,(IX+6)
        PUSH AF
                          ;SAVE AC
;
                         ;PUT CMD PROCESSOR # IN DCB
        LD
             (IX+6),00H
        CALL DEVIN
                          SET UP INPUT TO CMD PROC
        POP AF
                         GET OLD AC
                          RESTORE TO DCB
             (IX+6),A
        LD
        SCF
                          ;
        RET
                          ;
;
        ROUTINE TO HANDLE COMMANDS
;;
;
;
        ENTER- IX=DCB, HL=BUF
        EXIT- COMMAND PROCESSED
2
CMDPRC:
        PUSH HL
                          ;
        POP
             IΥ
                          ; IY=BUF
        INC
                          ;
             HL
        INC
                          ;
             HL
        INC
             HL
        INC
             НL
                          ;HL=DATA POINTER
        LD
             A,(HL)
                         ;GET FIRST CHAR
             ' A '
        СP
                         ; = A
                         ;YES, ASIGN # TO DCB(SRC AC)
        JR
             Z,ASGN-$
             • D •
                         ·;= 'D'
        CP
                         ;YES, DISCONNECT DEVICE FROM HOS
        JR
             Z, DDIS-$
                         ;= 'CNTL A'
        CР
             01H
                         ;YES, SEND SELECT TO HOST & ASGN
        JR
             Z, DSEL-$
;
FINC:
                          ;
        XOR
             Α
        RET
                          ;
;
```

E 2 5

;; ASGN: PTAC: ;	JR OR JR CP JR PUSH CALL POP BIT JR POP	C FINC-S A Z CKHST-S O4H NC NONE-S AF DEVDCB HL ONLN (HL) Z NODEV-S	<pre>;CONVERT ASCII CHAR TO HEX BYTE ;CARRY, IF ERROR ;A=HOST ;YES, CHECK HOST ;A<04H ;NO, NO SUCH DEVICE ;SAVE AC ;GET DCB ADDR ;HL=DCB ;DCB(ONLN)=1 ;NO, DEV NOT ON LINE ;YES, GET AC ;PUT NEW AC IN DCB ;</pre>
СКНЅТ:	BIT JR PUSH POP INC	ONLN,(HL) NZ,PTAC-\$ IY	;GET HOST DCB ADDR ;HCB(ONLN)=1 ;YES, HOST IS ON LINE ; ;
;		B,3AH A,00H ERROUT~\$;B=REMAINING BYTES IN BUF ;A=ERROR MSG # ;SEND ERROR MSG
; NODEV:	INC LD CALL	IY HL HL (HDRID),A IDHDR A,O1H ERROUT-\$; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ; ;
NONE:	C ALL L D		; ; ; ; A=ERROR MSG # ; SEND ERROR MSG
DDIS:	JR OR RET CP CP CALL	Z)4H NC DISGEN IX • XQUE PUT	CONVERT ASCII CHAR TO HEX BYTE A=O YES, RETURN A <o4h NO, RETURN GENERATE DISCONNECT MSG PUT MSG ON QUEUE ;</o4h

.

•

RET ; ; CONVERT ASCII CHAR TO HEX BYTE DSEL: CALL AHBYTE JR C.FINC-\$; 0 R ;A=0 Α ;YES, RETURN RET Z ;A<04H CP 04H RET NC ;NO, RETURN GENERATE SELECT MSG CALL SELGEN LD IX,XQUE CALL PUT PUT MSG ON QUEUE XOR A ; RET ; ;; AHBYTE: LD B,0 ;B=0 ;GET FIRST NIBBLE JR GTNBL-\$; **;**ADD PREVIOUS VALUE*4 NXNBL: ADD A / B ADD A / B ; LD B/A SAVE IN B ; ;GET CHAR FOR NIBBLE GTNBL: INC HL LD A/(HL) ; CALL ASCHEX CONVERT TO HEX NIBBLE NC, NXNBL-S **;CARRY**, IF NON-HEX CHAR JR ; LD A, ODH ; CP (HL) ; = C R SCF ; RET ΝZ **;NO, RET WITH CARRY** ;YES, CLEAR CARRY XOR Α ;PUT VALUE IN A LD A B RET ; ; ;; 30H ;A<'0' ASCHEX: SUB ;YES, RET WITH CARRY RET C CP 17H ;A>'F' CCF RET ;YES, RET WITH CARRY C CР ; 0<=A<=9 0AH CCF ; RET NC ;YES, RET NO CARRY ;ADJUST CHAR SUB 07H ;NO CARRY IF, 'A'<=A<='F' СP OAH RET ; ; ;;; ROUTINE TO OUTPUT ERROR MESSAGES ; ; ENTER- HL=BUF DATA POINTER, IY=BUF ADDR ; A=ERROR CODE, B=REMAINING BYTES, IX=DST DCB ; EXIT- ERROR MSG PUT IN BUF AND DEVOUT ENTERED

; FRROUT.	рисн	BC	SAVE BYTE COUNT
LKKOUT.	PUSH	HI	SAVE DATA POINTER
			HL=ERROR TABLE ADDR
	ADD		A = A + A
		B,00	;
			BC=ERROR CODE OFFSET
			;HL=ERRTAB+OFFSET
		E,(HL)	;
	INC		;
	LD	D, (HL)	JDE=ERROR MSG ADDR
;			
		A, (DE)	A=MSG BYTE COUNT
	LD		BC=ERR MSG BYTE COUNT
			DE=ERR MSG POINTER
			HL=BUF DATA POINTER
		DE,HL	IMOVE EDDOD MEC TO DUE
;	LDIR		MOVE ERROR MSG TO BUF
,	POP	BC	GET REMAINING BYTES
		8	
		A, JAH	CALCULATE TOTAL BYTE COUN
			PUT IN BUF
			;SET DCB(BSY)= 0
	RES	RD,(IX+0)	;SET DCB(RD)=0
	PUSH	IX	;
	POP		;HL=DST DCB
	JP	DEVOUT	;OUTPUT MESSAGE
;		5000	
ERRTAB:			FRROR MSG DD
			;ERROR MSG 01 ;ERROR MSG 02
			ERROR MSG D2
			ERROR MSG 02
			ERROR MSG D2
			FERROR MSG 02
		ERRO2	FERROR MSG 02
;			
ERROO:	DEFB	13H	BYTE COUNT
		"HOST"	;
		NOT!	;
		• ON •	
		'LINE'	
	DEFM		
-	DEFW	OADDH	; L F , C R
	DEFB	050	BYTE COUNT
;		Urn	DITE COUNT
		• NOT •	
	DEFM	• NOT • • ON •	;
	D E F M D E F M	' ON '	;
	DEFM DEFM DEFM	' ON ' 'LINE'	; ; ;
	DEFM DEFM DEFM DEFM	' ON ' 'LINE'	; ; ; ; L F , C R
	DEFM DEFM DEFM DEFM	ON LINE	; ; ; L F , C R

```
DEFM ' NOT'
                           ;
        DEFM ' DEF'
                           ;
         DEFM 'INED'
                           ;
        DEFM '... '
                           ;
                           JLF/CR
         DEFW OAODH
;
;
   FILE INIT
        ROUTINE TO INIT SYSTEM
;
;
       SET UP RESTART JUMPS
;
;
            SP, 2FFFH
                        STACK=TOP OF MEMORY
       LD
;
            A,03CH
                         ;CODE FOR JUMP
       LD
       LD
            BC,START
                         JUMP ADDR FOR RST OOH
            (00),A
                         STORE JUMP
       LD
       LD
            (01H),BC
                         ;STORE ADDR
                         JUMP ADDR FOR RST 66H
       LD
            BC, MAIN
       LD
            (66H)/A
                         STORE JUMP
       LD
            (67H),BC
                         STORE ADDR
;
       CALL CIO
                         ;SET UP I/O, INTR & TIMERS
       RET
;
        ROUTINE TO SET UP INTERRUPTS
;
        SET UP OF CLOCK DEVICE (CTC)
        & INTR HANDLER (8259)
    ********
;
CIO:
             С,18н
                           ;CHANNEL #0
        LD
             DE,951FH
                           ;(TIMR, 25MS, INTR)
        LD
        OUT
             (C),D
                           ;
        OUT
             (C),E
                           ;
                           CHANNEL #1
        INC
             C
        LD
             DE,953EH
                           ;(TIMR, 5MS, INTR)
        OUT
             (C) \cdot D
                           ;
        OUT
             (C) • E
                           ;
        INC
                           CHANNEL #2
             C
             DE,9570H
        LD
                           ;(TIMR,1MS,INTR)
        OUT
             (C),D
                           :
         OUT
             (C)/E
         INC
                           CHANNEL #3
             С
         LD
             DE,OC5FAH
                           ;(CNTR,CNT=250,INTR)
         OUT
             (C),D
         OUT
             (C) \cdot E
                           ;INTR VECTOR ADDR
         LD
             BC.IVEC
         LD
             A / B
        LD
             I / A
                           SET UP UPPER BYTE
         LD
             A . C
                           SET UP LOWER BYTE
         OUT
              (18H) • A
;
        LD
             A,12H
                           SET UP INTR DEVICE (8259)
             (1EH) / A
         OUT
                           ;
```

E29

XOR JCLEAR A Α OUT (1FH) / A ; **;INIT MASK** LD (MASK)/A CPL ;INVERT A OUT (1 FH) / A ; ; ;SET INTR MODE 2 ΙM 2 ; ; SET UP INO CHANNELS ; LD BC,4011H ;reset usart, i\o chan(10,11) LD DE, OFA35H ;SET MODE & FUNCTIONS CALL MSET **;OUTPUT COMMANDS** ; I \ O CHAN(12,13) C,13H LD CALL MSET **JOUTPUT COMMANDS** C,15H ; I \ O CHAN(14,15) LD CALL MSET **;OUTPUT COMMANDS** C•17H LD ; I\O CHAN(16,17) CALL MSET ;OUTPUT COMMANDS CALL CIO **;**SET UP INTR & TIMERS RET ; ; MSET: OUT (C),B ;OUTPUT B OUT (C),D ;OUTPUT D OUT (C) • E ;OUTPUT E RET ; ; INTO: NOP ; INTERRUPT LEVEL O NOP ; NOP ; ; ΕI ; RETI ; INT1: EQU INTR ;INTERRUPT LEVEL 1 ; NOP ;INTERRUPT LEVEL 2 INT2: NOP ; ; NOP ΕI ; RETI ; ; INT3: NOP ;INTERRUPT LEVEL 3 NOP ; NOP ; ; ΕI ; RETI ; ;;; FILE SUBS ; ; FILE CONTAINING FLAG BIT ASSIGNMENTS AND ; SMALL SUBROUTINES USED BY MORE THAN ONE OTHER ROUTINE. ; ;

BSY:	EQU	0	QUEUE OR BUFFER BUSY
FULL:	EQU	1	FULL QUEUE OR BUFFER
RD:	EQU	1	;DEVICE READ BIT
EMTY:	EQU	2	SEMPTY QUEUE OR BUFFER
ACPT:	EQU	2	<pre>; RNP ACCEPTING CALL</pre>
ACK:	EQU	3	<pre>; PREVIOUS MSG ACKNOWLEDGED</pre>
HOST:	EQU	3	BUFFER TO OR FROM HOST
NSC:	EQU	4	;NEW SEQUENCE COUNT
ONLN:	EQU	4	;RNP OR DEVICE ON LINE
PAR:	EQU	4	; BAD PARITY IN BUFFER
	EQU	5	SERVICE MESSAGE REQUEST
ACTV:		5	BUFFER PRESENTLY ACTIVE
ETB:		5	END OF BUF, NOT END OF TEXT
LOG:		6	LOG MSG IN RBUF
LNK:		6	BUF LINKED TO NEXT BUF
IDM:			ID MESSAGE REQUEST
SYNC:		6 7	SYNC RD/WR IN SYNC
ERR:	EQU	/	FRROR IN BUFFER
;;			
:			
DSTDCB:	PUSH	н	;MOVE BUF ADDR
	POP		TO IY
			GET DST AC
			;GET DCB, (SP)=DCB
	POP	HL	;HL=DST DCB
	RET		;
;;			
;			
INITAQ:			;IX=AQUE (AVAILABLE BUFFERS)
			BC=64 (SIZE OF EACH BUF)
•	LD	HLABUF	;HL=ABUF (START OF BUF AREA)
; PTBF:	CALL	DHT	; PUT ADDR ON AQUE
FIDE			RETURN WHEN AQUE FULL
		HL,BC	INEXT BUF
	JR	PTBF-\$	CONTINUE TILL AQUE FULL
;;	0.11		
;			
DEVDCB:	ADD	A , A	; A = A + A
	LD	C . A	; C = O F F S E T
	LD	B,00H	;
	LD	HL, DCBTAB	;HL=DCBTAB ADDR
	ADD	HL,BC	;HL=DCBTAB+OFFSET
	LD		;DCB L
	INC	HL	;
	LD	B,(HL)	;DCB U
	POP		GET RETURN ADDR
	PUSH		SAVE DCB ADDR
	PUSH	ΗĻ	RESTORE RETURN ADDR
	RET		;
;;			
, HLDE:	LD	A,E	LOWER BYTE
11 L V C 🕯	<u> </u>	n / L	FLVWER DITE

SUB ; Ł LD AD JUPPER BYTE ;CARRY SET IF HL>DE SBC A / H RET ; ;; ; OFH **;CLEAR UPPER 4 BITS** HEXASC: AND A . 90H ;A=A+90H ADD DAA JDECIMAL ADJUST ; A = A + 4OH + CARRYADC A / 40H **;**DECIMAL ADJUST DAA RET **FASCII VALUE IN A** ; ;;; FILE QUE ; INPUT-OUTPUT QUEUE ROUTINES ; VERSION 1.0 REV C ; ROUTINE TO PUT AN ADDRESS INTO QUEUE. ; ; ENTRY- IX CONTAINS QUEUE CONTROL BLOCK ADDR ; AND HL CONTAINS DATA TO BE PUT ON QUEUE. EXIT- IF QUEUE IS FULL, RETURN WITH CARRY SET ; ; NORMAL RETURN IS WITH CARRY CLEAR. ; ;SET CARRY FLAG PUT: SCE BIT BSY (IX+0) FTEST FOR BSY QUEUE RET ΝZ ;YES, RETURN **;**TEST FOR FULL QUEUE BIT FULL (IX+O) RET NZ ;YES, RETURN $BSY_{\prime}(IX+0)$;SET BSY BIT SET RES EMTY, (IX+O) ; CLEAR EMPTY FLAG DEAHL ; DE = ADDRΕX LOWER BYTE BQP $L_{(IX+3)}$ LD LD H,(IX+4) JUPPER BYTE BQP ; (HL)/E **;**PUT LOWER BYTE LD INC ;NEXT BYTE HL LD (HL),D ; PUT UPPER BYTE INC HL ;HL=BQP+2 ; ; IF BQP=BB, SET BQP=TB CALL CHKBB ; LOWER BYTE TOP LD A (IX+1) CP L ;= L ;NO, QUEUE NOT FULL JR NZ, PUTR-\$;YES, UPPER BYTE TQP LD A (I X + 2) CP ;= H Н ;NO, QUEUE NOT FULL NZ, PUTR-S JR FULL (IX+O) ;YES, SET QUEUE FULL SET ; PUTR: LD (IX+3),L LOWER BYTE BQP LD (IX+4),H JUPPER BYTE BQP ΕX DE,HL ;HL=ADDR XOR CLEAR CARRY FLAG Α

BSY (IX+O) ; RESET BSY FLAG RES RET ;; ; ; ROUTINE TO GET ADDRESS FROM QUEUE ENTRY- IX CONTAINS QUEUE CONTROL BLOCK ADDR ; EXIT- HL CONTAINS DATA REMOVED FROM QUEUE. ; ; IF QUEUE IS EMPTY, RETURN WITH CARRY SET ; NORMAL RETURN IS WITH CARRY CLEAR. ; ;SET CARRY FLAG GET: SCF FTEST FOR BSY QUEUE BIT BSY (IX+0) ;YES, RETURN RET ΝZ EMTY (IX+0)**;**TEST FOR EMPTY QUEUE BIT RET ;YES, RETURN ΝZ SET BSY_{IX+0} **;**SET BSY FLAG RES FULL (IX+O) **;**CLEAR FULL FLAG LD L_{IX+1} LOWER BYTE TOP LD $H_{\prime}(IX+2)$ **;UPPER BYTE TQP** ; LD E,(HL) GET LOWER BYTE INC HL ;NEXT LOCATION GET UPPER BYTE LD D/(HL) INC HL ;HL=TQP+2 ; CALL CHKBB **;IF TQP=BB**, **SET TQP=TB** ; LD LOWER BYTE BQP $A_{\prime}(1X+3)$ CP L ;= L ;NO, QUEUE NOT EMPTY JR NZ/GETR-\$;YES, UPPER BYTE BQP LD $A_{\prime}(IX+4)$ CP ;= H Н ;NO, QUEUE NOT EMPTY JR NZ,GETR-S EMTY, (IX+O) ;YES, SET QUEUE EMPTY SET ; GETR: LD (IX+1)/L **;LOWER BYTE TQP** (IX+2).H **;**UPPER BYTE TQP LD ЕΧ DE.HL ;HL=ADDR CLEAR CARRY FLAG XOR A RES BSY (IX+0) FRESET BSY FLAG RET ;; ; CHKBB: LD $A_{r}(1X+7)$ LOWER BYTE BB CР L ;= L RET **;NO**, RETURN ΝZ ;YES, UPPER BYTE BB LD $A_{(1X+8)}$ CΡ ;= H H RET ΝZ **;NO**, RETURN LD L (I X + 5) LOWER BYTE TB LD $H_{(IX+6)}$ **;UPPER BYTE TB** RET ; ;;; FILE HBUFS

E33

;	-	ER STORAGE AND	QUEUES **********
,******			*****
	ORG	2000H	
;			
•;	HOST	XMIT AND RECE	IVE BUFFERS
;			
RBUF:	EQU	\$;RECEIVE BUFFER (FROM HOST)
	DEFB	04H	STATUS BYTE
	DEFW	0000н	BYTE COUNT
	DEFS	400H-03H	BUFFER STORAGE LOCATIONS
RBFN:	EQU	\$;END OF RBUF
;			
;			
XBUF:	EQU	\$;XMIT BUFFER (TO HOST)
	DEFB		STATUS BYTE
		0000H	
	DEFB		;
			; S Y N . S Y N
			; S Y N . S Y N
	DEFB		SOH
	DEFB		FC (ACK/NAK MSG)
XBSC:	DEFB		SC (41H OR 42H)
ADJU.	DEFB		· JAC
хвос:	DEFB		; OC (ACK, NO INST)
XBUC:	DEFB		;IC (41H+#MSG)
	DEFB		STX
	DEFB		; ETX
хвсс:	DEFB		BCC (BLOCK CHECK CHAR)
ADUL.		16168	SYN, SYN
L0G1:	DEFW		SOH (START OF FIRST LOGICAL MSG
LUGI.			BUFFER STORAGE LOCATIONS
XBFN:	EQU	\$ \$	END OF XBUF
;		.c.	
SBUF:	EQU	¢	SERVICE MESSAGE BUFFER
3001.	DEFB		STATUS BYTE
		000EH	BYTE CNT
	DEFB		
		1616н	SYN, SYN
		1616H	; SYN, SYN
	DEFB		; SOH
	DEFB		FC (SRV MSG)
S85C:	DEFB		SC (41H OR 42H)
3036:	DEFB		
SBOC:	DEFB		; OC (40H+ACK/NAK+SVM)
3000.	DEFB		JIC (1 MSG)
•	DEFB		STX (NO TEXT)
	DEFB		ETX
SBCC:	DEFB		;BCC (42H+SC+OC)
3011:	DEFB		;EOT
;	VERD	UH N	
	EOU	¢	Q-FRAME BUFFER
QFRM:	EQU DEFB		
			STATUS BYTE
	VEFW	0006H	;

•

.

DEFB OOH ; DEFW 1616H ; SYN, SYN DEFW 1616H ; SYN, SYN DEFB 01H SOH DEFB 04H ;EOT ;;; FILE QUES QUEUE CONTROL AND STORAGE BUFFERS ; ORG 281CH ;RECEIVE QUEUE EQU \$ RQUE: ;RQUE FLAG DEFB 04H DEFW TRB **;**TOP RQUE POINTER DEFW TRB **;**BOTTOM RQUE POINTER DEFW TRB ;TOP RQUE BUFFER DEFW TRB+20H BOTTOM RQUE BUFFER DEFS 20H **;**RQUE BUFFER TRB: ; XQUE: EQU **FRANSMIT QUEUE** \$ DEFB 04H ;XQUE FLAG DEFW TXB **;**TOP XQUE POINTER DEFW TXB ;BOTTOM XQUE POINTER DEFW TXB ;TOP XQUE BUFFER DEFW TXB+20H BOTTOM XQUE BUFFER DEFS 20H XQUE BUFFER TXB: ; EQU **;**AVAILABLE BUFFERS QUEUE AQUE: \$ DEFB 04H ;AQUE FLAG DEFW TAB ;TOP AQUE POINTER **;**BOTTOM AQUE POINTER DEFW TAB DEFW TAB **;**TOP AQUE BUFFER BOTTOM AQUE BUFFER DEFW TAB+20H TAB: DEFS 20H ;AQUE BUFFER ; ;WAIT QUEUE WQUE: EQU \$;WQUE FLAG DEFB 04H DEFW TWB **;**TOP WQUE POINTER DEFW TWB ;BOTTOM WQUE POINTER DEFW TWB **;**TOP WQUE BUFFER DEFW TWB+20H ;BOTTOM WQUE BUFFER TWB: DEFS 20H ;WQUE BUFFER ; **;**ERROR QUEUE EQUE: EQU \$ DEFB 04H ;EQUE FLAG DEFW TEB **;**TOP EQUE POINTER DEFW TEB ;BOTTOM EQUE POINTER DEFW TEB ;TOP EQUE BUFFER DEFW TEB+20H BOTTOM EQUE BUFFER DEFS 20H ;EQUE BUFFER TEB: ; ;;; FILE STRG ; GENERAL STORAGE AREA ;

E 3 5

HDRID: DEFB OOH ;HEADER ID SAVE AREA ; SAVSC: DEFB OOH SAVE PRESENT SC (RBUF) DEFB OOH SAVE PRESENT IC (RBUF) SAVIC: **;**SAVE PRESENT OC (RBUF) SAVOC: DEFB OOH SAVE MSG LENGTH & FLAG (RBUF) DEFW 0000H SAVMSG: DEFW OOOOH SAVE MSG LOCATION DEFW 0000H SAVE BUF LOCATION ; DEEW OOOOH SAVE LOCATION POINTER (XBUF) SAVLOC: SAVBCC: DEFB OOH ;SAVE NEW BCC (XBUF) ; LEVEL 3 INTR MASK DEFB OFFH MASK: ENDCHR: DEFB ODH ;BUFFER TERMINATING CHAR <CR> ; ORG 2900H ;AVAILABLE BUFFER AREA ABUF: EQU \$ DEFS 10H*40H ;16 BUFS * 64 BYTES PER BUF ; IVEC: EQU \$ **;INTERRUPT VECTOR TABLE** DEFW INTO **;INTERRUPT LEVEL O** DEFW INT1 **;INTERRUPT LEVEL 1 ;INTERRUPT LEVEL 2** DEFW INT2 DEFW INT3 JINTERRUPT LEVEL 3 ; DEFW DCBOO ;HOST DCB (HCB) INTTAB: DEFW DCB01 CONSOLE DCB DEFW DCB01 CONSOLE DCB DEFW DCBO2 ;RMC1 DCB DEFW DCBO2 ;RMC1 DCB ;RMC2 DCB DEFW DCB03 DEFW DCB03 ;RMC2 DCB DEFW DCB01 **;**EXTERNAL INTR DCB DEFW DCBOO ; DEVOO DCB (HOST) DCBTAB: ; DEVO1 DCB (CONSOLE) DEFW DCB01 DEFW DCBO2 ; devo2 dcb (rmc01) ;DEVO3 DCB (RMCO2) DEFW DCB03 DEFW DCBXX ;DEVO4 DCB DEFW DCBXX ;DEVO5 DCB **;**DEVO6 DCB DEFW DCBXX ;DEVO7 DCB DEFW DCBXX ; DCBXX: EQU OFFFFH JUMMY DCB ; HCB: EQU \$;HOST DCB = DCBOODCBOO: DEFB 02H **JDEVICE STATUS FLAGS** SAVE BYTE COUNT DEFW 0000 SAVE BUFFER POINTER DEFW 0000 ; DCB01: DEFB O2H JEVICE STATUS FLAGS DEFB 02H **;**DEVICE MASK DEFB 10H ;DEVICE # DEFB 01H ;AC (LU#)

;	DEFB 00H DEFW 0000 DEFB 0000 DEFW 0000 DEFW 0000	;DST AC (READ ONLY) ;SAVE MAIN BUF ADDR ;SAVE REMAINING BYTES ;SAVE DE (STATUS, TOTAL BYTES) ;SAVE HL (POINTER)
		DEVICE STATUS FLAGS
DCBUZ:	DEFB 02H	DEVICE MASK
	DEFB 14H	DEVICE #
	DEFB 02H	JAC (LU#)
		DST AC (READ ONLY)
	DEFW 0000	SAVE MAIN BUF ADDR
	DEFB 0000	
		;SAVE DE (STATUS, TOTAL BYTES)
	DEFW 0000	;SAVE HL (POINTER)
;		
DCB03:	DEFB O2H	DEVICE STATUS FLAGS
	DEFB 20H	DEVICE MASK
	DEFB 16H	;DEVICE #
	DEFB O3H	;AC (LU#)
	DEFB OOH	;DST AC (READ ONLY) ·
	DEFW 0000	SAVE MAIN BUF ADDR
	DEFB 0000	;SAVE REMAINING BYTES
	DEFW 0000	;SAVE DE (STATUS, TOTAL BYTES)
	DEFW 0000	;SAVE HL (POINTER)
;		
	END 0000	END OF RNP

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I. Getting on line with TSS

The modem line can be in any one of several states, depending on what has happened since it was last used. The following routine will usually get the time sharing system connected, regardless of the state the line is presently in.

- 1) Turn on power to CPU and CONSOLE. (Green light on CPU should light if power is ok.)
- 2) Press lower push button on CPU panel to reset CPU. (Illustrated in App. A) The following message should appear on the console: <u>Zee OS V1.2.</u> If it does not appear, press button again.
- 3) You are now in the monitor program (described in Chapter 2). Next press "K" on the console keyboard. This should cause the CPU to enter the HOST I/O routine (described in Chapter 3). The following message should appear on the console: <u>ZIO+</u>. If it does not appear, go back to step 2.
- 4) You are now ready to connect to the TSS system. Press "CR" on the console. The TSS system may answer with the sign-on message. If so, continue as if you were on a dial-up terminal, (Refer to Honeywell Time-Sharing System Pocket Guide, BS12, page 13, TSS Terminal Operation.)
- 5) If there is no response from the previous input, press a "CTL A". The system may respond with: <u>Program Name?</u> If so, answer with <u>TSS</u>, and continue like a dial-up terminal.

- 6) If there is still no response, press the "ESC" key on the console and go back to step 4. This serves as the <u>"break"</u> key on this system.
- 7) If, after repeating the above procedure (from 4 to 6) several times, there is still no response from the HOST system, disconnect the power to the data modem and reconnect (this should reset the data line to the HOST system). Then go back to step 4 and start over.
- 8) If none of the above procedures get a response from the HOST system, then there is probably something wrong with the HOST system. In this case, you should check with the Computing Center on the status of TSS.

II. Loading and debugging programs

Since the HOST system does not require continous I/O to remain connected, you can jump from the I/O routine back into the monitor and then back into the I/O program and still remain connected to TSS. This capability allows you to load programs from the HOST, jump to the monitor for debugging and then jump back to the I/O program without disconnecting from TSS. However, you must remember that TSS will time out if it sees no input within 10 minutes. Therefore, you should go into the IDLE mode if you plan to be off longer, or sign off completely. To load a program from the HOST, the program must be in absolute HEX format. (This format is explained in the Zapple Monitor Operating Manual.) The following shows how to load a program from the HOST and then go to the monitor.

- 1) You must be in the I/O program and connected to TSS.
- 2) Press "CTL B" on the console and the console will return a ">".
- 3) You must now type in an offset value in HEX, or if no offset is required, just press "CR".
- 4) You now type in the name of the file which holds the assembled program, followed by a "CR". The console will then print <u>LIST <file></u>, after which the HOST system will start sending the program code. The program code is both placed in memory and printed on the console so you can tell if an error occurs. (ERROR recovery will be explained later.)
- 5) After the program has been loaded, you may press "CTL N" to return to the monitor. .
- 6) While in the monitor you may run the program just loaded, display the code, modify the code, etc.(All monitor functions are explained in Appendix B and the Zapple Monitor Operating Manual.)
- 7) When finished with the monitor, you press "K", which will go to the I/O program, and you can use the HOST system without signing back on to TSS.

ERROR recovery - if an error is detected during the the file transfer, the program will exit to the monitor routine and print "*" to indicate an error. Any non-Hex character detected during the transfer of a block or a bad checksum will cause an error exit. If an error is indicated it means you have returned to the monitor program. Therefore, you must re-enter the HOST I/O routine by entering a "K" from the keyboard before you can continue with TSS operations.

III. Creating and assembling microprocessor programs.

Programs for the microprocessor are written on the HOST system using the standard text editor. They must be written in standard ZILOG or MOSTEK Z-80 neumonics and include only the pseudo ops given in the assembler instruction manual (MOSTEK XFOR-80 CROSS ASSEMBLER MANUAL).

The file created with the editor is used as the input file to the assembler (XFOR-80), which is run using the standard FORTRAN system on the HIS 66/60. The input file number is 05, the line printer output file is 06, and the assembled Hex code file is 03. Two temporary files, 02 and 04, are also generated, but are of no use and need not be saved.

The following command is the format for assembling a file called MYSRC (Z80 SOURCE CODE), putting the listing into a file called MYLP (LINE PRINTER LISTING) and putting the assembled code into a file called MYCODE (Z80 HEX CODE). RUNY Z80#MYSRC"05"; MYLP"06"; MYCODE"03"

The results of the assembly can be displayed on the console, using LIST or EDIT, to check for errors, etc.

Naturally, since the assembler is written in FORTRAN, it can also be run under batch and the I/O files can be any medium acceptable to the FORTRAN system. (For detailed information on the FORTRAN system, refer to Honeywell FORTRAN Pocket Guide, DD82.)

IV. Console Baud Rate Considerations

Since the I/O program is run in real time and uses no buffering, the console must not be run slower than the modem link to the HOST system. In fact, to prevent loss of characters, it should be faster than the modem speed (ie. if the modem is set at 1200 baud, run the console at 2400 baud).

The console baud rate is selected on the CPU by the rotary switch on the AUX card (illustrated in App. A). The setting here should match the setting on the rear of the console.

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