A FEASIBILITY STUDY IN APPLYING READ-ONLY MEMORIES TO THE ANALYSIS OF NONLINEAR ANALOG FUNCTION RESPONSES

A Thesis

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Presented to

the Faculty of the Department of Electrical Engineering

University of Houston

In Partial Fulfillment

of the Requirements for the Degree Master of Science in Electrical Engineering

by

Edgar Lee Dohmann December, 1972

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ABSTRACT

The measurement of physical qualities such as temperatures, pressures, and flow rates is both desirable and essential in many industrial processes. Devices such as thermocouples and other transducers have been developed which yield a predictable voltage or current response to the external stimulus. The use of such devices which produce voltage and current signals provide economical and convenient capabilities to measure the desired physical qualities because many signals may be terminated and monitored at a central location such as a control room.

Thermocouples for temperature measurement are the most widely used devices in such industrial process applications because of their simplicity, accuracy, and low cost. As most such devices which respond to an external stimulus, the relationship between the voltage generated between thermocouple junctions and the temperature of the junctions is nonlinear.

This nonlinear relationship poses the greatest problem in designing and constructing accurate temperature indicating devices.

Several analog-to-digital conversion techniques have been developed which compensate for the nonlinear function being measured and produce temperature readouts in proper engineering units. All techniques developed for commercial applications strive to attain satisfactory degrees

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of speed, simplicity, accuracy, economics, and reliability. Most techniques offer some advantages in two or more of these qualities but no presently available techniques offers a clear advantage in all five qualities.

Recent developments in the field of semiconductor memory devices have made it possible to design a digital temperature indicator using read-only memories (ROMs) with accuracies approaching those attainable with computer based techniques. The speed, simplicity, and reliability of such a ROM technique would be better than any other presently available technique. Predicted trends in future semiconductor developments also make such a technique very attractive economically.

A computer program has been developed to generate ROM truth tables for a nonlinear response system. The truth tables generated by this program are based on the system of quadratic equations which describe the nonlinear function. The accuracies are limited only by the inherent accuracy of the equations used, the size of the ROM, and the dynamic range of the stimulus to be measured.

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CHAPTER I

INTRODUCTION

Many natural processes are neither linear nor quantitative but behave in continuously smooth, nonlinear ways. It is desirable and necessary to measure such nonlinear natural processes in many applications such as industrial process control and operation, laboratory experimentation, and data analysis.

Many devices have been developed which yield a predictable voltage or current response to an external stimulus. Thermocouples are widely used as temperature transducers because of high inherent accuracy, wide measurement range, fast thermal response, ruggedness, reliability, and low cost. Thermocouple circuits develop an electromotive force (emf) at the measurement terminals whose magnitude and polarity depend on properties of the metals used and the temperature difference between the thermocouple junctions.

Although many materials can be combined to produce a thermoelectric effect, certain pairs of metals have become standard. The four most commonly used types of thermocouples are: Type J (Iron/constantan), Type T (copper/ constantan), Type K (chromel/alumel), and Type E (chromel/ constantan). The type used depends on such factors as desired operating range, external environment, desired emf output range, and installation considerations. As with most transducer devices which respond to an external stimulus, the relationship between the voltage generated between thermocouple junctions and the temperature of the junctions is nonlinear. Several techniques have been developed to measure thermocouple emf, compensate for the nonlinear function, and produce temperature readouts in proper engineering units. All techniques developed for commercial applications strive to attain satisfactory degrees of speed, simplicity, accuracy, economics, and reliability.

The purpose of this investigation was to evaluate the feasibility of using read-only memory (ROM) devices in the linearization of such quantities as thermocoupleproduced emf and develop techniques of applying ROMs to these tasks. While many of the principles and techniques presented and discussed in this report apply to measurement and analysis of nonlinear functions in general, this investigation was limited to the primary field of thermocouple circuits due to their common and widespread use and the availability of background information.

CHAPTER II

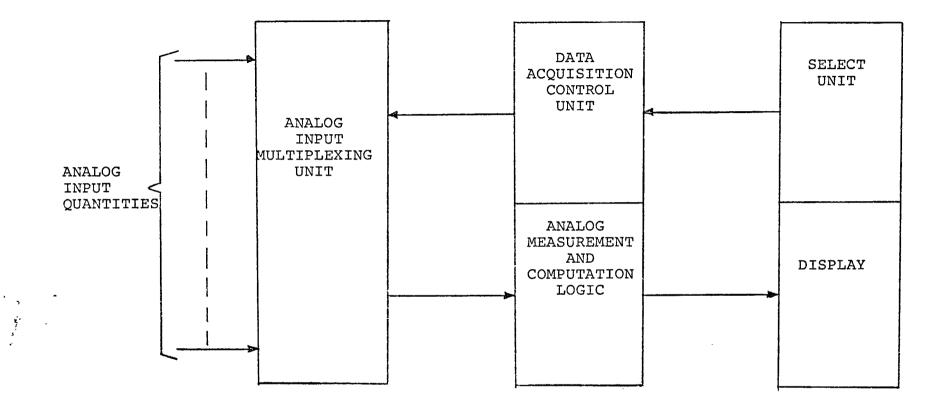
CONVENTIONAL LINEARIZATION TECHNIQUES

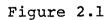
The simplest method of measuring temperature using thermocouple circuits is to use an analog or digital voltmeter to measure the emf produced and then to refer to a National Bureau of Standards (NBS) thermocouple conversion table to determine the corresponding temperature.

While this manual table-look-up method may be simple it can also be extremely time consuming and impractical in industrial process applications where hundreds or thousands of thermocouple circuits may be terminated in a single control room and it is often nesessary to read and/or record many temperature values in a small amount of time.

Figure 2.1 shows a generalized block diagram of a data acquisition system. Such a system permits an operator to select any one of many input points to be monitored and displayed. The control logic of the data acquisition system multiplexes the selected input point onto a single analog channel then measures the selected quantity and displays the resultant value to the operator.

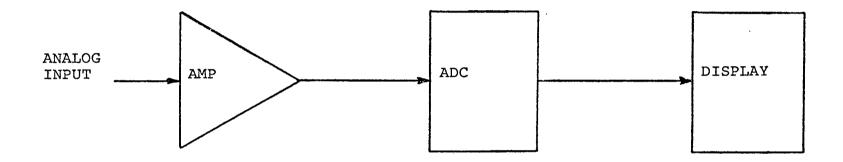
The generalized block diagram is simplified even further in Figure 2.2. This figure shows the basic component parts of the Analog Measurement/Computation Logic and the resultant display output. Figure 2.2 shows only one analog input value as would be the case after the selected point is

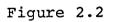


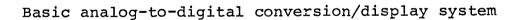


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Typical data acquisition system







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multiplexed onto a single analog input channel. This basic block diagram is referred to throughout the remainder of this chapter as some variation of this block diagram is used in almost every conventional linearization technique.

The two most basic means of linearizing an input function are by analog techniques or by digital techniques although some techniques employ a combination of both. It is not the intent or purpose of this chapter to present and discuss every linearization technique that has ever been developed but rather to discuss in general terms the characteristics of some of the most commonly used techniques to provide a basis for comparison of the read-only memory linearization technique to be discussed in a subsequent chapter.

The basic goals of a linearization technique should be to provide a system which is accurate, simple, economic, fast, and reliable. Each of these qualities is used as a basis for comparing the desirability of one technique over another. Often trade-offs must be made between two or more of these qualities to optimize the adherence to another.

Accuracy, when applied to linearization techniques, normally refers to the ability of the linearization network or device to conform to the published NBS tables without regard to the inherent inaccuracies of the other system elements such as the sensor itself, the signal conditioning

amplifier, or the analog-to-digital converter.

Simplicity of a linearization technique includes the space required by the technique, the maintenance required, the versatility of the technique, and the difficulty to design and implement the technique. Versatility is very important because some applications require the use of a particular type of ADC or amplifier but if a linearization technique is highly dependent on some other type of ADC or amplifier, it may not be possible to use that technique.

Economics of a linearization technique includes both the cost of the components used and the cost of implementing the technique.

The speed of a linearization technique refers to the time required for the measured quantity to be adjusted to conform to the actual value of the stimulus.

The reliability of a linearization technique includes the expected life of the components used in the technique and the ability of the technique to continue to conform to the published NBS tables over an extended period of time.

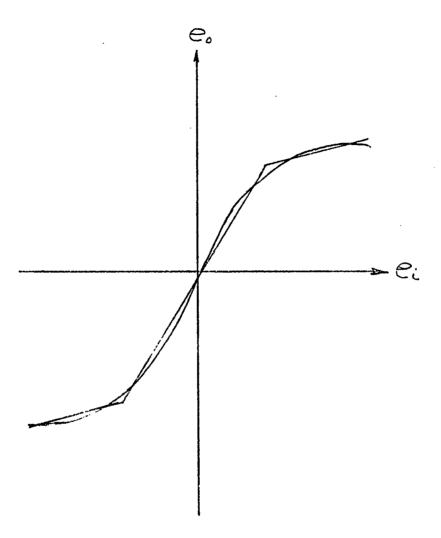
The simplest linearization technique based on using a circuit like that shown in Figure 2.2 would be to scale the displayed data to a value that could be conveniently referenced in an NBS look-up table. This technique would be similar to that mentioned at the beginning of this chapter with the exception that automatic analog multiplexing

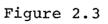
is implied by the system shown in Figure 2.2. Even though automatic signal multiplexing is provided by this technique, the procedure of conversion by reference to a table is not simple and is certainly not reliable since each reference is subject to human error. The method is also very slow and costly in terms of time and effort expended. Since all five of the desired qualities are severely lacking in this method it is readily evident why automatic linearization techniques are preferred over manual methods.

AMPLIFIER, DIODE FUNCTION GENERATOR FEEDBACK TECHNIQUE

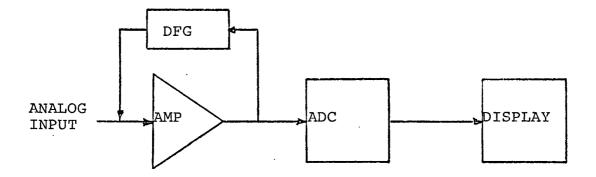
One of the most commonly used linearization techniques is the Amplifier DFG feedback technique. This technique uses the very popular piecewise linear approximation as do most other conventional linearization techniques. The accuracy of any method using piecewise linear approximation is determined by the number of line segments used and the inherent accuracy of the circuitry generating the line segments. Figure 2.3 illustrates the basic piecewise approximation of a nonlinear curve.

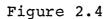
The Amplifier DFG feedback technique employs a diode function generator (DFG) network in the amplifier feedback circuit to generate a function which is the inverse to the input voltage. The block diagram of such a system is shown in Figure 2.4 and the typical response curves are shown in Figure 2.5.



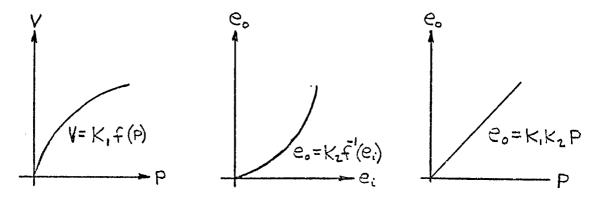


Piecewise approximation of a nonlinear function





Amplifier diode-function-generation feedback technique



(a) Transducer output (b) DFG function

(c) Composite transfer curve

Figure 2.5

Linearization of curves by a DFG circuit

The complexity of such a system depends on the number of line segments used and the capabilities provided for temperature-compensation of the breakpoint diodes. Precision limiter circuits using operational amplifiers to simulate ideal diodes can greatly improve the accuracy but would also greatly increase the cost and complexity of the system. This method has the significant advantage of being very versatile and essentially independent of the type of amplifier and ADC used. This system usually depends very much on the reference voltage used to bias the DFG circuit and, normally, each segment of the curve requires a separate adjustment because of inherent variations in diode characteristics. This method has the advantage of being very fast since the time required for the output voltage from the amplifier to stabilize depends only on such factors as amplifier slew rate and settling time and diode switching time.

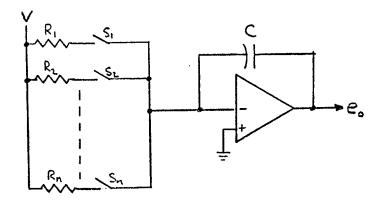
RAMP LADDER LINEARIZATION TECHNIQUE

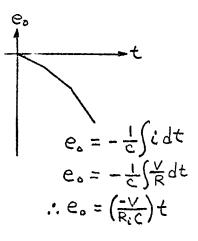
One analog-to-digital conversion technique is to allow a counter to accumulate pulses at a fixed rate while a ramp generator is running; then, when the ramp output matches the unknown input signal, the total count accumulated is proportional to that input signal. If the input current to the ramp generator is changed in steps depending

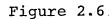
on the number of pulses counted, a piecewise linear ramp can be generated as shown in Figure 2.6. If the polarity and magnitude of the ramp is properly calculated, the resultant ramp can approximate the transducer response curve as shown in Figure 2.7.

Figure 2.8 illustrates how such a nonlinear ramp generator might be used in a linearizing data acquisition system. The basic accuracy of such a system depends on the number of line segments used and the inherent accuracies of the components used. Precision components can be used in the ramp ladder circuit to simplify implementation and calibration and to increase reliability. The cost also depends on the desired accuracy and simplicity of calibration requirements.

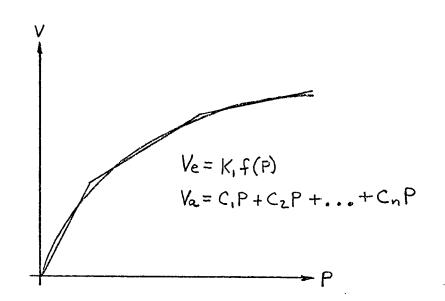
The two major disadvantages of the ramp ladder linearizer are speed and flexibility. The ADC counter must start at zero at the beginning of the conversion cycle and must count up to the unknown value. For a full scale analog input voltage 2ⁿ-1 clock pulses are required before the conversion is complete (n is the number of bits in the digital word). The frequency of the clock pulses driving the ADC counter is limited also by the response time of the switches in the ramp ladder circuit and the slew rate of the ramp generator.

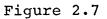




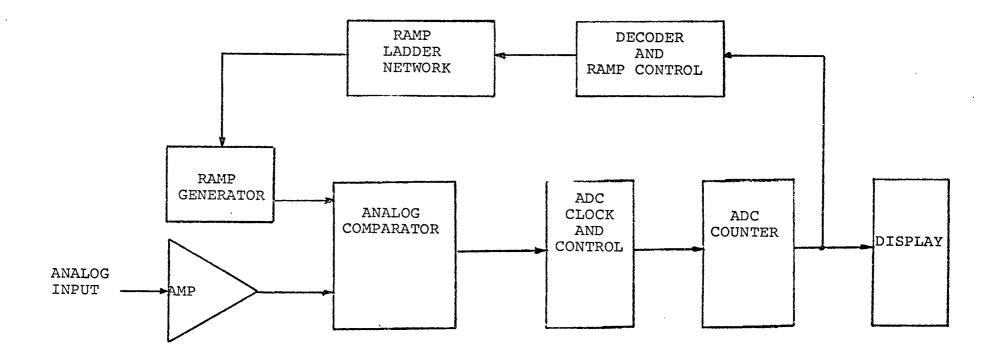


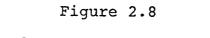
Typical nonlinear ramp generator





Piecewise linear approximation of transducer output





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Ramp ladder network technique

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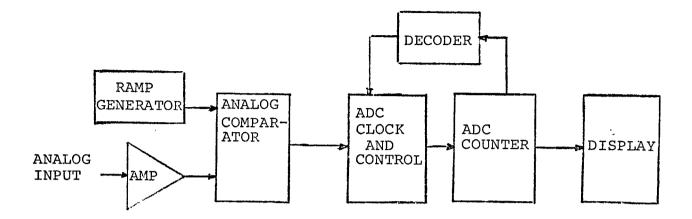
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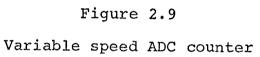
The ramp ladder linearizer technique requires that a counter-ramp type of ADC be used. This type of ADC must be either a specially designed ADC or one which provides access to the digital count during the conversion cycle and access to the ramp generator input so the decoder/ramp control circuitry and ramp ladder circuitry can be added externally.

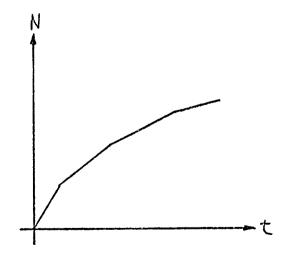
VARIABLE SPEED ADC COUNTER TECHNIQUE

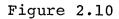
Another linearization technique using the basic counter-ramp ADC is the variable speed ADC counter. This is a digital linearization technique as shown in Figure 2.9 using a fixed ramp to control the conversion time but using more than one clock frequency to control the speed of the ADC counter. The resultant output of such a device is shown in Figure 2.10. The number of pulses accumulated in the ADC counter varies with time to approximate the response function.

The accuracy and complexity of the circuitry used in this technique depends on the number of clock frequencies available and the complexity of the decoder circuit used to control the selection of the frequency to drive the counter. The theoretical accuracy obtainable by this method is dependent also on the number of line segments used in the piecewise approximation. The number of precision components needed to make a circuit which is easy to calibrate and









Piecewise linear approximation by varying counter speed

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maintain is much less than required by the ramp ladder technique. A severe limitation of this technique is the requirement that the ADC be specially designed for this application. Another limitation is the speed of the system. In general, however, it should be possible to operate the ADC counter for this method at faster speeds than for the ramp ladder method because the response time of the digital switches controlling the counter pulse frequency is less than the response time of the analog switches on the ramp ladder network. A one-digit ambiguity exists each time the oscillators are switched if discrete oscillators are used to vary the ADC counter speed. If a resistor is switched in the frequency determining network of a single oscillator, this ambiguity does not arise.

ACCUMULATED ADC COUNT CORRECTION TECHNIQUE

Another digital linearization technique accomplishes the same overall effect in a slightly different manner. This method which corrects the accumulated pulse count of the ADC counter has essentially the same general composition as the method shown in Figure 2.9. A fixed ramp generator produces a linear output until a comparison is made with the unknown input. During the conversion cycle the ADC counter is driven at a constant rate but when the decoder/control circuitry detects a "correction point", a pulse is either

added to or subtracted from the total accumulated count.

The counter correction circuit can either increment or decrement the ADC counter when the correction point is detected if an up/down counter is used or the total of all additions and subtractions can be maintained until the end of the conversion cycle and the final accumulated value corrected before presentation to the display. The net effect of either method of correcting the accumulated value of the ADC counter is to approximate the response curve by shifting line segments with the same slope.

The major advantage of the two digital techniques is that for the same or less degree of complexity and cost than analog techniques, a piecewise approximation can be realized with many more line segments thereby yielding a more accurate device.

The reliability of the digital techniques discussed here is very good since highly accurate, stable, and reliable components are available to construct these devices and very little calibration of the system would be required.

Even though the ADCs discussed for the counterramp conversion techniques used a ramp generator to compare with the unknown input the dual-slope ADC technique could be applied to all 3 linearization methods discussed here. The dual-slope method employs an integrating ramp on the analog input to charge a capacitor to the unknown value. The

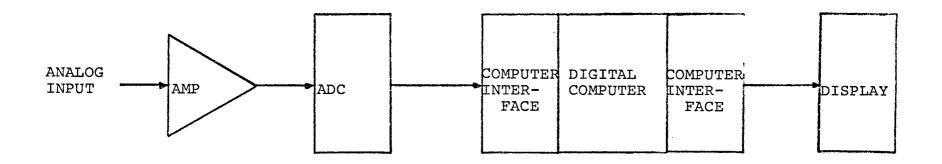
capacitor is then discharged at a constant rate until the unknown voltage reaches zero and the accumulated pulses in the ADC counter during this discharge time is proportional to the unknown input. If the unknown value were discharged at a non-linear rate or if the clock pulse frequency were varied, or if the accumulated value were adjusted by appropriate values this ADC technique could be used in essentially the same manner as the counter-ramp method.

DIGITAL COMPUTER PROGRAMMING TECHNIQUES

Another commonly used technique is to use a digital computer program to perform the linearization. Such a method as shown in Figure 2.11 would merely examine the digitized value of the analog input and either evaluate a series of polynomial or straight-line equations describing the response function or use a table look-up technique.

The computer program linearization offers the advantage of being very versatile. The method is essentially independent of the type of amplifier and ADC used.

To justify the computer program linearization approach economically, however, the computer would have to be an inherent part of an overall system and the measured data would have to be required by other programs in the overall system. The cost of a stand-alone system as shown in Figure 2.11 would, in general, be prohibitive.



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Figure 2.11

Computer program linearization

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In addition to the conversion time of the ADC used in the method of Figure 2.11, the device access time and program execution time would be prime factors in the overall speed of the system. In general a table look-up technique would be faster than a polynomial equation evaluation technique but would also occupy more memory space.

The straight forward processing technique of the computer program linearizer presents a very simple design concept but the actual implementation of the system may be very complicated depending on the structure and purpose of the overall system.

Of all the commonly used linearization techniques presented here, none possesses extreme advantages over any other. Perhaps the most commonly used techniques in the past have been manual table-look-up and DFG feedback circuits. The ramp ladder and variable clock techniques have been used when simpler circuitry and less calibration were desired. The computer program technique has been widely used in systems in which computers are employed for other uses. The digital techniques have become increasingly popular in the last two years because continuing advances in digital technology make the circuitry required to support such techniques simpler and more economical to use.

Indications are that the application requirements will continue to dictate the technique used but continuing

advances in digital technology should make digital techniques more and more attractive in the future.

Recent innovations in economical modular analog circuitry components such as power supplies, multiplexers, sample and hold devices, analog-to-digital converters, digital-to-analog converters, and amplifiers make it increasingly more attractive to search for linearization techniques which are as independent of the other circuit modules as possible.

CHAPTER III

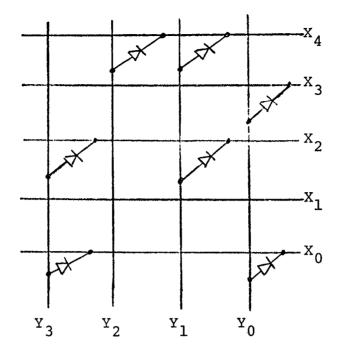
READ-ONLY MEMORY CHARACTERISTICS AND APPLICATIONS

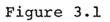
A read-only memory (ROM) is a circuit which can accept a digital code at its input terminals and provide a unique digital code on its output terminals. The relationship between the input and output codes is fixed and is usually alterable only by replacing all or part of the circuitry.

Many advances in read-only memory technology have been made since the mid-1960's and although ROMs have received much publicity recently, they certainly are not new devices. Read-only memories have been used in digital circuitry for many years but their usage in the past has been rather limited.

The most common form of read-only memory is the diode matrix which has been used since the early days of electronic digital computers. A typical example of such a diode matrix is shown in Figure 3.1. The matrix shown has 4 input bits and 5 output bits which makes the circuit a 20bit ROM, meaning that the circuit has 20 memory calls which may be programmed with a logic "0" (no diode) or a logic "1" (diode present).

The most common use of such diode matrix circuits has been the implementation of Boolean logic equations. The truth table shown in Figure 3.2 illustrates how the matrix





Diode matrix read-only memory

INP	UT E		ATTERN		PUT	BIT	PATT	ERN
^Ү з	¥2	Yl	Y ₀	×4	^х з	^X 2	xı	х ₀
0 0 0 0 0 0 0 1 1 1	0 0 0 1 1 1 0 0 0	0 0 1 1 0 0 1 1 0 0 1	0 1 0 1 0 1 0 1 0 1 0	 0 0 1 1 1 1 1 0 0	0 1 0 1 0 1 0 1 0 1 0	0 0 1 0 0 1 1 1 1 1	0 0 0 0 0 0 0 0 0 0 0 0 0	0 1 0 1 0 1 0 1 1 1 1
	0 1 1 1 1	1 0 0 1 1	0 1 0 1 1	1 1 1 1 1	0 1 0 1 0 1		0 0 0 0	1 1 1 1 1

Figure 3.2

A read-only memory truth table

of Figure 3.1 solves the following set of logic equations:

$$X_0 = Y_0 + Y_3$$
$$X_1 = 0$$
$$X_2 = Y_1 + Y_3$$
$$X_3 = Y_0$$
$$X_4 = Y_1 + Y_2$$

Such diode matrix circuits have also been popularly used in code-conversion circuits such as BCD-to-decimal, binary-to-BCD, and BCD-to-binary. Each of these examples and many others qualify as read-only memory circuits since a unique digital code can be applied to the input terminals and a unique digital code obtained on the output terminals.

Typical applications for devices which can generate a group of unique output bits for every group of unique input bits include implementing logic and counting functions, performing code conversion, function generating, character generating, micro-programming of digital processes, and performing table look-up functions.

Many devices which qualify as read-only memories have been developed to perform these tasks, although most of them, such as the diode matrix, are not commonly referred to as ROMs. A very popular technique is to store the bit pattern of a particular function in a common read-write memory (core memory, semiconductor memory, drum storage, etc) then activate a write-disable circuit to prevent inadvertent destruction of the stored information but still permit access to the stored information. Such a device would properly be classified as a ROM while the write-disable circuit remained active but would still provide the option of changing the contents of the memory at a later time.

Even though many devices qualify as read-only memories, the term "read-only memory" or ROM, is in general used to refer to semiconductor integrated circuit devices satisfying the definition of a read-only memory. Semiconductor memories incorporate high-gain and high-speed digital elements. Consequently they are a faster and more ideal storage device than their magnetic counterparts which depend on the analog properties of a slower device and exhibit no gain.

Memories built with semiconductor devices are becoming cost-competitive with ferrite-core memories. In addition they offer superior performance and greater freedom of organization. Many types of semiconductor memories, based on both metal-oxide silicon (MOS) and bipolar technologies, are becoming available. These memories differ primarily in cost, organization, and performance.

MOS technology is presently the most popular for the fabrication of read-only memories. The high density possible in this memory matrix cannot be achieved by any other present technology. In addition the circuitry required for address decoding and input-output buffering can be included

on the same substrate as that which contains the memory matrix. The result is a very dense, low-cost array that can be used easily.

A general diagram of a MOS read-only memory is shown in Figure 3.3. The ROM is organized as a 16-word 2-bitsper-word memory. The basic storage element of the MOS ROM is the MOS transistor. The ROM shown in Figure 3.3 has 32 possible transistor locations or memory "cells". A logic "1" is coded by inserting a transistor in a particular position and a logic "0" by omitting the transistor. The truth table for this ROM is illustrated in Figure 3.4.

Semiconductor memories (except read-only memories) are, in general, volatile storage devices. For a volatile storage device to retain stored data, power must be applied continuously. Read-only memories, while difficult to alter, eliminate this storage volatility problem. Most ROMs are fabricated by using a photographic mask with a customized memory matrix on it which stores the desired bit patterns. ROMs manufactured in this manner must be replaced if the bit pattern is to be changed later. Some ROMs are termed "programmable" read-only memories (PROMs) meaning that they can be programmed after the integrated circuit wafer has been packaged. Most PROMs are manufactured with logic "1"s in all bit positions and the programming consists of removing certain "1"s by applying a large reverse-biased voltage

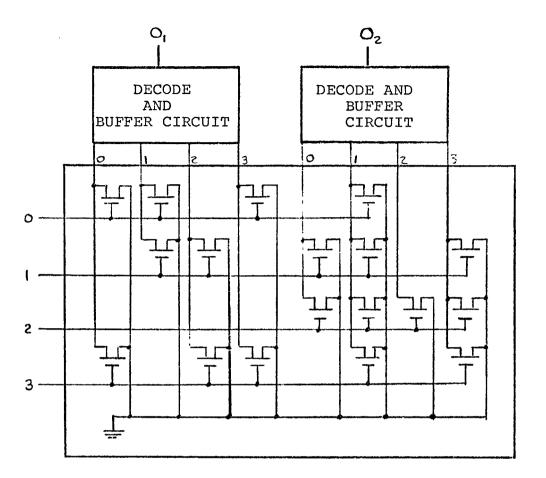


Figure 3.3

MOS read-only memory (16 word x 2 bit)

ROM WORD NO.	01	0 ₂
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16	1 0 1 0 1 0 0 0 0 1 0 1 0 1	0 1 0 1 1 1 1 1 1 1 1 0 1 0 1 0

Figure 3.4 MOS read-only memory truth table

to certain terminals in order to destroy the corresponding transistor. Some PROMs are even erasable (E-PROMs) by x-rays or ultra-violet rays. Applying such energy sources to an E-PROM returns all bit positions to logic "1" permitting the device to be re-programmed.

Some factors significantly affecting the cost of ROM devices are the type, size, method of programming, and quantity purchased. All of these factors must be considered when purchasing ROMs because an absolute comparison between devices cannot be made unless these factors are defined. MOS memories are, in general, less expensive than bipolar memories and can be obtained with more memory capacity in the same physical area. Memory costs are often compared on a price-per-bit basis. Read-only memories in the 1024 to 4096 bit range may be priced as low as 1¢ per bit in guantities of 1000 devices but may cost 5¢ per bit in unit quantities. Also a ROM in the 64 to 256 bit range may cost more per bit than a 1024 to 4096 bit ROM, because the actual device costs may be approximately equal. Another factor greatly affecting the cost of ROMs is the programming of the ROM. ROMs which are produced by photomasks are in general less expensive than equivalent ROMs which are PROMs or E-PROMs. If a custom mask must be made for a particular application, about 50 to 100 such ROMs would have to be produced from the mask to bring the total overall cost below that of a PROM

or E-PROM. In general, for very small quantities of custom memory patterns the PROM or E-PROM is the most economical. It is necessary to evaluate the actual costs and the projected requirements of a particular memory configuration to determine the most economical type of ROM to use in a particular application.

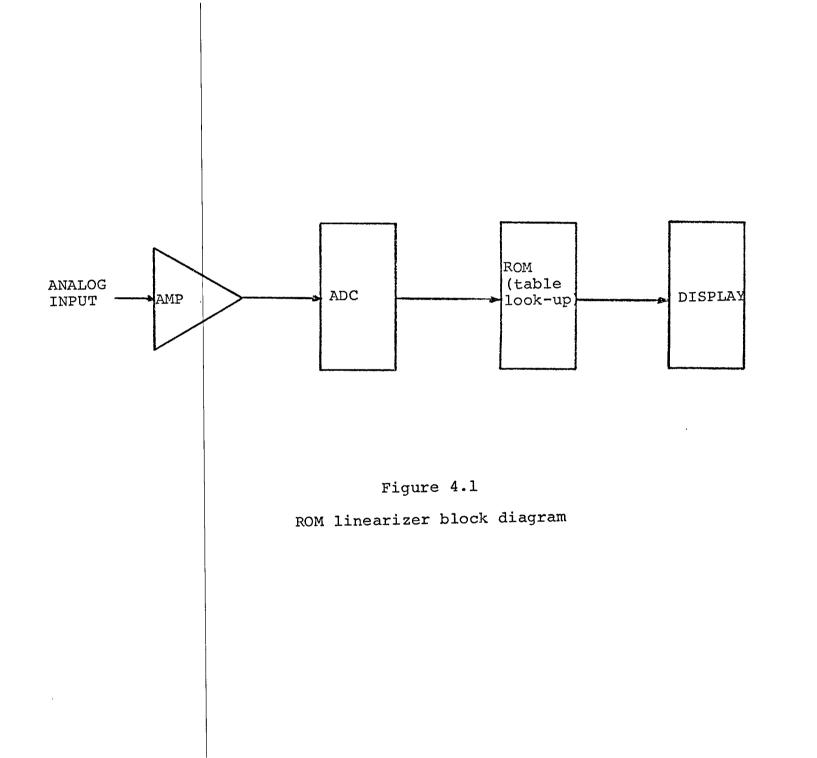
CHAPTER IV

READ-ONLY MEMORY LINEARIZATION APPLICATION

The basic characteristic of read-only memories (ROMs) is that a unique digital code on the input terminals results in a unique digital code on the output terminals. This makes them extremely useful for performing code conversion and table look-up functions. One of the many applications of such functions is the conversion of digitized values of nonlinear analog functions to suitably coded engineering units.

By effectively programming a NBS thermocouple table into a ROM, a temperature can be determined by measuring the emf produced by a thermocouple. Referring to an NBS "manually" (as discussed in Chapter 2) can be performed automatically. A simplified block diagram of such an automatic device is shown in Figure 4.1. If in such a device as shown in Figure 4.1, the amplifier and ADC are used to generate a binary code proportional to the input analog value, the ROM can be used to generate a corresponding human-oriented engineering unit value for each discrete ADC value. The display is used to present the ROM output in alpha-numerical form to the person desiring to measure the physical quantity.

A technique such as that shown in Figure 4.1 offers many of the qualities desirable in linearization techniques. The technique is very simple in concept and is fairly simple to implement once it is designed. There are several important



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considerations in the design of the system and the programming of the ROM; however, these considerations are not extremely difficult and not altogether different from those required in the design of systems using other linearization techniques. The considerations peculiar to the ROM linearizer are discussed in detail in Chapter 5.

The system is very versatile since the type of amplifier and ADC used may be almost any type as required by the particular application the system is to be used in. The ROM may require different programming for different types of ADCs but the technique in general is virtually ADC-independent. A variety of display devices may also be used as the output code of the ROM may be programmed in many ways although some methods require more bits per word, hence a "larger" ROM which may result in a more costly ROM.

Another advantage of the ROM linearizer is the simplicity of maintenance and calibration. An important characteristic of this system is the absence of feedback loops in the overall block diagram. This eliminates most of the calibration and maintenance problems associated with many of the techniques discussed in Chapter 2. Another important characteristic of the system is that once a ROM has been programmed (truth table generated) many identical ROMs can be manufactured with very little testing time needed. Once an overall system is designed, many can be built requiring only minor calibration of the amplifier and ADC.

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The ROM linearizer is probably the fastest linearizing technique available. Since there are ROMs available with access speeds of one microsecond or less, the response time of the overall system is essentially dependent only on the amplifier, ADC, and display response times.

The use of a ROM linearizer is an extremely reliable technique. One inherent characteristic of read-only memories is that they never forget and have a normal life expectancy of many years. Extremely reliable amplifier, ADC, and display devices are also available so it is possible to construct highly reliable ROM linearizer systems.

The accuracy of a ROM linearizer is dependent on the inherent accuracies of the amplifier and the ADC and the ROM itself. The programming of the ROM is the most important consideration in the overall accuracy of the system. If a ROM is chosen which has the same number of "words" as the ADC has discrete values and each "word" has an adequate number of bits to code the engineering unit value to be represented, then the overall accuracy depends on the number of discrete ADC values used to represent the analog input and the method used to generate the ROM truth table. The considerations involved in determining the accuracy of a ROM linearizer are discussed in detail in Chapter 5.

The economics of the ROM linearizer may range from a very economical system to a very expensive system

depending on a number of factors. The two most important considerations in determining the cost of a ROM are the programming technique used and the size of the ROM. The size of the ROM is usually determined by the range and accuracy desired and the method of coding the output data.

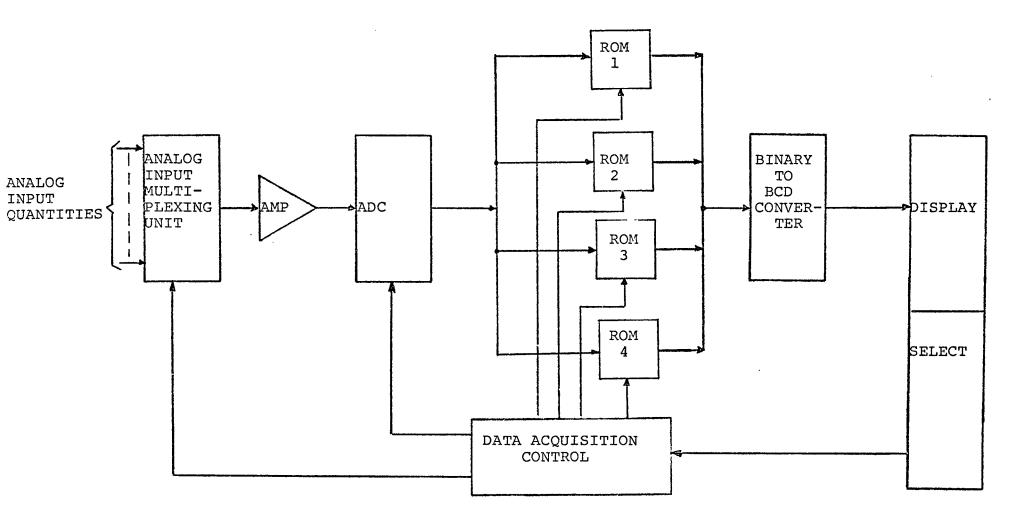
The two most popular methods of coding the output data would probably be binary and BCD. Binary is the most commonly used code in digital applications so the ROM addresses (word numbers) are usually coded in binary. BCD is a very popular code for display purposes since each BCD quad may be decoded to drive a numeric display using the familiar decimal numbers. There are a number of BCD-to-7 segment and BCD-to-decimal decoder/driver/display devices that are very economical, reliable, and simple to use, and which can be adapted to a wide variety of applications.

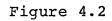
The ROM output data can be coded in binary or BCD format with essentially the same degree of effort. Coding the data in BCD eliminates the necessity of converting the binary data to BCD for display purposes but also requires more bits per word for any numbers greater than 9. For example a 2048 ROM arranged 256 words by 8 bits per word, could be used to generate output data from 0 to 99 if coded in BCD but could generate output data from 0 to 255 if coded in binary.

If it is necessary to generate output displays from 0 to 99, 8 bits would be required for BCD coding but only 7 for binary coding. If the ROM has to have 256 words, then a 2048-bit ROM would be required for BCD output code and a 1792-bit ROM would be required for binary output. If the ROM used costs 5 cents per bit then the cost difference is \$12.80. If it is necessary to convert the binary data to BCD, the cost of the ROM should include the cost of performing the conversion externally if the binary coded ROM is used.

Another consideration in performing the binary-to-BCD conversion externally to the ROM is the fact that doing so introduces more components and more complexity into the system. However, as mentioned previously, ROMs are ideally suited for code conversion applications so the binary-to-BCD conversion could be another highly-reliable easilydesigned ROM. There are also some very economical 4-bit binary-to-BCD and BCD-to-binary converter integrated circuits available which can be configured in arrays to perform n-bit binary-to-BCD and BCD-to-binary conversion.

Figure 4.2 shows a typical application where it might be desirable to use a binary-to-BCD conversion external to the ROM performing the linearization. The system shown has four analog input types each with a different response curve therefore requiring four ROMs for linearizing. If the ROMs





ROM linearizing data acquisiton system



used permit having their outputs tied together, they may all be used to drive a common binary-to-BCD converter and the number of bits per word saved in each linearizing ROM may be more than enough to justify the external binary-to-BCD conversion.

CHAPTER V

GENERATING READ-ONLY MEMORY TRUTH TABLES

The procedure of using read-only memories (ROMs) to convert digitized values of nonlinear analog functions to corresponding engineering units is very similar to the procedure described in Chapter II using a computer program to perform the linearization. The ROM actually performs as a special-purpose computer program for the purpose of linearizing a particular function.

There are two very popular techniques used in the computer program linearization method to determine the physical value to attach to the measured analog quantity. The first procedure involves comparison of voltages by tabular look-up in an appropriate table (such as an NBS thermocouple table)which has been stored in the computer memory. In the second procedure, the table or response curve may be approxmiated by polynomial equations either by a single power series or a set of second order equations.

The single equation power series technique uses an equation of the form:

$$T = C_0 + C_1 E + C_2 E^2 + C_3 E^4 + C_4 E^5 + C_6 E^6$$
 (5.1)

The string of second order equations uses equations of the form:

 $T=C_0+C_1E+C_2E^2$, each equation having a (5.2) a specified range of applicability.

For each measuring device as thermocouples, the accepted standard is the set of tables published in NBS circular 561. Accuracies of linearization techniques for thermocouples are usually measured in terms of degrees of variation from these tables. For thermocouple curve approximation by equations (5.1) and (5.2) the values of the coefficients are determined by curve-fit techniques. Equation (5.1) is usually used when it is desired to use one equation over a broad range of the thermocouple curve. Equation (5.2) is usually used over narrow ranges of thermocouple curves either when the values of interest are confined to a narrow range or when several second order equations are more desirable to evaluate than one high ordered equation. Coefficients for equations (5.1) and (5.2) can be calculated to conform to NBS table values within 0.5°F for most common thermocouple types.

For thermocouple curves the most accurate technique would be to store the NBS look-up tables in a computer program. However, due to the amount of memory storage and time required to store the tables in memory this may not be a very economical approach. Using equations like (5.1) or (5.2) to calculate temperatures from measured emf is a more general approach to nonlinear function measurement. Many nonlinear functions can be easily expressed easily by polynomial equations as in equations (5.1) or (5.2). By merely changing a

few coefficients the same program can be used to evaluate a different nonlinear function. If the table look-up technique is used the entire table would have to be recalculated if a new function is to be evaluated.

Even though the ROM linearizer employs a table lookup technique to perform the linearization, the same arguments apply to calculating the ROM truth table. The truth table for a thermocouple curve can be calculated by hand for a readonly memory but the same procedure would have to be used for every truth table to be generated.

The Fortran program in Appendix A was developed to generate truth tables for ROMs by evaluating a series of equations like equation (5.2). Data required by the program includes the coefficients C_0 , C_1 , and C_2 of equation (5.2), the range for which the equation is valid, and the number of equations used to approximate the curve.

The values for the coefficients and the best ranges for the equations used to describe the function may be determined by any of several commonly used curve-fit techniques. Tables 5.1 and 5.2 give the results obtained by using the simple trial-and-error technique of entering breakpoints into a computer program, permitting the program to generate a second-order equation, then adjusting breakpoints and recalculating until an acceptable small maximum error was reached. The best technique to use may depend on the type

COEFFICIENTS		RANGE OF EOUATION	MAX. DEV. FROM NBS	
с _о	cl	с ₂	(°F)	TABLES (°F)
31.94469 33.91171 43.19788 18.28229 -46.70140 -40.58614 308.71375	34.33704 32.26427 34.47177 39.30495 39.18289	-0.28722 -0.10397 0.00762 -0.04064 -0.13029 -0.13151 0.05607	0-200 200-400 400-700 700-960 960-1220 1220-1400 1400-1600	0.308 0.301 0.289 0.291 0.335 0.297 0.395

Table 5.1

Type J (iron/constantan) thermocouple calibration curve data

COEFFICIENTS		RANGE OF	MAX. DEV.	
с _о	C ₁	C ₂	EQUATION (°F)	FROM NBS TABLES(°F)
31.68816 36.65155 18.50809 45.33263 114.41992 169.81726 272.45447 447.60522	46.15694 42.12250 47.10442 43.60689 37.89775 34.45381 29.44804 21.93095	-0.65195 0.16779 -0.14561 -0.03162 0.08633 0.13980 -0.20083 0.28114	0-140 140-280 280-700 700-1080 1080-1380 1380-1800 1800-2240 2240-2480	0.356 0.292 0.313 0.371 0.347 0.370 0.432 0.349

Table 5.2

Type K (chromel/alumel) thermocouple calibration curve data

of curve to be linearized. Some functions which respond as, for instance, the square root of a stimulus may have a very simple characteristic equation which can be described without requiring use of a curve-fit technique.

The data generated by the ROM Truth Table Generator of Appendix I is a ROM truth table with the ROM address given in decimal and binary format and the data for each ROM word given in decimal and BCD. The decimal representation is provided for easy reference and the binary and BCD codes are provided as the actual bit pattern that would be programmed into the ROM. BCD is chosen for this discussion since it readily lends itself to numeric display in decimal form. The truth table could be easily changed from BCD to binary values by changing just one statement in the main program from a call to a BCD conversion routine to a call to a binary conversion routine.

Besides the calibration curve data, other input data to the ROM Truth Table Generator is the number of ROM words and the ADC resolution (counts per analog input unit). The number of ROM words (NWDS) is used by the program to determine the stopping point of the program. The ADC resolution (NCTS) is used as a scaling factor for the set of equations used to approximate the response curve. For example, the coefficients for Types J and K thermocouples given in Tables 5.1 and 5.2 are based on the units of E for equation (5.2) being in millivolts. The value of NCTS used in the program would represent the number of ADC counts per millivolt of analog input.

The value chosen for NCTS is the key to the accuracy of the ROM table and the eventual size of ROM required for the desired range of the physical stimulus to be measured. The derivation of the choice for NCTS is given in Appendix B along with some examples of how the choice of NCTS affects the inherent system accuracy.

As an analysis example, assume that a Type K thermocouple is to be used in the range of 32°F to 150°F and that it is desired to measure temperatures to the nearest degree in this range. The calibration curve data given in Table 5.2 could be used to calculate the ROM truth table. A Type K thermocouple circuit referenced to 32°F will generate a nonlinear emf from 0 millivolts to 2.66 millivolts for this range.

If the response curve of a Type K thermocouple were linear, only 119 discrete ADC values would be required to display each temperature value between 32°F and 150°F with the amplifier and ADC of Figure 4.1 scaled so 0 ADC counts would represent 32°F and 118 ADC counts would represent 150°F.

However, the thermocouple curve is nonlinear for this range so a ROM is chosen to adjust the ADC values from 0 to 1.18 to "linearize" the response curve. The derivation in

Appendix B reveals that the best choice of the approximate slope S used to determine the value of NCTS for the program of Appendix A is one which approximates the curve as closely as possible with all points of the curve falling below the approximation slope.

Figure 5.1 shows how such a curve might be selected for a Type K thermocouple. The slope of the curve shown in Figure 5.1 would result in an NCTS value of 45 ADC counts per millivolt. The results of using this value for a 512 word ROM are given in Appendix C.

Tables 5.3 through 5.7 show the ROM truth tables for Type K thermocouples that would be obtained for 36-word ROMs using values of 20, 40, 45, 50, and 60 for NCTS. The Type K thermocouple response curve is nonlinear between 32°F and 68°F but the nonlinearity is very slight and the results obtained by using a value of 45 for NCTS shows that the temperature over this small range could be measured within one degree of accuracy without any linearization techniques However, most practical application require temperaat all. ture measurement over a much broader range than 46°F. Sometimes accuracies greater than one degree may be required and many transducer response curves deviate more from a straight line approximation than thermocouples do. Tables 5.3 through 5.7 are presented for purposes of illustrating the importance of a good choice for the value of NCTS.

Using a value too low such as 20 ADC counts per millivolt results in a number of "gaps" in the truth table output data. For example an emf of 0.37 millivolts which should correspond exactly to 49°F would be displayed as 48°F. It is readily obvious that many temperatures could not be displayed since they are not programmed into the ROM.

Using a value too high such as 60 ADC counts per millivolt results in a number of "repeated" values in the truth table output data. While this increases the accuracy of the ROM rather than decreases it, using excessively high values for NCTS may increase the size of the ROM unnecessarily and may increase the cost of the ROM linearizer significantly.

As a comparison, a value of 20 for NCTS permits a temperature range of 32°F to 113°F to be linearized by a 36-word ROM. A value of 45 for NCTS permits a temperature range of 32°F to 68°F to be linearized by a 36-word ROM. A value of 60 for NCTS permits a temperature range of 32°F to 59°F to be linearized by a 36-word ROM. Thus it is readily obvious that the desired inherent accuracy directly affects the required "size" of the ROM and often the cost of the ROM is heavily dependent on the size of the ROM.

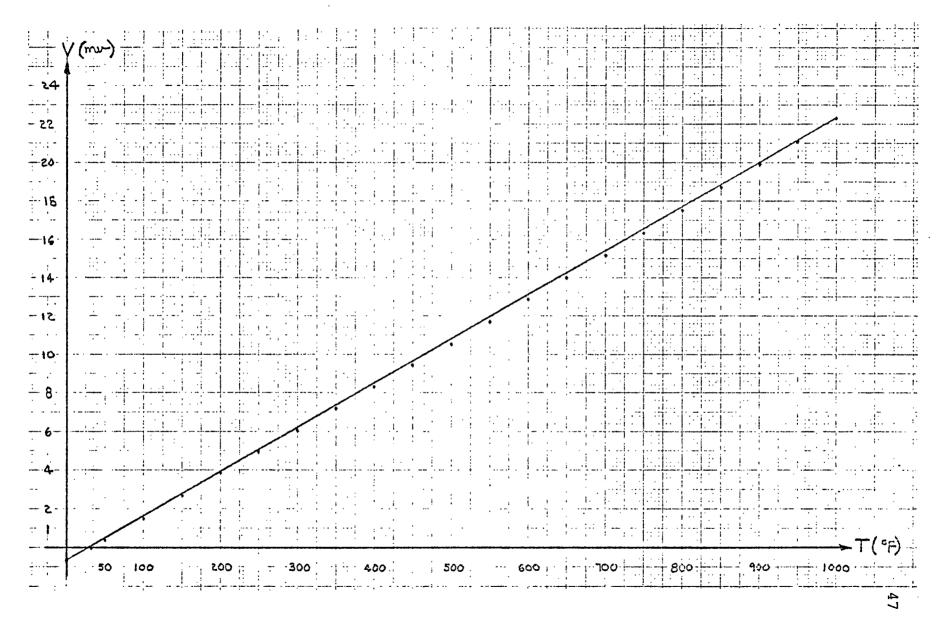


Figure 5.1

Type K (chromel alumel) response curve approximation

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****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
0	32	0000000000	0000000000110010
1	34	0 0000000001	0000000000110100
2	36	8888888818	00000000000110110
3	39	0000000011	00000000000111001
4	41	0000000100	8888888881888881
5	43	0000000101	00000000000000011
6	45	0000000110	0000000001000101
7	48	0000000111	0000000001001009
· 8	50	8666631666	00000 0000101.0000
9	52	0000001001	0000000001010010
10	55	0000001010	0000000001010101
11	57	8888881811	0000000001010111
12	59	0000001100	0000000001011001
13	61	0000001101	0000000001100001
14	64	6666661116	00000000001100100
15	66	0000001111	6666666661166116
16	68	0000010000	0000000001101000
17	70	0000010001	00000000001110000
18	73	0000010010	0 000000001110011
19	75	6666616611	0 000000001110101
20	77	8888616188	00000000001110111
21	79	0000010101	0000000001111001
22	82	0000010110	0000000010000010
23	84	0000010111	0000000010000100
24	86	0000011000	0000000010000110
25	88	0000011001	0000000010001600
26	91	0000011010	0 000000010010001
27	93	0000011011	0000000010010011
28	95	0000011100	0000000010010101
29	97	0000011101	0000000010010111
30	99	0000011110	8080888816811681
31	102	0000011111	0000000100000010
32	104	0000100000	000000010000100
33	106	8888188881	0000000100000110
34	108	8686168616	0000000100001000
35	110	0000100011	6666666166616666

Table 5.3

NCTS = 20

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****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
Ø	32	0000000000	0000000000110010
1	33	00 00000001	00000000000110011
2	34	00 0000010	888888888888118188
3	35	0000000011	0000000000110101
4	36	0000000100	88688888888118118
5	37	0000000101	00000000000110111
6	39	0000000110	0 0000000000111001
7	40	8688666111	8 868888888 1 888888
8	41	0000001000	0000 00000100001
9	42	0 066001601	0 0000000001000010
10	43	0 000001010	000 00000001000011
11	44	0000001011	0000000001000100
12	45	0000001100	8888888888888888
13	47	0000001101	00000000001000111
14	48	0000001110	0 0000000001001000
15	49	8888881111	0 000000001001001
16	50	0000010000	0000000001010000
17	51	0000010001	00 00000001010001
18	52	0 000010010	00000000001010010
19	53	0000010011	0000000001010011
20	55	0000010100	0 000000001010101
21	56	0000010101	0000000001010110
22	57	8898818118	0000000001010111
23	58	8006610111	66666666661611666
24	59	0000011000	0000000001011 091
25	60	00 00011001	666666666611866666
26	61	0000011010	88888888881188881
27	63	0000011011	00000000001100011
28	64	0000011100	0000000001100100
29	65	0000011101	66666666661166161
30	66	0000011110	00000000001100110
31	67	0000011111	6666666661166111
32	68	0090100000	666666661161666
33	69	0000100001	0000000001101001
34	70	6000100010	66666666611166666
35	72	0000100011	00000000001110010

Table 5.4

NCTS = 40

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*****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
Ø	32	8888888888	00000000000110010
1	33	0000000001	88888888888118811
2	34	0000000010	88888888888118188
3	35	0000000011	00000000000110101
4	36	0000000100	0000000000110110
5	37	0000000101	0000000000110111
6	38	6666666116	0600000000111000
7	39	0000000111	0000000000111001
8	40	0000001000	0000000001000000
9	41	0000001001	0000000001000001
10	42	0000001010	88888888881888818
11	43	8888881811	888888888888888888
12	44	0000001100	ØØØØØØØØØ1ØØØ1ØØ
13	45	8888881181	0000000001000101
14	46	8666661118	0000000001000110
15	47	8888881111	0000000001000111
16	48	6666616666	0000000001001000
17	49	0000010001	00000000001001001
18	50	0000010010	6666666661616666
19	51	0000010011	6666666661616661
20	52	8888618168	0000000001010310
21	53	0000010101	6666666661616611
22	54	0000010110	0000000001010100
23	55	8088818111	00000000001010101
24	56	8888811888	0000000001010110
25	57	0000011001	00000000001010111
26	58	8888811818	0000000001011000
27	59	0000011011	0000000001011001
28	60	0000011100	0000000001100000
29	61	0000011101	6666666661166661
30	62	0000011110	0000000001100 010
31	63	8888811111	88888888881188811
32	64	0000100000	0000000001100100
33	65	0000100001	6666666661166161
34	66	0000100010	8666666661166118
35	67	0000100011	0000000001100111

Table 5.5

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NCTS = 45

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****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
Ø	32	66666666666	0000000000110010
1	33	8888888881	0000000000110011
2	34	0000000910	0000000000110100
3	34	0000000011	0000000000110100
4	35	0000000100	0000000000110101
. 5	36	0000000101	0000000000110110
6	37	8868888118	0000000000110111
7	38	0000000111	0000000000111000
8	39	6666661666	6666666666111661
9	40	0000001001	66666666616666666
10	41	8666691816	0000000001000001
11	42	0000001011	8666666661666616
12	43	0000001100	00000000001000011
13	44	0000001101	00 000000001000100
14	45	0000001119	0 0000000001000101
15	45	0000001111	88888888881888181
16	46	0000010000	00000000001000110
17	47	0000010001	0 000000001000111
18	48	0000010010	0000000001001000
19	49	0000010011	0000000001001001
20	50	0000010100	0 000000091010000
21	51	8888819191.	8888888881818881
22	52	0000010110	0000000001010010
23	53	0000010111	6666666661616611
24	54	8888811888	68888888881818188
25	55	0000011001	6666666661616161
26	56	0000011010	88888888881818118
27	56	0000011011	0000000001010110
28	57	8888811188	0000000001010111
29	58	0000011101	0000000001011000
30	59	0000011110	0000000001011001
31	60	6666611111	0000000001100000
32	61	0000100000	0000 000001100001
33	62	0000100001	0000000001100010
34	63	0000100010	0000000001100011
35	64	0000100011	86868888881188188

Table 5.6

NCTS = 50

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****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
Ø	32	00000000000	0000000000110010
1	32	. 8888888881	0000000000110010
2	33	0000000010	8688888888118811
3	34	00000000 11	00000000000110100
4	35	0000000100	0000000000110101
5	36	0000000101	0000000000110110
6	36	0000000110	00000000000110110
7	37	0000000111	0000000000110111
8	38	0000031008	0000000000111000
9	39	6666661661	00000000000111001
10	39	0000001010	0000000000111001
11	40	0000001011	88888888881888888
12	41	0000001100	0000000001000001
13	42	0000001101	0000000001000010
14	42	0000001110	0000000001000010
15	43	0000001111	0000000001000011
16	44	0000010000	0000000001000100
17	45	0000010001	00000000001000101
18	45	0000010010	88688888881688181
19	46	0000010011	0000000001000110
20	47	6868616168	0000000001000111
21	48	0000010101	06666333661391666
22	49	0000010110	6666666661661661
23	49	8888818111	0000000001001001
24	50	0000011000	00000000001010000
25	51	0000011001	0000000001010001
26	52	0000011010	0000000001010010
27	52	0000011011	8888888881818818
28	53	0000011100	0 0000000001010011
29	54	0000011101	0020006001010100
30	55	0000011110	0000000001010101
31	55	0000011111	0000000001010101
32	56	0000100000	0000000001010110
33	57	0000100001	0 000000001010111
34	58	0000100010	0000000001011000
35	58	0000100011	0000000001011000

Table 5.7

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NCTS = 60

CHAPTER VI

RESULTS AND CONCLUSIONS

Using the formulas derived in Appendix B and the results discussed in Chapter V it can be determined readily that to "linearize" a Type K thermocouple curve from 32°F to 500°F requires a 512 word by 11 bits per word if the result is to be in BCD to the nearest degree. A truth table for such a ROM is given in Appendix C. The required 512 x 11 ROM is a 5632-bit ROM. If the output data of the ROM were in binary and converted to BCD externally to the linearizing ROM, the required ROM would be 512 words by 9 bits per word or a 4608bit ROM.

It is relatively simple to implement either a 512 x 11 or a 512 x 9 ROM using one or more presently available semiconductor ROM integrated circuits. Some of the more popular types available are 1024-bit ROMs arranged 128 x 8; 2048-bit ROMs arranged 256 x 8 or 512 x 4; 4096-bit ROMs arranged 512 x 8 or 1024 x 4; and 8192-bit ROMs arranged 1024 x 8.

The present cost of such semiconductor ROMs is between 1¢ and 5¢ per bit in quantities of 100 or more while unit quantities may range from 3¢ to 10¢ per bit.

Measurement accuracies within 1°F for Type K thermocouples can be obtained with one straight-line approximation segment between 0°F and 500°F with such conventional techniques as the ramp ladder network. This same technique can

attain the same accuracy from 0°F to 2000°F for Type K thermocouples with 6 straight-line approximation segments.

The ROM linearization technique presented here requires 512 words x 11 bits (BCD output) or 512 words x 9 bits (binary output) for the same degree of accuracy between 0°F and 500°F and a 2048 x 14 ROM (BCD output) or 2048 x 11 ROM (binary output) for the same degree of accuracy between 0°F and 2000°F.

In most industrial process applications accuracies within one or two degrees Farenheit of the actual value are acceptable. In such cases it is probably better in most applications to use a conventional linearization technique, such as one of those discussed in Chapter II, rather than the ROM linearization technique presented here. This is particularly true since the largest single-chip ROM presently available is 8192 bits and the cost of ROMs is still high enough that conventional linearization techniques are more economical for wide-range, low degree of accuracy applications.

Semiconductor industry predictions estimate that by 1975 single chip memories of 16,000 bits and greater will be available between 0.1¢ and 1¢ per bit. If these predictions prove to be correct the ROM linearization technique presented here may become economically attractive for some applications between now and 1975.

Presently for most industrial process applications the advantages of the ROM linearization technique are probably outweighed by the cost of implementing such a system. While this may be true in particular for thermocouples and other "almost linear" response curves, there may be some applications that more readily lend themselves to the characteristics of the ROM linearizer.

Several articles have been written about techniques using ROMs to synthesize complex waveforms. Such devices are very similar to implementing the ROM linearizer in reverse by using an up/down counter to sequentially address the inputs of a ROM programmed with a complex function look-up table and using the ROM output to drive a digital-to-analog converter.

The derivation given in Appendix B shows how a ROM can easily be programmed to analyze complex equations. The program developed here for generating ROM truth tables is based on describing the response function by one or more quadratic equations. The program could easily be changed to evaluate almost any type of simple or complex characteristic equation of complex function to generate truth tables.

If it is desired to analyze highly complex functions it may be more desirable to use a ROM linearizer rather than a conventional technique just as it has already been shown that is is sometimes more desirable to use ROMs to synthesize the functions.

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APPENDIX A

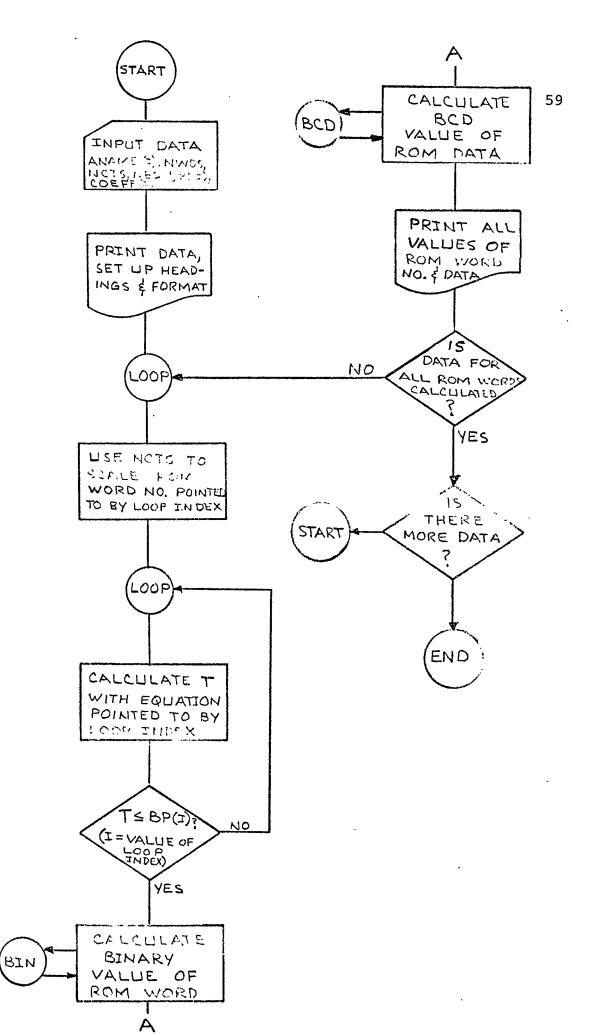
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FORTRAN COMPUTER PROGRAM FOR ROM TRUTH TABLE GENERATOR



С С ROM TRUTH TABLE GENERATOR 60 C С С ANAME DESCRIBES FUNCTION TO LINEARIZE ¢ NWDS IS NUMBER OF ROM WORDS C NCTS SPECIFIES SCALING OF ADC С NEAN IS NUMBER OF EQUATIONS REPRESENTING CURVE С COEFF 1RE THE A, B, C COEFFICIENTS OF THE EQUATIONS C T IS THE CALCULATED VALUE OF THE EQUATIONS С DIMENSION NBIN(10), NBCD(16), BP(10), COEFF(30), ANAME(20) 1 READ (3,105) (ANAME(I), I=1,20) WRITE (4,105) ANAME READ(3, 100) NWDS, NCTS, NEQN WRITE (4, 107) NEQN DO 2 I=1, NEQN JJ≈3*I-2 JK=3*I READ(3,102)(COEFF(J), J=JJ, JK) WRITE(4,108) (COEFF(J), J=JJ, JK) 2 CONTINUE WRITE (4,109) READ(3,101) (BP(I), I=1,10) WRITE(4,101) (EP(I), I=1,10) WRITE (4,110)NCTS WRITE (4, 111) NWDS CALL FORM WRITE (4, 105) ANAME WRITE (4,103) WRITE (4,106) DO 5 I=1, NWDS N=1-1 X=FLORT(N)/FLORT(NCTS) DO 3 J=1,10 JBP=J*3 T=COEFF(JBP-2)*X**2+COEFF(JBP-1)*X+COEFF(JBP) IF(EP(J)-T)3,4,4 3 CONTINUE 4 CONTINUE NT=IFIX(T+. 5) CALL BIN(N, NBIN) CALL BCD(NT, NBCD) WRITE(4, 104) N, NT, (NBIN(K), K=1, 10), (NBCD(K), K=1, 16) 5 CONTINUE 1 GO TO 1 100 FORMAT (315) FORMAT(10F8.1) 101 102 FORMAT(3F20.10) FORMAT(2228X) 38H****READ ONLY MEMORY TRUTH TRELE*****/2) 103 FORMAT (8X, 14, 8X, 14, 8X, 1011, 10X, 1611) 104 105 FORMAT (20A4) FORMAT(7%, SHROM WORD, 4%, SHROM DATA, 4%, 13HROM WORD(BIN), 8%, 106 1 13HROM DATA(BCD)/)

107 FORMAT(///3X, 14, 1X, 9HEQUATIONS//)

108 FORMAT(3X, F14, 8, 6HX**2 +, F14, 8, 3HX +, F14, 8)

109 FORMAT(///14HBREAKPOINTS AT/)

110 FORMAT(///3X, 14, 26H ADC COUNTS PER INPUT UNIT)

111 FORMAT(///5X,13HSIZE OF ROM =,14,6H WORDS) END

END

0,*SIZE 001323

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1,*SIZE 000104

С	
С	INTEGER TO 10-BIT BINARY CONVERTER
С	
C	
С	NN IS NUMBER TO CONVERT
С	NBIN IS BINARY VALUE OF NUMBER
С	
	SUBROUTINE BIN(NN, NBIN)
	DIMENSION NBIN(10)
	KR=NN
	DO 40 I=1,10
	IJ=11-I
	K=KR/2
	KN=K*2
	NBIN(IJ)=KR-KN
	KR=K
40	CONTINUE
	RETURN
	END
	END
0 46170	E 000076
1) #0120	E 000024

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с с с	INTEGER TO 16-BIT BCD (4 DIGIT) CONVERTER 63
. C C C	NTN IS NUMBER TO CONVERT NBCD IS BCD VALUE OF NUMBER IN BIT FORMAT
	SUBROUTINE BCD(NTN,NBCD) DIMENSION NBCD(16) NTTH=NTN/1000
	N1=NTTH DO 51 I=1,4 IJ=5-I
	K=N1/2 K1=K*2
	NBCD(IJ)=N1-K1 N1=K
51	CONTINUE NTTH1=NTTH*1000 NTTH2=NTN-NTTH1
	NTH=NTTH2/100 N1=NTH
	DO 52 I=1,4 IJ=9-I K=N1/2
	K1=K*2 NBCD(IJ)=N1-K1
52	N1=K CONTINUE NTH1=NTH*100
	NTH2=NTH2-NTH1 NTT=NTH2/10
	N1=NTT D0 53 I=1,4
	IJ=13-I K=N1/2 K1=K*2
	NBCD(IJ)=N1-K1 N1=K
53	CONTINUE NTT1=NTT*10 NTO=NTH2-NTT1
	N1=NTO D0 54 I=1,4
	1J=17-I K=N1/2
	K1=K*2 NBCD(IJ)=N1-K1 N1=K
54	CONTINUE RETURN END
0,*SI	END

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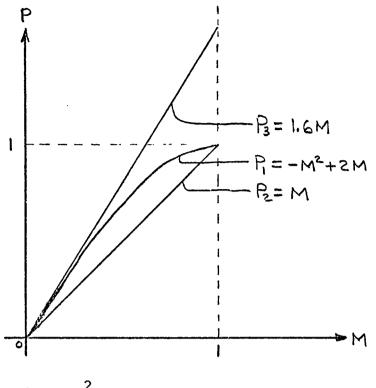
DERIVATION OF ADC RESOLUTION AND ROM LINEARIZER ACCURACY

APPENDIX B

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 $P = -M^2 + 2M$

P=range of physical stimulus M=range of measured quantity S=approximate slope of curve R = ADC resolution V = value of ADC counts T = total ADC counts required N = number of ADC bits (physical units-°F,°C,GPM, etc) (measured units-mw,ma, etc) (slope of P₂ or P₃) (ADC counts per physical unit) (ADC counts per measured unit) (ADC counts, ROM words) (binary coding) Assume that P is the range of the physical stimulus to be measured and some device responds to this stimulus generating a measurable quantity M and that the characteristic equation for this response from 0 to 1 units of P is

$$P = -M^2 + 2M$$

Assume also for analysis purposes that the units of M are in millivolts and the units of P are in gallons per minute (GPM).

The appropriate slope of this curve is given by

$$S = P/M$$

where P is a full scale value of P and M is a full scale value of M. The figure given shows two approximation for S, $S_1 = 1$ and $S_2 = 1.6$. Both approximations are highly inaccurate but one shows all true values of P above the approximation and the other shows all values of P below the approximation. Any value of S between 1 and 1.6 would yield an approximate curve with some real values of P above and some values below the approximation line.

2. Let R represent the resolution of the ADC in counts per physical unit. In this case the resolution would be expressed in counts per GPM. Assume that it is desired to display values of GPM to the nearest tenth between 0 and 1, then the value of R would be 10 counts per GPM.

3. Let V represent the measurable value of each ADC count.

This is given by

$$V = RS$$

which has the units of counts per measured value or in this case counts per millivolt.

Thus it follows that

 $V_1 = S_1 R = (1)(10) = 10$ counts per mv $V_2 = S_2 R = (1.6)(10) = 16$ counts per mv

4. Let T represent the number of ADC counts required to represent the maximum desired value of P. This is given by

$$T = VM$$

or in this case the value of T is

$$T_1 = V_1 M = 10$$
 counts
 $T_2 = V_2 M = 16$ counts

This points out the fact that as the slope of the approximation curvegets steeper more ADC counts, and hence, the more ROM words will be required to generate the displayed values. But also the accuracy of the ROM linearizer is more accurate with a 16 count ADC than with 10.

Let N = the number of bits representing a full-scale
ADC word. This is given by

$$N = \log_2 T$$

For this example

$$N_1 = \log_2 T_1 = \log_2 10 = 3.322$$

 $N_2 = \log_2 T_2 = \log_2 16 = 4.0$

. Hence, to fully represent the desired values of P, using slope S₁ would require a 4-bit ADC and a 10 word ROM and using slope S₂ would require a 4-bit ADC and a 16 word ROM.

6. The results of using a 10-word and a 16-word ROM to "linearize" the curve of this example are given in Tables B.1 and B.2. Note that in Table B.1 the approximation does not permit displays to the nearest tenth of the full scale range but as shown in Table B.2, using a 16-word ROM does permit displays to the nearest tenth.

Figures B.1 and B.2 illustrate the relative accuracies of using 10 and 16-word ROMs for this example. Note that a 10-word ROM results in a maximum error of about 0.12 while a 16-word ROM results in a maximum error of 0.08. Thus the 16-word ROM with a 4-bit ADC satisfies the design criteria for this example.

The accuracy of the ROM linearizer can be improved by using a larger ADC to get more counts per input unit and a larger ROM with more words and perhaps also with another BCD quad per word to display the nearest hundredth of the linearized data.

TYPICAL QUADRATIC RESPONSE CURVE

****READ ONLY MEMORY TRUTH TABLE*****

ROM	NORD	ROM	DATA	ROM WORD(BIN)	ROM DATA(BCD)
	Ø		Ø	0000000000	000000000000000000
	1		2	0 000000001	000000000000000000000000000000000000000
	2		4	0000000010	00000000000000100
	3		5	0000000011	00000000000000101
	4		6	0000000100	00000000000000110
	5		8	0000000101	88 888888888888881888
	6		8	0000000110	000000000000001000
	7		9	0000000111	00000000000001001
	8	1	.0	0000001000	0 000000000010000
	9	1	.0	0000001001	8666666666616666

Table B.1 NCTS = 10

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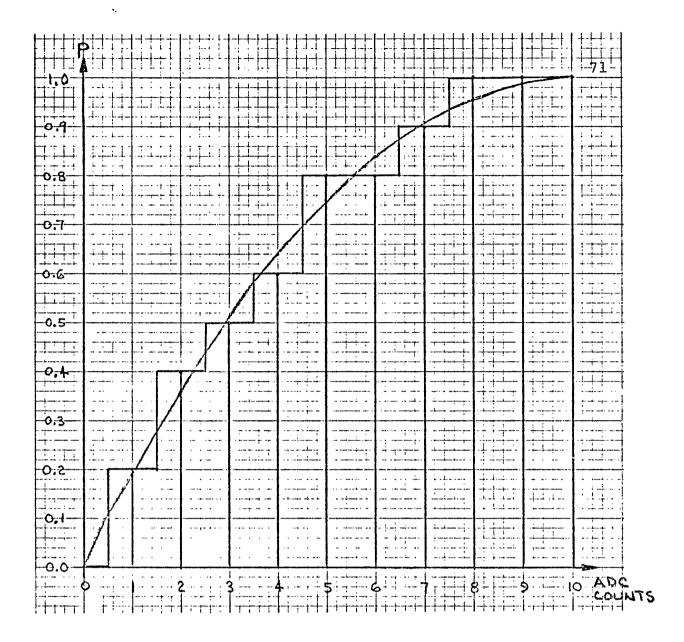
TYPICAL QUADRATIC RESPONSE CURVE

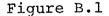
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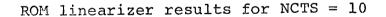
****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
Ø	Ø	8008089089	00000 0000000000000000000000000000000
1	1	0000000001	00 0000000000000000001
2	2	0000000010	00 000000000000010
З	3	0000000011	00 0000000000000011
4	4	0000000100	0000000000000000 00000000000000000000
5	5	0000000101	00000000000000101
6	6	0000000110	8080808080808 08118
7	7	0000000111	00 000000000000111
8	8	0000001000	88 888888888881888
9	8	0000001001	00 00000000001000
10	9	0000001010	00 000000000001001
11	9	0000001011	86888888888888
12	9	0000001100	00000000000001001
13	10	0000001101	88 866666666616666
14	1.0	6666661116	00 00000000016000
15	10	0000001111	0 000000000010000

Table B.2 NCTS = 16







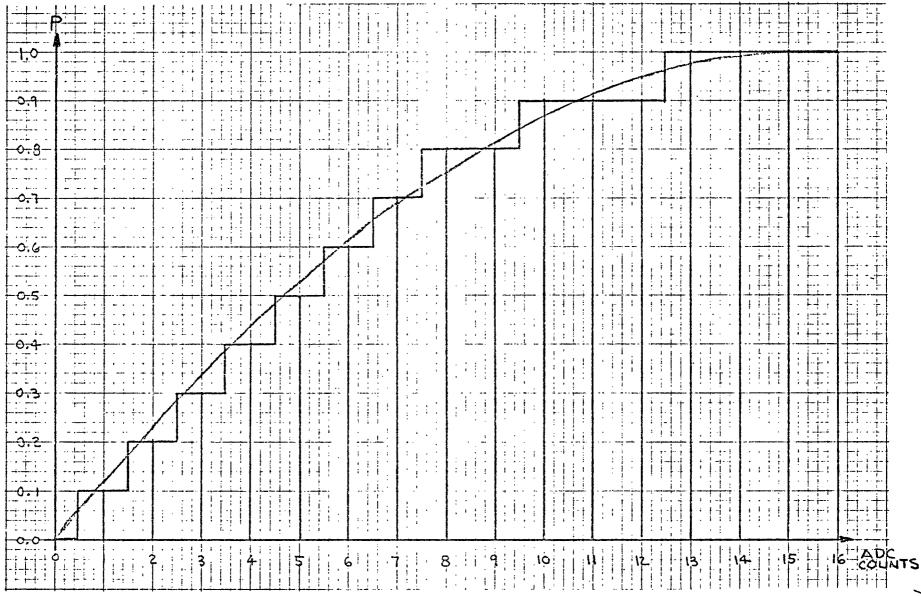


Figure B.2

ROM linearizer results for NCTS = 16

APPENDIX C

.

TYPE K THERMOCOUPLE LINEARIZATION RESULTS

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TYPE K THERMOCOUPLE

8 EQUATIONS

65194988X**2	+	46. 15693665X	+	31. 68815613
. 16779000X**2	+	42. 12249756X	+	36. 65153503
14560997X**2	+	47. 10441589X	+	18. 50808716
03162000×**2	+	43. 60688782X	+	45. 33262634
.08633000×**2	+	37.89773560X	+	114. 41989136
. 13979995×**2	+	34. 45379639X	+	169. 81719971
. 20082998×**2	+	29. 44803619X	+	272. 45446777
28113997X**2	+	21. 93094635X	+	447.60510254

BREAKPOINTS AT

140.0 280.0 700.0 1080.0 1380.0 1800.0 2240.0 2480.0 0.0

45 ADC COUNTS PER INPUT UNIT

SIZE OF ROM = 512 WORDS

*****READ ONLY MEMORY TRUTH TABLE*****

ROM WORD	ROM DATA	ROM WORD(BIN)	ROM DATA(BCD)
ø	32	88888888888	6666666666116616
1	33	00000000001	0000000000110011
2	34	6666666618	666666666666118166
3	35	0000000011	0000000000110101
4	36	0000000100	0000000000110110
5	37	8888888181	8888888888118111
6	38	0000000110	8888888888111888
7	39	0000000111	00000000000111001
8	40	0000001000	8888888881888888
9	41	0000001001	00000000001000001
10	42	0000001010	66666666661668616
11	43	0000001011	6666666661666611
12	44	0000001100	666666666666661666166
13	45	6666661161	0000000001000101
14	46	0000001110	00000000001000110
15	47	6666661111	8686666661866111
16	48	0000010000	8000000001001000
17	49	0000010001	8888888881881881
18	50	0000010010	0000000001010000
19	51	0000010011	00000000010100 01
20	52	0000010100	6666666661616616
21	53	0000010101	00000000001010011
22	54	0000010110	666666666666666666
23	55	0000010111	00000000001010101
24	56	0000011000	0000000001010110
25	57	0000011001	0000000001010111
26	58	0000011010	0000000001011000
27	59	0000011011	0000000001011001
28	60	0000011100	00000000001100000
29	61	9999911191	0000000001100001
30	62	0000011110	000000001100019
31	63	0000011111	99999999911999 11 99999999911999 11
32	64	0000100000	00000000001100100 00000000001100101
33	65 62	0000100001	0000000001100101 0000000001100110
34 35	66 67	0000100010 0000100011	000000000000000000000000000000000000000
30	68 67	0000100011 0000100100	0000000001100111 0000000001100111
36 37	69	0000100100 0000100101	000000000000000000000000000000000000000
38 38	69 70	0000100101 0000100110	000000000000000000000000000000000000000
39	70 71	0000100110	00000000001110000 0000000001110000
37 40	72	0000100111	0000000001110001 0000000001110010
70	14	O O O O T O T CARDO	SOCCESSION TITCOTS

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.1.1	77	00004 04 004	
41	73	0000101001	868688881116811
42	74	0000101010	00000000001110100
43	75	0000101011	0000000001110101
44	76	0000101100	00000000001110110
45	77	0000101101	8888888881118111
46	78	0000101110	0000000001111000
47	79	0000101111	0000000001111001
48	80	0000110000	00000000100000000
49	81	0000110001	6666666616666666
50	82	0000110010	6666666616686618
51	83	0000110011	0000000010000011
52	84	0000110100	000000010000100
53	85	0000110101	6666666616666181
54	86	0000110110	6666666616668116
55	87	0000110111	6666666616666111
56	88	0000111000	0000000010001000
57	89	0000111001	0000000010001001
58	98	0006111010	0000000010010000
59	91	0000111011	8666666616616861
60	92	0000111100	0000000010010010
61	93	0000111101	0666666616916911
62	94	0000111110	0000000010010100
63	95	0000111111	0000000010010101
64	96	0001000000	0000000010010110
65	97	0001000001	66666666166161 11
66	98	8861888618	0000000010011000
67	99	0001000011	0000000010011001
68	100	8661666166	86666661666666666
69	101	0001000101	6666666166666666
70	102	0001000110	86686661866666618
71	103	8881688111	0 00000010000001 1
72	104	0001001000	0000000100000100
73	105	<u>8681681681</u>	0000000100000101
74	106	<i>0</i> 001001010	0000000100000110
75	107	0001001011	000000100000111
76	108	0001001100	6666666166661666
77	109	0001001101	666666616661661
78	110	0001001110	0000000100010000
79	111	0001001111	00000010001000 1
80	112	0001010000	0000000100010010
81	113	0001010001	0000000100010011
82	114	0001010010	0000000100010100
83	115	0001010011	8686666168616181
84	116	0001010100	000000100010110
85	117	0001010101	8888888188518111
			•

86	118	0001010110	000000100011000
87	118	00010101 11	0000000100011000
88	119	0001011000	0000000100011001
89	120	8881811881	0000000100100000
90	121	0001011010	0000000100100001
91	122	0001011011	0000000100100010
92	123	0001011100	0000000100100011
93	124	8881811181	0000000100100100
94	125	0001011110	0000000100100101
95	126	0001011111	0000000100100110
96	127	0001100000	0000000100100111
97	128	0001100001	0000000100101000
98	129	8661166618	0000000100101001
99	130	0001100011	0000000100110000
100	131	0001100100	0000000100110001
101	132	8661186161	0000000100110010
102	133	0001100110	0000000100110011
103	134	0001100111	0000000100110100
104	135	0001101000	0000000100110101
105	136	<u>8881181881</u>	6666666166116116
106	137	0001101010	0000000100110111
107	138	0001101011	0000000100111000
108	139	0001101100	0000000100111001
109	140	0001101101	00000001010000000
110	141	0001101110	0 000000101000001
111	142	0001101111	000000101000010
112	143	0001110000	0000000101000011
113	143	0001110001	0000000101000011
114	144	0001110010	0000000101000109
115	145	0001110011	0000000101000101
116	146	0001110100	0000000101000110
117	147	0001110101	0000000101000111
118	148	0001110110	888888899181881898
119	149	0001110111	6666666161661661
120	150	0001111000	0000000101010000
121	151	0001111001	8888888181818881
122	. 152	0001111010	0000000101010010
123	153	0001111011	0000000101010011
124	154	0001111100	0000000101010100
125	155	0001111101	6666666161616161
126	156	0001111110	8666666161616116
127	157	0001111111	8888888191818111
128	158	0010000000	000000101011000
129	159	0010000001	0000000101011001
130	160	0010000010	0000000101100000

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131	161	0010000011	6666666161166661
132	162	0010000100	6666666161166616
133	163	0010000101	0000000101100011
134	164	0010000110	0000000101100100
135	165	0010000111	0000000101100101
136	165	0010001000	0000000101100101
137	166	8818881881	0000000101100110
138	167	0010001010	88888888181188111
139	168	0010001011	00000000101101000
140	169	0010001100	0000000101101001
141	170	8818881181	88888888181118888
142	171	0010001110	0000000101110001
143	172	9010001111	0000000101110010
144	173	0010010000	0000000101110011
145	174	0010010001	0000000101110100
146	175	0010010010	0000000101110101
147	176	0010010011	0000000101110110
148	177	0010010100	0000000101110111
149	178	0010010101	0000000101111000
150	179	0010010110	88888888181111881
151	180	0010010111	0000000110000000
152	181	0010011000	0000000110000001
153	182	6618611861	6666666116666616
154	183	6616611616	8888888118888811
155	184	0010011011	8868666118866188
156	185	0010011100	0000000110000101
157	186	0010011101	8668868116886118
158	187	6616611116	0000000110000111
159	188	0010011111	0000000110001000
160	189	0010100000	00000000110001001
161	190	0010100001	0000000110010000
162	190	0010100010	0000000110010000
163	191	0010100011	66666666116616661
164	192	0010100100	86666666116616616
165	193	0010100101	66666669116616611
166	194	0010100110	0000000110010100
167	195	6616166111	8868888116818181
168	196	881818188	0000000110010110
169	197	0010101001	0000000110010111
170	198	9919191919	8888888118811888
171	199	0010101011	0000000110011001
172	200	8616161166	6666666166666666666
173	200	0010101100	8086661866666666
174	201	8616161116	00000010000000010
175	202 203	9919191118 9919191111	00000010000000010
IT U	202	مله علم من الله الله الله الله الله الله الله الل	

176	204	0010110000	0000001000000100
177	205	0010110001	0000001000000101
178	206	0010110010	0000001000000110
179	207	8818118811	0000001000000111
180	208	0010110100	000000100001000
181	209	0010110101	66666616666661661
182	210	0010110119	8886661666616668
183	211	0010110111	0000001000010001
184	212	0010111000	0060001000010010
185	213	0010111001	0000001000010011
186	214	0010111010	0000001000010100
187	215	0010111011	0000001000010101
188	216	0010111100	0000001000010110
189	217	0010111101	0000001000010111
190	217	0010111110	6666661666616111
191	218	<i>001011111</i>	0000001000011000
192	219	0011000000	0000001000011001
193	220	8011888881	0000001000100000
194	221	0011000010	888888188818881
195	222	6611666611	0000001000100010
196	223	0011000100	8888881888188811
197	224	0011000101	0000001000100100
198	225	6611666116	0000001000100101
199	226	0011000111	0000001000100110
200	227	0011001000	6699661999166111
201	228	0011001001	6000001000101000
202	229	0011001010	0000001000101001
203	230	0011001011	0000001000110000
204	231	0011001100	6896691688118661
205	232	0011001101	0000001000110010
206	233	0011001110	0000001000110011
207	234	0011001111	6666661666116166
208	235	881181888	0000091000110101
209	236	0811818881	8888881898118118
210	237	0011010010	8000001000110111
211	238	0011010011	8888881888111888
212	239	6611616166	88686691666111661
213	240	0011010101	00000010010000000
214	241	6611616118	0000001001000001
215	242	0011010111	6000001661666916
216	243	. 0011011000	8666881981888811
217	244	0011011001	6666661661666166
218	245	0011011010	0000001001000101
219	246	0011011011	8666691861869118
220	247	0011011100	6666661661666111

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221	248	8811811181	0000001001001000
222	249	0011011110	000001001001001
223	250	0011011111	6666661661616668
224	250	0011100000	0000001001010000
225	251	0011100001	0000001001010001
226	252	0011100010	0000001001010010
227	253	0011100011	0000001001010011
228	254	0011100100	0000001001010100
229	255	0011100101	0000001001010101
230	256	0011100110	0000001001010110
231	257	0011100111	0000001001010111
232	258	0011101000	0000001001011000
233	259	0011101001	6666661661611661
234	260	0011101010	0000001001100000
235	261	0011101011	0000001001100001
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238	264	0011101110	0000001001100100
239	265	0011101111	6666661661166161
240	266	0011110000	6666691661166116
241	267	0011110001	0000001001100111
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246	272	0011110110	0000881081110810
247	273	0011110111	0000001001110011
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249	275	6611111661	0000001001110101
250	276	0011111010	0000001001110110
251	277	0011111011	0000001001110111
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253	279	001111101	0000001001111001
254	280	001111110	888888181888888
255	281	0011111111	00000010100000001
256	282	0100000000	0000001010000010
257	283	0100000001	0000001010000011
258	284	0100000010	0000001010000100
259	285	8188889911	0000001010000101
260	286	0100000100	0000001010000110
261	287	0100000101	0000001010000111
262	288	0100000110	8888881016861688
263	289	0100000111	000001010001001
264	290	0100001000	0000001010010000
265	291	0100001001	0000001010010001

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266	292	0100001010	0000001010010010
267	293	0100001011	0000001010010011
268	294	0100001100	0000001010010100
269	295	0100001101	0000001010010101
270	296	0100001110	0000001010010110
271	297	0100001111	0000001010010111
272	298	0100010000	0000001010011000
273	299	0100010001	0000001010011001
274	300	0100010010	00000011000000000
275	301	0100010011	00000011000000001
276	302	0100010100	0000001100000010
277	303	0100010101	0000001100000011
278	304	8188818118	0000001100000100
279	305	0100010111	8888881188888181
280	306	0100011000	0000001100000110
281	307	0100011001	0000001100000111
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283	309	0100011011	8888881188881881
284	310	0100011100	0000001100010000
285	311	6166611161	8888881188818881
286	312	0100011110	0000001100010010
287	313	0100011111	8886681186818811
288	314	0100100000	0000001100010100
289	315	0100100001	0000001100010101
290	316	0100100010	0886891168818110
291	317	0100100011	0000001100010111
292	318	0100100100	8868661188811888
293	319	0100100101	8988891188811881
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296	322	0100101000	0000001100100010
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306	332	0100110010	8888881188118818
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310	336	0100110110	0000001100110110

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311	337	0100110111	0000001100110111
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313	339	0100111001	0000001100111001
314	340	0100111010	0000001101000000
315	341	0100111011	0000001101000001
316	342	0100111100	0000001101000010
317	343	0100111101	0000001101000011
318	344	0100111110	0000001101000100
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321	347	0101000001	0000001101000111
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323	349	6161666611	0000001101001001
324	350	0101000100	0000001101010000
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326	352	0101000110	00000011010100010
327	353	0101000111	0000001101010011
328	354	0101001000	888888118181818
329	355	0101001001	8888881181818181
330	356	0101001010	0000001101010110
331	357	0101001011	0000001101010111
332	358	0101001100	8888881181811888
333	359	0101001101	0000001101011001
334	360	0101001110	6868381191166666
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337	363	0101010001	6666661161166611
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339	365	0101010011	0000001101100101
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347	373	0101011011	6666661161116611
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353	379	0101100001	0000001101111001
354	380	0101100010	0000001110000000
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357	383	6161166161	8866661116666611
358	384	8181188118	0000001110000100
359	385	0101100111	8888881118888181
360	386	0101101000	0000001110000110
361	387	0101101001	8888881118888111
362	388	0101101010	0000001110001000
363	389	0101101011	0000001110001001
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366	392	0101101110	0000001110010010
367	393	0101101111	0000001110010011
368	394	0101110000	0000001110010100
369	395	0101110001	0000001110010101
370	396	0101110010	8686881118818110
371	397	0101110011	0000001110010111
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374	400	0101110110	888881868888888888
375	401	8181118111	6666616666666666
376	402	0101111000	8888818888888888
377	403	8181111881	8668618888668611
378	404	0101111010	0000010000000100
379	405	0101111011	00000100000000101
380	406	0101111100	0000010000000110
381	407	0101111101	0000010000000111
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383	409	0101111111	0000010000001001
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386	412	0110000010	88888198881988
387	413	0110000011	0000010000010011
388	414	0110000100	88888188881818188
389	415	0110000101	6666616666616161
390	416	0110000110	8888818888818118
391	417	0110000111	8868818888818111
392	418	0110001000	6666616666611666
393	419	0110001001	888881888811881
394	420	0110001010	888881888818888
395	421	0110001011	000001000010000 1
396	422	0110001100	000001000010001 0
397	423	0119991191	0000010000100011
398	424	0110001110	6666616666166166
399	425	0110001111	0000010000100101
400	426	0110010000	0000010000100110

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401	427	8118818881	6666616666166111
492	428	0110010010	0000010000101000
403	429	0110010011	8888818888181881
404	430	0110010100	0000010000110000
405	431	0110010101	8888818688118881
406	432	0110010110	8666916669116616
407	433	0110010111	6666616666116611
408	434	0110011000	8888818688118188
409	435	0110011001	8868818889118181
410	436	0110011010	0000010000110110
411	437	0110011011	8966916666116111
412	438	0110011100	0000010000111000
413	439	0110011101	0000010000111001
414	440	8116611116	0000010001000000
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416	442	0110100000	0000010001000010
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421	446	0110100101	000010001000110
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423	448	0110100111	0000018881881888
424	449	8118181888	<i>000001001001001</i>
425	450	B11 0101001	0000010001010000
426	451	0110101010	ØØØØØ10001010001
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431	456	0110101111	8686818681818118
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436	461	0110110100	0000010001100001
437	462	0110110101	0000010001100010
438	463	0110110110	<u>6666616661166611</u>
439	464	0110110111	0000010001100100
440	465	0110111000	8888818881188181
441	466	0110111001	0000010001100110
442	467	0110111010	0000010001100111
443	468	6116111611	0000319091191993
444	469	0110111100	0000010001101001
445	470	0110111101	0000010001110000
446	471	0110111110	0000010001110001
447	472	0110111111	0000010001110010
448	473	0111000000	0000010001110011
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450	475	0111000010	0000010001110101

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451	476	0111000011	0000010001110110
452	477	0111000100	0000010001110111
453	478	0111808181	0000010001111000
454	479	0111000110	0000010001111001
455	480	0111000111	66666166166666666
456	481	0111 001000	888881881886888
457	482	8111881881	000001001000000 10
458	483	8111881818	0 0808198188888 11
459	484	0111001011	6666616616666166
460	485	0111001100	00 000100190001 01
461	486	0111001101	0000010010000110
462	487	0111001110	0 00082160166861 11
463	488	0111001111	0000010010001000
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465	490	0111010001	0000010010010000
466	491	0111010010	0000010010010001
467	492	0111010011	0000010010010010
468	493	6111616168	0000010010010011
469	494	0111010101	0000010010010100
470	495	0111010110	0000010010010101
471	496	6111616111	8666618616616118
472	497	8111811888	0000010010010111
473	498	0111011001	0000010010011000
474	499	0111011010	0000010010011001
475	499	0111011011	0000010010011001
476	500	0111011100	0000010100000000
477	501		8888918188988891 888894 84 8888894 8
478	502		88888181888888818 888884 84 88838844
479 499	503 504	0111011111 0111100000	66666161666666611 6666664 646666664 66
480 481	504 505	0111100000 0111100001	8888818188888198 88888181888888181
482		0111100001 0111100010	88666161666666116
483	-000 507	8111188818	8888818188888119 8888818188888111
484	508	8111166166	0000010100001111 00000010100001111
485	509	8111166161	89868181888861881
486	510	6111166116	0000010100001000
487	511	8111168111	8898818188818881
488	512	0111101000	0000010100010010
489	513	0111101001	8888918188818911
490	514	8111181818	8866616188616168
491	515	8111181811	00000101000010101
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493	517	0111101101	0000010100010111
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496	520	0111110000	8886618188188888
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499	523	0111110011	0000016106100011
500	524	0111110100	0000010100100100
501	525	0111110101	8888818188188188181
502	526	0111110110	0000010100100110
503	527	0111110111	<i>0</i> 000010100100111
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