

RELIABILITY ANALYSIS AND OPTIMIZATION
OF COMPLEX SYSTEMS

A Dissertation
Presented to

the Faculty of the Department of Chemical Engineering
University of Houston

In Partial Fulfillment
of the Requirements for the Degree
Doctor of Philosophy in Chemical Engineering

by
Satish Loonkaran Gandhi
December, 1973

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ABSTRACT

Most of the earlier literature on system reliability optimization consider only the simple series-parallel systems subject to one or two constraints. Practical systems have complex rather than the simple series parallel configurations. With a view towards solving these complex system reliability optimization problems, an efficient computer algorithm based on the path enumeration method has been developed. An important feature of the method is the module representation of the reliability graph which considerably simplifies the calculation of reliability and sensitivity functions of complex systems. A modified integer gradient method is used for system optimization. Although the method does not insure a global optimum, it does find various near-optimum solutions. From a practical consideration, this could provide for a wider choice during the design phase.

In this research an effort has also been made to apply basic reliability concepts to process plants. A new formulation of the optimal reliability design of process plants which takes into account the quantitative aspects of systems throughput is proposed. It is based on the k-out-of-n configuration instead of the conventional parallel redundancy configuration. The problem is so formulated that determining the optimum configuration also determines the optimum capacity of units to be used at each stage of the system. A computer program based on a pseudo-Boolean algorithm is used to solve this non-linear integer programming problem.

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CHAPTER I

INTRODUCTION

Considerable literature has appeared in the last two decades in the area of system reliability analysis and optimization, particularly by researchers in electrical engineering and operations research. However, there are two areas where very little progress has been made.

Most of the earlier literature on system reliability optimization consider only the simple series-parallel systems subject to one or two constraints. The reason only this particular system has received so much attention is the separability and concavity of the reliability function. This makes it amenable to be solved using the techniques of dynamic programming [6,23,30,33,38,43,61], discrete maximum principle [19,55], integer programming [24,41,44,54,56], Lagrangian multipliers [2,3,18,38,40], and the dominating sequence concept [3,8,30,38,48]. Practical systems have complex rather than the simple series-parallel configurations, and, therefore, require methods capable of solving complex system optimization problems.

In the design of electrical, electronic, and safety systems, the flow of information is the most important factor, whereas in process systems the flow of materials and energy is also very important. Therefore, some innovations and modifications are required before the previous research in the area of reliability engineering can be effectively applied to chemical process systems. The vital role reliability plays in the operation of today's large, complex and expensive chemical plants has been widely recognized [10,17,26,35,58].

In this research an effort has been made to apply basic reliability concepts to process plants. In Chapter II a general computer algorithm based on path enumeration method has been developed to calculate the

reliability of complex systems. It also evaluates an important parameter; namely, the sensitivity of the system reliability to individual system units. The introduction of the module concept considerably simplifies the evaluation of the reliability and sensitivity functions. In Chapter III this algorithm is used in conjunction with a modified integer gradient method [50] for solving a large class of reliability and availability optimization problems. In Chapter IV an optimal reliability design of process systems is proposed which takes into account the quantitative aspects of systems throughput or capacity. The problem is so formulated that finding the optimum configuration, also determines the optimum capacity of units to be used at each stage of system.

All the procedures developed in this research have been programmed in FORTRAN IV. A number of illustrative examples are included in each chapter.

CHAPTER II

RELIABILITY AND SENSITIVITY ANALYSIS OF COMPLEX SYSTEMS USING FLOW GRAPH METHODS

In this chapter a general computer algorithm has been developed to calculate the system reliability when given reliabilities of each unit in the system and relationships between the units in the form of a reliability graph. Besides giving the system reliability, the programs calculate the sensitivity of the system reliability to individual system units.

Reliability

Reliability has been defined in many ways. A widely accepted definition [5] reads: "Reliability is the probability of a device performing its purpose adequately for the period of time intended under the operating conditions encountered."

The mathematical derivation of the reliability function is treated in most reliability texts [5,25,49,59]. Here the general formula is given without proof

$$R(t) = \exp \left[- \int_0^t z(\tau) d\tau \right] \quad (2.1)$$

where, the function $z(t)$ is called the failure rate function, hazard function, or hazard rate.

If the system has a failure time distribution with density function $f(t)$, then the failure time distribution function is given by

$$F(t) = \int_0^t f(\tau) d\tau \quad (2.2)$$

called the unreliability of the system at time t . Thus, the system reliability is given by

$$R(t) = 1 - F(t) = \exp \left(- \int_0^t z(\tau) d\tau \right) \quad (2.3)$$

From Equations (2.2) and (2.3) we obtain

$$z(t) = f(t)/R(t) \quad (2.4)$$

Typical failure time density functions with their corresponding reliability and failure rate functions are:

(1) Exponential

$$f(t) = \lambda e^{-\lambda t} \quad R(t) = e^{-\lambda t} \quad z(t) = \lambda$$

(2) Weibull

$$f(t) = \lambda \alpha t^{\alpha-1} e^{-\lambda t^\alpha} \quad R(t) = e^{-\lambda t^\alpha} \quad z(t) = \lambda \alpha t^{\alpha-1}$$

(where $\alpha, \lambda > 0$)

(3) Gamma

$$f(t) = \frac{\lambda}{(r-1)!} (\lambda t)^{r-1} e^{-\lambda t} \quad R(t) = \sum_{j=0}^{r-1} \frac{(\lambda t)^j e^{-\lambda t}}{j!} \quad z(t) = \frac{\lambda^r t^{r-1}}{(r-1)!} \frac{e^{-\lambda t}}{R(t)}$$

(r is a positive integer)

In the following it is assumed that individual units forming a system have exponential failure time density functions. The analysis, however, with other failure time density functions is an extension of the techniques presented in this thesis.

Methods for Computing Complex System Reliability

1. State Enumeration or Event-Space Method

In this method we first make a list of all possible mutually exclusive states of the system. A state is defined by listing the successful

and failed elements in the system. In general the total number of states will be given by 2^n , where n is the number of units or elements in the system. Next, we identify the states which result in successful system operation and determine the probability of occurrence of each successful state. Finally, we sum up all the successful state probabilities which gives the system reliability. Although the method is easily programmed [9] it is computationally not feasible for systems having large number of elements.

2. Network Reduction Method

In this method the series, parallel and series-parallel subsystems are combined until we end up with a non-series parallel system which cannot be further reduced. Reference [45] then suggests the use of a factoring theorem. A particular element x is selected, and the two networks are obtained and generated when x is replaced by a short circuit (perfect connection) and an open circuit. If these two networks are simple series-parallel, they can be reduced. Otherwise the next block, y , must be selected and the procedure repeated.

Suppose block x is selected and the two networks are generated. The system reliability is then [45]

$$R = R_x * R \Big|_{R_x=1} + (1-R_x) * R \Big|_{R_x=0} \quad (2.5)$$

where R_x is the reliability of block x . The method is discussed in [1,13,14,39,45].

3. Path Enumeration and Cut Enumeration Methods

These methods can be used to determine the reliability of any system not containing dependent failures [16,29,52]. A path is a set of elements which form a connection between input and output when traversed

in a stated direction, A minimal path is one in which no node is traversed more than once in tracing the path.

Let P_i ($i=1, 2, \dots, M$) be the i th minimal path in the system. If any path is operable, the system performs adequately. Thus, the system reliability is

$$R = P_r \left\{ \bigcup_{i=1}^M P_i \right\} \quad (2.6)$$

where \cup denotes the union.

By use of the expansion rule for the probability of the union of M events [21], we have the formula

$$\begin{aligned} R = & \sum_{i=1}^M P_r \{P_i\} - \sum_{i=1}^M \sum_{j>i}^M P_r \{P_i \cap P_j\} \\ & + \sum_{i=1}^M \sum_{j>i}^M \sum_{k>j}^M P_r \{P_i \cap P_j \cap P_k\} + \dots + (-1)^{M-1} P_r \left\{ \bigcap_{i=1}^M P_i \right\} \end{aligned} \quad (2.7)$$

where \cap denotes the intersection.

A cut is defined as a set of elements that if they fail, the system will fail regardless of the condition of the other elements in the system. A minimal cut is one in which there is no proper subset of elements whose failure alone will cause the system to fail.

Let C_j ($j=1, 2, \dots, M$) be the minimal cuts in the system. If any cut is inoperable, the system fails. The system reliability is thus given by

$$R = 1 - P_r \left\{ \bigcup_{j=1}^M \bar{C}_j \right\} \quad (2.8)$$

$$\begin{aligned}
= 1 - \sum_{j=1}^M P_r \left\{ \overline{C}_j \right\} + \sum_{j=1}^M \sum_{k>j}^M P_r \left\{ \overline{C}_j \cap \overline{C}_k \right\} - \sum_{j=1}^M \sum_{k>j}^M \sum_{l>k}^M P_r \left\{ \overline{C}_j \cap \overline{C}_k \cap \overline{C}_l \right\} \\
+ \dots + (-1)^j P_r \left\{ \bigcap_{i=1}^M \overline{C}_j \right\}
\end{aligned} \tag{2.9}$$

where \overline{C}_j denotes the complement of the event C_j ; i.e., \overline{C}_j denotes the failure of all elements of the cut C_j .

The algorithm developed in this chapter is based both on the network reduction and path enumeration methods. Advantages and disadvantages of this approach have been discussed and compared with typical algorithms based on the state enumeration or event space methods.

Module Representation of Reliability Graphs

A reliability flow graph consists of nodes, branches, and modules with directed arrows between them. Consider the modules in the complexly connected reliability flow graphs shown in Figures 2.1, 2.2, and 2.4. The reliability flow graph of Figure 2.1 represents a series-parallel-series system distinct from simple series, parallel or series-parallel system. The reliability flow graph of Figures 2.2 and 2.4 represent non-series-parallel systems. In the module representation of reliability graphs more than one module between any two arbitrary nodes is permitted. A module is defined to be a single unit or simply connected parallel units. A configuration permitted as a module is limited to either of the following:

1. Single Unit Module

A single unit module is one which consists of a single unit. The module reliability is the reliability of the unit itself.

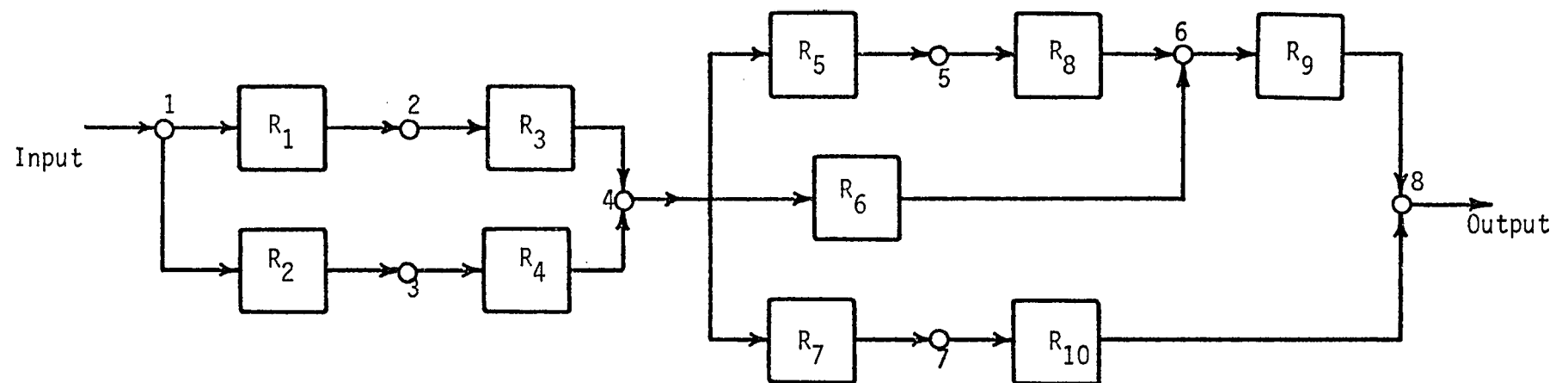


Figure 2.1

Reliability Graph - Series-Parallel-Series System

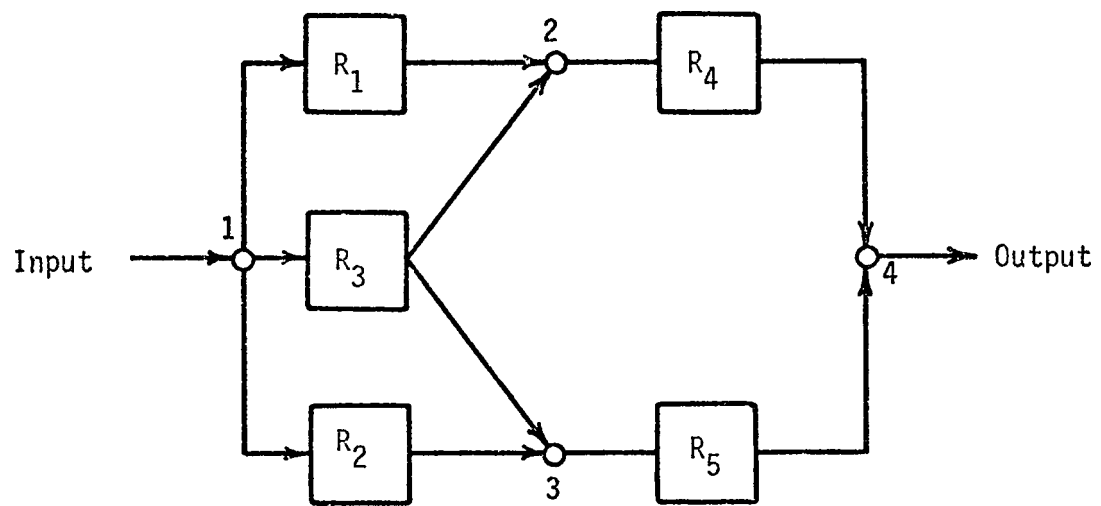


Figure 2.2

Reliability Graph - Non-Series-Parallel System

2. Multi-Unit Module

A module which consists of multiple parallel units is either an active redundancy module or a standby redundancy module as shown in Figure 2.3. We distinguish between three standby redundancy modules. A unit is called a cold standby when the failure probability in standby is zero. When the failure probability is the same in standby as in the service, it is called a hot standby. Cases that lie in between these two extremes are called warm standby. Depending on the type of redundancy, we have different expressions for module reliability as shown in Appendix A.

The advantages of the module representation of a reliability graph will become clear in the next few sections.

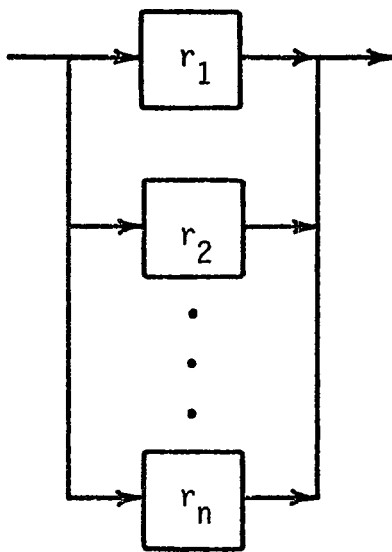
Reliability Calculation

The system reliability is defined here as the probability of the successful functioning of all the modules in at least one minimal path. A minimal path is one which contains a minimum number of modules and no node is traversed more than once in tracing the path.

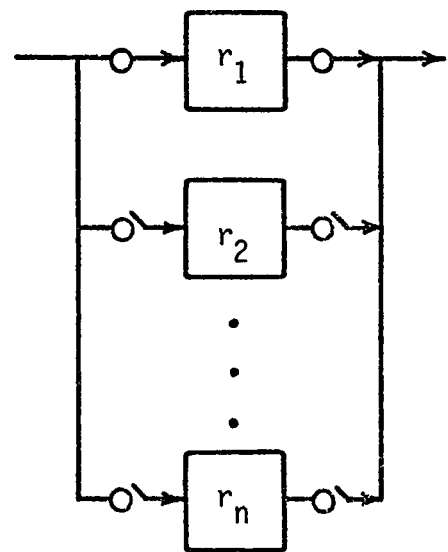
If P_i ($i=1, 2, \dots, M$) denotes the minimal paths in the system, the system reliability R is as shown in Equations (2.6) and (2.7).

The system reliability R is thus given in terms of module reliabilities by

$$\begin{aligned}
 R = & \sum_{i=1}^M \prod_{l \in P_i} R_l - \sum_{i=1}^M \sum_{j>i}^M \prod_{l \in P_i \cup P_j} R_l \\
 & + \sum_{i=1}^M \sum_{j>i}^M \sum_{k>j}^M \prod_{l \in P_i \cup P_j \cup P_k} R_l + \dots + (-1)^{M-1} \prod_{l \in \bigcup_{i=1}^M P_i} R_l
 \end{aligned} \tag{2.10}$$



Active Redundancy Module



Standby Redundancy Module

Figure 2.3
Multi-Unit Module

where the members of the i th path, the union of the i th and the j th paths, etc., are denoted by $l \in P_i$, $l \in P_i \cup P_j$, etc.

The total number of terms Z involved in equation (2.10) is given by

$$Z = 2^M - 1 \quad (2.11)$$

It should be noted in Equations (2.7) and (2.10) that if, for example, the path P_1 has modules R_1 , R_2 , and R_3 , and the path P_2 has modules R_2 , R_3 , R_4 , and R_5 , then $P_r \{P_1 \cap P_2\}$ is given by $R_1 R_2 R_3 R_4 R_5$, not $R_1 R_2^2 R_3^2 R_4 R_5$.

Sensitivity Calculation

It is very interesting and important that the system reliability expression given by Equation (2.10) is a bilinear function of each module's reliability. By use of this property, the sensitivity of the system reliability to the module reliability R_i can be obtained by the simple rule

$$S_{R_i} = \partial R / \partial R_i = R \Big|_{R_i=1} - R \Big|_{R_i=0} \quad i=1, 2, \dots, N \quad (2.12)$$

where N is the total number of modules in the system. We call S_{R_i} a module sensitivity.

Two types of sensitivities are of interest. The first is the sensitivity of the system reliability to a unit reliability. A module sensitivity is itself of this type if the module is a single-unit module. For a multiple-unit module, the sensitivity is given by

$$S_{r_{ij}} = \frac{\partial R}{\partial r_{ij}} = \frac{\partial R}{\partial R_i} * \frac{\partial R_i}{\partial r_{ij}} \quad (2.13)$$

where r_{ij} is the reliability of the j th unit of the i th module. For an active redundancy module with different units, we have the following expression:

$$S_{r_{ij}} = S_{R_i} \left[R_i \Big|_{r_{ij}=1} - R_i \Big|_{r_{ij}=0} \right] = S_{R_i} \prod_{\substack{j=1 \\ j \neq i}}^{N_i} (1-r_{ij}) \quad (2.14)$$

where N_i is the number of units in the i th module.

The second type of sensitivity is the sensitivity of the system reliability to the number of units in a module. This type of sensitivity is defined for the active redundancy module or the standby redundancy module with identical units.

$$S_{N_i} = \frac{\partial R}{\partial N_i} = \frac{\partial R}{\partial R_i} * \frac{\partial R_i}{\partial N_i} = S_{R_i} * \frac{\partial R_i}{\partial N_i} \quad (2.15)$$

The sensitivity for the active redundancy module with identical units is given by

$$S_{N_i} = S_{R_i} \left[- (1-r_i)^{N_i} \ln (1-r_i) \right] \quad (2.16)$$

The sensitivity for the standby modules can be approximately calculated by

$$S_{N_i} \approx S_{R_i} \left[R_i (N_i+1) - R_i (N_i) \right] \quad (2.17)$$

A large sensitivity value means that we may increase the system reliability greatly by increasing the corresponding individual reliability or by increasing the number of units in the corresponding module. Sensitivity, therefore, is a measure of system reliability improvement and provides important information for maximizing the reliability of complex systems. Sensitivity expressions for multi-unit modules are presented in Appendix A.

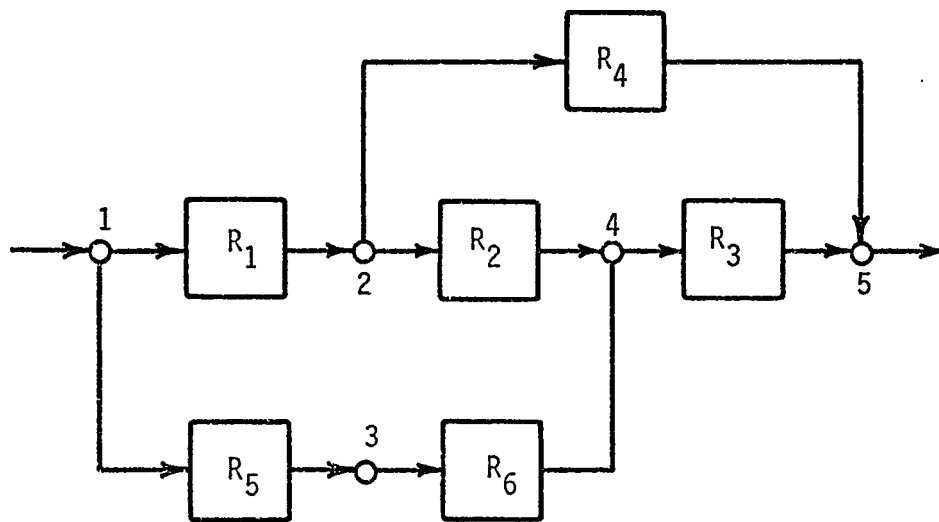


Figure 2.4

Reliability Graph - Non-Series-Parallel System

Basic Algorithm

The total system reliability and sensitivities can be obtained by the following six-step procedure:

1. Find all minimal paths using the reliability graph;
2. Find all the required unions of the paths;
3. Give each path union a reliability expression in terms of module reliability;
4. Sum up all the reliability expressions obtained above according to Equation (2.10);
5. Evaluate the system reliability by substituting numerical values of each module reliability; and,
6. Evaluate the desired sensitivities from Equations (2.12) to (2.17).

A brief description of each step is given and illustrated by way of a non-series-parallel reliability graph in Figure 2.4.

Path Finding Algorithm

We require a path-finding routine which detects all of the minimal paths connecting input and output when given a system reliability graph. The path finding algorithm developed by Henley and Williams [27] is ideally suited for this purpose. Having been given the reliability graph of Figure 2.4, the routine detects three paths in the node form shown in Table (2.1). By the nature of the algorithm only minimal paths are detected. Due to the possibility that a branch may have more than one module and/or a module may be present in more than one branch, the path finding algorithm [27] was modified to trace the modules in each minimal path as shown in Table 2.1. This allows a unique representation of each

PATH	TRACE NODES	TRACE MODULES
P1	1 \longrightarrow 2 \longrightarrow 4 \longrightarrow 5	R1 \longrightarrow R2 \longrightarrow R3
P2	1 \longrightarrow 2 \longrightarrow 5	R1 \longrightarrow R4
P3	1 \longrightarrow 3 \longrightarrow 4 \longrightarrow 5	R5 \longrightarrow R6 \longrightarrow R3

Table 2.1

PATH DETECTION

path which in general is not possible when paths are traced in the node form. Details of the modified path finding algorithm are given in Appendix B.

Another path finding algorithm suitable for reliability calculation and its Fortran listing are available in the literature [4,46].

Algorithm to Find Path-Unions

Each detected path is converted into the binary form shown in Figure 2.5. The locations containing 1's indicate the modules present in the path. N is the total number of modules in the system. Each path is given a positive sign. This is the sign of the first summation term in the reliability expression given by Equation (2.7) or (2.10).

The next step is to generate the binary representation of all the necessary unions of paths. This is easily done by using a bitwise logical OR operation in the order of ascending path number. That is, $(P_1) \text{ OR } (P_2)$, $(P_1) \text{ OR } (P_3)$, ..., $(P_1) \text{ OR } (P_M)$, $(P_2) \text{ OR } (P_3)$, ..., $(P_2) \text{ OR } (P_M)$, ..., $(P_{M-1}) \text{ OR } (P_M)$, which completes the 2-path unions. We then proceed to 3-path unions, $((P_1) \text{ OR } (P_2)) \text{ OR } (P_3)$, ..., $((P_1) \text{ OR } (P_2)) \text{ OR } (P_M)$, ..., $((P_{M-2}) \text{ OR } (P_{M-1})) \text{ OR } (P_M)$. Similar procedures continue until we reach M -path union, $(P_1) \text{ OR } (P_2) \text{ OR } \dots \text{ OR } (P_{M-1}) \text{ OR } (P_M)$. The sign changes with the generation of each higher order path union. Thus, it is negative for 2-union paths, positive for 3-union paths, ..., and $(-1)^{M-1}$ for M -union path. Details of the algorithm that generates all the necessary path unions are given in Appendix C.

The binary representation of the paths and path-unions for the reliability graph of Figure 2.4 are shown in Figure 2.6. Each path or path union corresponds to a term in Equation (2.10).

Sign		Module				
	1	2	3	4	...	N
+	1	0	0	1	...	1

Figure 2.5

Binary Representation of a Path or a Path Union

		MODULE					
	SIGN	1	2	3	4	5	6
P1	+	1	1	1	0	0	0
P2	+	1	0	0	1	0	0
P3	+	0	0	1	0	1	1
} MINIMUM PATHS							
P1 OR P2	-	1	1	1	1	0	0
P1 OR P3	-	1	1	1	0	1	1
P2 OR P3	-	1	0	1	1	1	1
} 2 PATH UNIONS							
P1 OR P2 OR P3	+	1	1	1	1	1	1
3 PATH UNION							

Figure 2.6

Binary Representation of 3 Paths and Unions of Paths

System Reliability

We now pick up every module that has a 1 in the corresponding location and form the product expressions for each path and path union. The collection of all the terms with their proper signs gives the required system reliability function as given in Equation (2.10).

For example, the system reliability expression in Figure 2.6 is given as

$$\begin{aligned}
 R = & (R_1 R_2 R_3 + R_1 R_4 + R_3 R_5 R_6) \\
 & - (R_1 R_2 R_3 R_4 + R_1 R_2 R_3 R_5 R_6 + R_1 R_3 R_4 R_5 R_6) \\
 & + R_1 R_2 R_3 R_4 R_5 R_6
 \end{aligned} \tag{2.18}$$

When given numerical values for the reliability of each module, we can obtain the system reliability. The module sensitivity S_{R_i} can also be calculated from Equation (2.12) by evaluating the system reliability twice, once with $R_i=1$ and once with $R_i=0$. The other sensitivities can also be calculated through the module sensitivities by Equations (2.13) to (2.17).

A number of terms in the reliability expression given in Equation (2.10) and generated as indicated above are identical and have opposite signs. A significant number of such terms can be eliminated by modifying the algorithm that generates the path-unions. Details of this modification together with an example problem are discussed in Appendix C. This reduces the storage requirement and computation time considerably, especially during system optimization discussed in the next chapter.

Comparison with State Enumeration Algorithm

An alternate method of calculating system reliability is the method of state enumeration. The algorithm for the state enumeration can

be briefly stated as follows [9]:

1. Find all the possible combinations of the states of the units (up or down);
2. Find each combination which connects input and output;
calculate the product of the unreliabilities of the down units and the reliabilities of the up units;
3. Sum up the products obtained in Step 2. This gives the system reliability expression.

The algorithm can be easily computerized by using Boolean algebra [9]. As an example, consider the reliability graph shown in Figure 2.7. The algorithm can be understood by referring to Table 2.2. From the table, the system reliability function is given by

$$\begin{aligned}
 R = & (1-r_1)(1-r_2)r_3(1-r_4)r_5 + (1-r_1)(1-r_2)r_3r_4r_5 \\
 & + (1-r_1)r_2(1-r_3)r_4(1-r_5) + (1-r_1)r_2(1-r_3)r_4r_5 \\
 & + (1-r_1)r_2r_3(1-r_4)r_5 + (1-r_1)r_2r_3r_4(1-r_5) \\
 & + (1-r_1)r_2r_3r_4r_5 + r_1(1-r_2)(1-r_3)r_4(1-r_5) \\
 & + r_1(1-r_2)(1-r_3)r_4r_5 + r_1(1-r_2)r_3(1-r_4)r_5 \\
 & + r_1(1-r_2)r_3r_4(1-r_5) + r_1(1-r_2)r_3r_4r_5 \\
 & + r_1r_2(1-r_3)r_4(1-r_5) + r_1r_2(1-r_3)r_4r_5 \\
 & + r_1r_2r_3(1-r_4)r_5 + r_1r_2r_3r_4(1-r_5) + r_1r_2r_3r_4r_5
 \end{aligned} \tag{2.19}$$

Paths and path unions required in the path enumeration algorithm are listed in Table 2.3. From Table 2.3 the system reliability function is easily derived as

$$R = r_1r_4 + r_2r_4 + r_3r_5 - (r_1r_2r_4 + r_1r_3r_4r_5 + r_2r_3r_4r_5) + r_1r_2r_3r_4r_5 \tag{2.20}$$

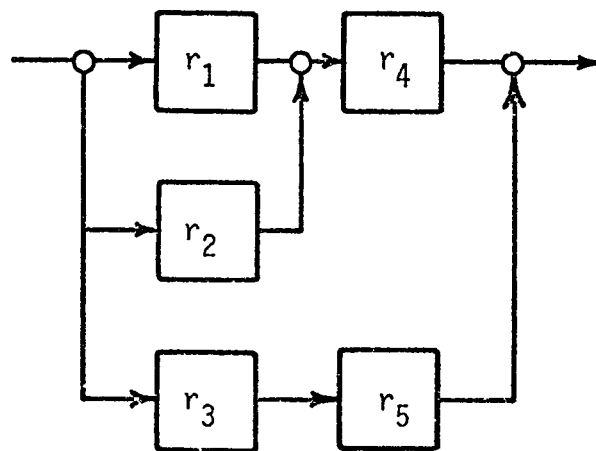


Figure 2.7
Reliability Graph

STATE NUMBER	BINARY NUMBER					GIVING A PATH?
	r ₁	r ₂	r ₃	r ₄	r ₅	
0	0	0	0	0	0	No
1	0	0	0	0	1	No
2	0	0	0	1	0	No
3	0	0	0	1	1	No
4	0	0	1	0	0	No
5	0	0	1	0	1	Yes
6	0	0	1	1	0	No
7	0	0	1	1	1	Yes
8	0	1	0	0	0	No
9	0	1	0	0	1	No
10	0	1	0	1	0	Yes
11	0	1	0	1	1	Yes
12	0	1	1	0	0	No
13	0	1	1	0	1	Yes
14	0	1	1	1	0	Yes
15	0	1	1	1	1	Yes
16	1	0	0	0	0	No
17	1	0	0	0	1	No
18	1	0	0	1	0	Yes
19	1	0	0	1	1	Yes
20	1	0	1	0	0	No
21	1	0	1	0	1	Yes
22	1	0	1	1	0	Yes
23	1	0	1	1	1	Yes
24	1	1	0	0	0	No
25	1	1	0	0	1	No
26	1	1	0	1	0	Yes
27	1	1	0	1	1	Yes
28	1	1	1	0	0	No
29	1	1	1	0	1	Yes
30	1	1	1	1	0	Yes
31	1	1	1	1	1	Yes

Table 2.2 All combination of the states

PATHS AND PATH UNIONS	SIGN	MODULE				
		r_1	r_2	r_3	r_4	r_5
P1	+	1	0	0	1	0
P2	+	0	1	0	1	0
P3	+	0	0	1	0	1
P1 OR P2	-	1	1	0	1	0
P1 OR P3	-	1	0	1	1	1
P2 OR P3	-	0	1	1	1	1
(P1 OR P2) OR P3	+	1	1	1	1	1

Table 2.3

Paths and Path Unions

Although the two expressions given by Equations (2.19) and (2.20) are different in form, it can be easily verified that both are equivalent.

We can conclude from the above comparison and from more general consideration that the advantages of the path enumeration method are:

1. Steps to reach the system reliability function are much less if the number of paths is fairly small and is unaffected by the total number of modules in the system;
2. The number of terms in the system reliability function are fewer; and,
3. Any term in the system reliability function contains equal or fewer multipliers.

The module structure introduced earlier in the chapter brings about a further reduction in the number of paths. Thus, the units r_1 and r_2 are combined and treated as the active redundancy modules R_1 as shown in Figure 2.8. The path enumeration algorithm then gives the path and path unions shown in Table 2.4. The system reliability is given by

$$R = R_1 r_4 + r_3 r_5 - R_1 r_3 r_4 r_5 \quad (2.21)$$

where

$$R_1 = 1 - (1-r_1)(1-r_2) = r_1 + r_2 - r_1 r_2 \quad (2.22)$$

In this way, the previously stated advantages of the path enumeration method are enhanced by introducing the module structure. Since it is customary to use parallel redundancy configurations to increase the reliability of a weak unit, the module concept significantly reduces the number of paths in the system thus reducing the computing time as shown by the following examples:

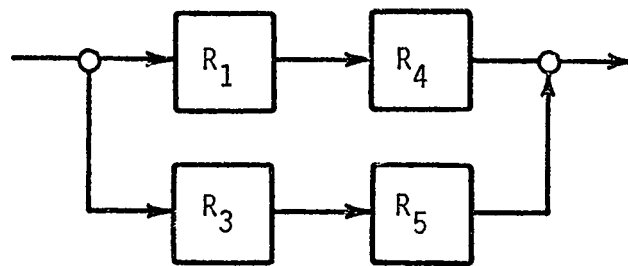


Figure 2.8

Module Representation of Figure 2.7

PATHS AND PATH UNION	SIGN	MODULE			
		R_1	r_3	r_4	r_5
P1	+	1	0	1	0
P2	+	0	1	0	1
P1 OR P2	-	1	1	1	1

Table 2.4

Paths and Path Union

Examples

At first, consider the fairly complicated reliability graph [46] shown in Figure 2.9. The reliability of each unit is given in Table 2.5. In this original graph, there are 16 units and 55 paths. The number of path and path unions required for calculating the system reliability, therefore, will be $2^{55}-1$ from Equation (2.11), which is far in excess of 10^{16} . The application of the path enumeration method to the original graph is impossible unless some approximations are introduced.

By using the module concept, the graph is automatically converted to Figure 2.1, where R_1 , R_3 , R_5 , R_9 , and R_{10} are actively connected multi-unit modules which consist of units (r_1, r_2) , (r_4, r_5) , (r_7, r_8) , (r_{12}, r_{13}, r_{14}) and (r_{15}, r_{16}) , respectively; R_2 , R_4 , R_6 , R_7 , and R_8 are single unit modules consisting of units r_3 , r_6 , r_9 , r_{10} , and r_{11} , respectively. The module representation of the reliability graph is shown in Table 2.6.

The converted graph has only six paths as shown in Table 2.7. The binary representation of the six paths is shown in Table 2.8. Since the number of paths is greatly reduced, the path enumeration algorithm discussed earlier can be applied easily and efficiently to produce the system reliability value of 0.97043 in a few seconds.

The sensitivities are listed in Table 2.9, where attention is focused on the first type of sensitivity, $\partial R / \partial r_{ij}$ given by Equation (2.14). The table shows that the module R_4 , that is, unit r_6 , has the larger sensitivity value (0.085), and the first unit of module R_{10} , unit r_{15} , has the smallest. It is, therefore, more efficient to increase the unit reliability of r_6 in order to increase the total reliability of the system.

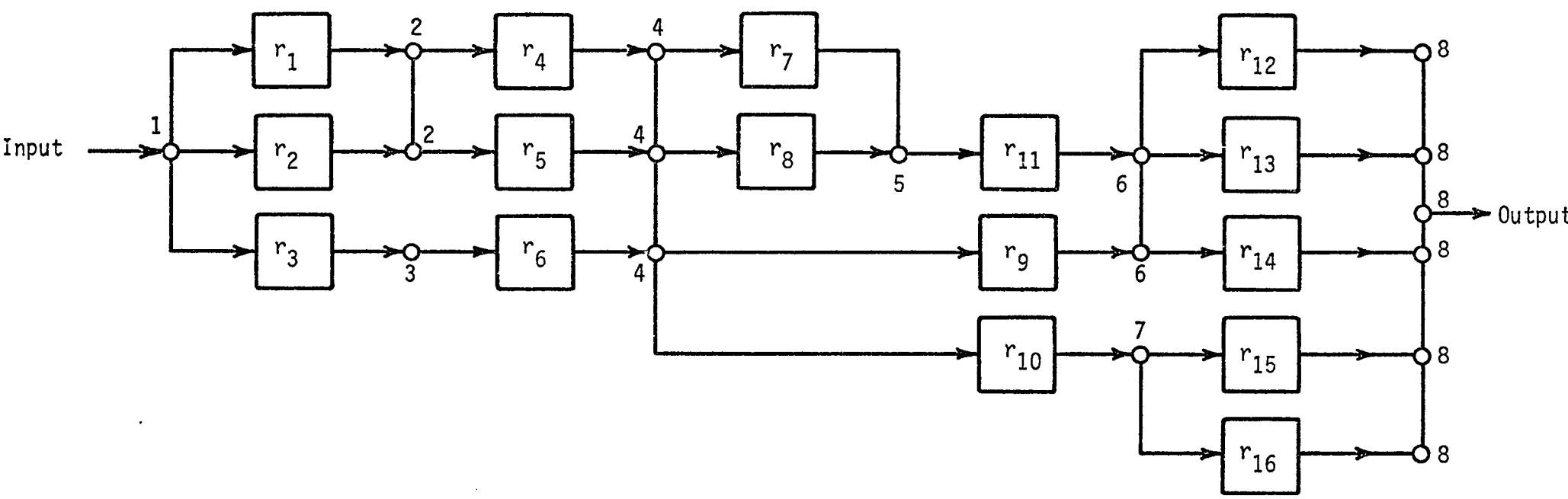


Figure 2.9

Reliability Graph - 55 Paths Between Input and Output

UNIT	RELIABILITY
r_1	0.80
r_2	0.70
r_3	0.90
r_4	0.75
r_5	0.85
r_6	0.87
r_7	0.82
r_8	0.62
r_9	0.88
r_{10}	0.75
r_{11}	0.89
r_{12}	0.85
r_{13}	0.75
r_{14}	0.65
r_{15}	0.70
r_{16}	0.90

Table 2.5
Unit Reliabilities

MODULE	BRANCH	UNIT RELIABILITY	NUMBER OF UNITS
1	1 → 2	0.800000 0.700000	2
2	1 → 3	0.900000	1
3	2 → 4	0.750000 0.850000	2
4	3 → 4	0.870000	1
5	4 → 5	0.820000 0.620000	2
6	4 → 6	0.880000	1
7	4 → 7	0.750000	1
8	5 → 6	0.890000	1
9	6 → 8	0.850000 0.750000 0.650000	3
10	7 → 8	0.700000 0.900000	2

Table 2.6

Module Representation

PATH	TRACE NODES	TRACE MODULES
P1	1 → 2 → 4 → 6 → 8 →	1 → 3 → 6 → 9
P2	1 → 3 → 4 → 6 → 8 →	2 → 4 → 6 → 9
P3	1 → 2 → 4 → 5 → 6 → 8 →	1 → 3 → 5 → 8 → 9
P4	1 → 3 → 4 → 5 → 6 → 8 →	2 → 4 → 5 → 8 → 9
P5	1 → 2 → 4 → 7 → 8 →	1 → 3 → 7 → 10
P6	1 → 3 → 4 → 7 → 8 →	2 → 4 → 7 → 10

Table 2.7

Path Detection

PATH	MODULE									
	1	2	3	4	5	6	7	8	9	10
P1	1	0	1	0	0	1	0	0	1	0
P2	0	1	0	1	0	1	0	0	1	0
P3	1	0	1	0	1	0	0	1	1	0
P4	0	1	0	1	1	0	0	1	1	0
P5	1	0	1	0	0	0	1	0	0	1
P6	0	1	0	1	0	0	1	0	0	1

Table 2.8

Binary Representation of Paths

MODULE	MODULE RELIABILITY	MODULE SENSITIVITY $\partial R / \partial R_i$	UNIT SENSITIVITY $\partial R_i / \partial r_{ij}$	SYSTEM SENSITIVITY $\partial R / \partial r_{ij}$
1	0.940000	0.206964	0.300000 0.200000	0.620891E-01 0.413928E-01
2	0.900000	0.821141E-01	1.00000	0.821141E-01
3	0.962500	0.202126	0.150000 0.250000	0.303189E-01 0.505314E-01
4	0.870000	0.849457E-01	1.00000	0.849456E-01
5	0.931600	0.281274E-01	0.380000 0.180000	0.106884E-01 0.506293E-02
6	0.880000	0.450027E-01	1.00000	0.450027E-01
7	0.750000	0.316913E-01	1.00000	0.316913E-01
8	0.890000	0.294421E-01	1.00000	0.294421E-01
9	0.986875	0.261395	0.875000E-01 0.525000E-01 0.375000E-01	0.228721E-01 0.137233E-01 0.980232E-02
10	0.970000	0.245036E-01	0.100000 0.300000	0.245036E-02 0.735107E-02

Table 2.9
Sensitivities $\partial R / \partial r_{ij}$

Moreover, the sensitivity tells us that the increase in total reliability will amount to approximately 0.0085 when the increases in the unit reliability of r_6 is 0.1.

As the second example, we consider a safety system [53].

Figure 2.10 shows a boiler DDC safety system (pressure safety system). Reliabilities of each unit are listed in Table 2.10. The original graph has 26 paths. After being converted to a module configuration shown in Figure 2.11 and Table 2.11, the number of paths is reduced to only nine as shown in Tables 2.12 and 2.13.

The total reliability of the system is computed to be 0.99927. The sensitivities are listed in Table 2.14. In this case, the sensitivities are evaluated with respect to the number of units in a module, that is, $\partial R / \partial N_i$ defined by Equation (2.16). The table shows that the module R_4 has the largest sensitivity and the module R_5 has the smallest sensitivity. Adding one more unit (Converter 1) to the module R_5 is of no use, as the increase in the total reliability will only be 0.0000009.

Advantages of Module Structure

It can be seen from the above examples and more general consideration that the introduction of module structure has the following advantages:

1. It considerably reduces the number of paths and, hence, provides a means for the efficient evaluation of the reliability of complex systems;
2. It greatly simplifies the evaluation of the sensitivity function, an important design parameter that will be used in system optimization;

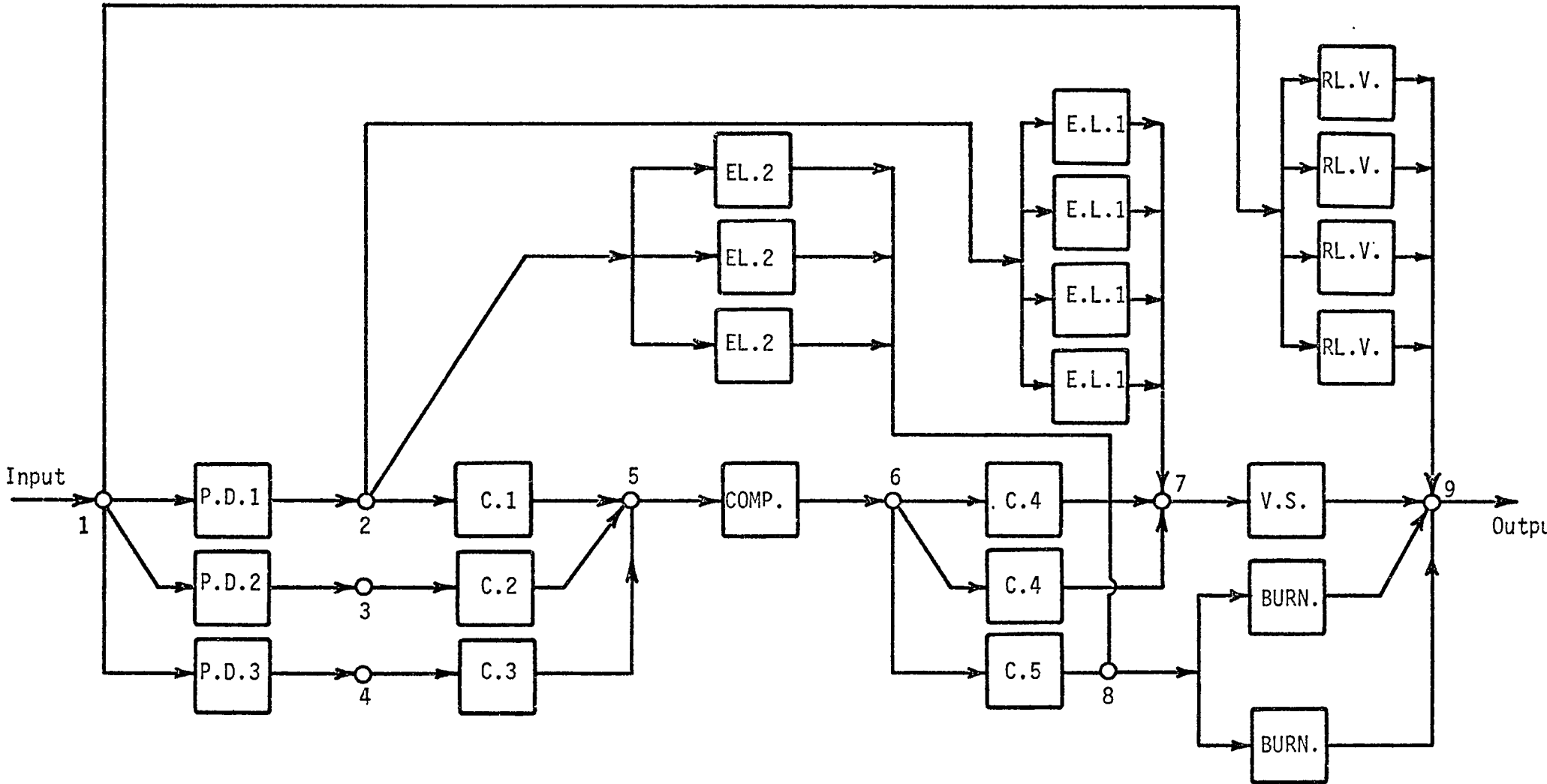


Figure 2.10

Reliability Graph - Boiler DDC Safety System

UNIT	RELIABILITY
P.D. 1: Pressure Detector 1	0.68
P.D. 2: Pressure Detector 2	0.75
P.D. 3: Pressure Detector 3	0.75
C.1 : Converter 1	0.70
C.2 : Converter 2	0.70
C.3 : Converter 3	0.70
C.4 : Converter 4	0.70
C.5 : Converter 5	0.70
COMP. : Computer	0.70
V.S. : Valve Servo	0.75
BURN. : Burner	0.68
RL.V. : Relief Valve	0.75
E.L. 1: Emergency Line 1	0.68
E.L. 2: Emergency Line 2	0.68

Table 2.10

Unit Reliabilities

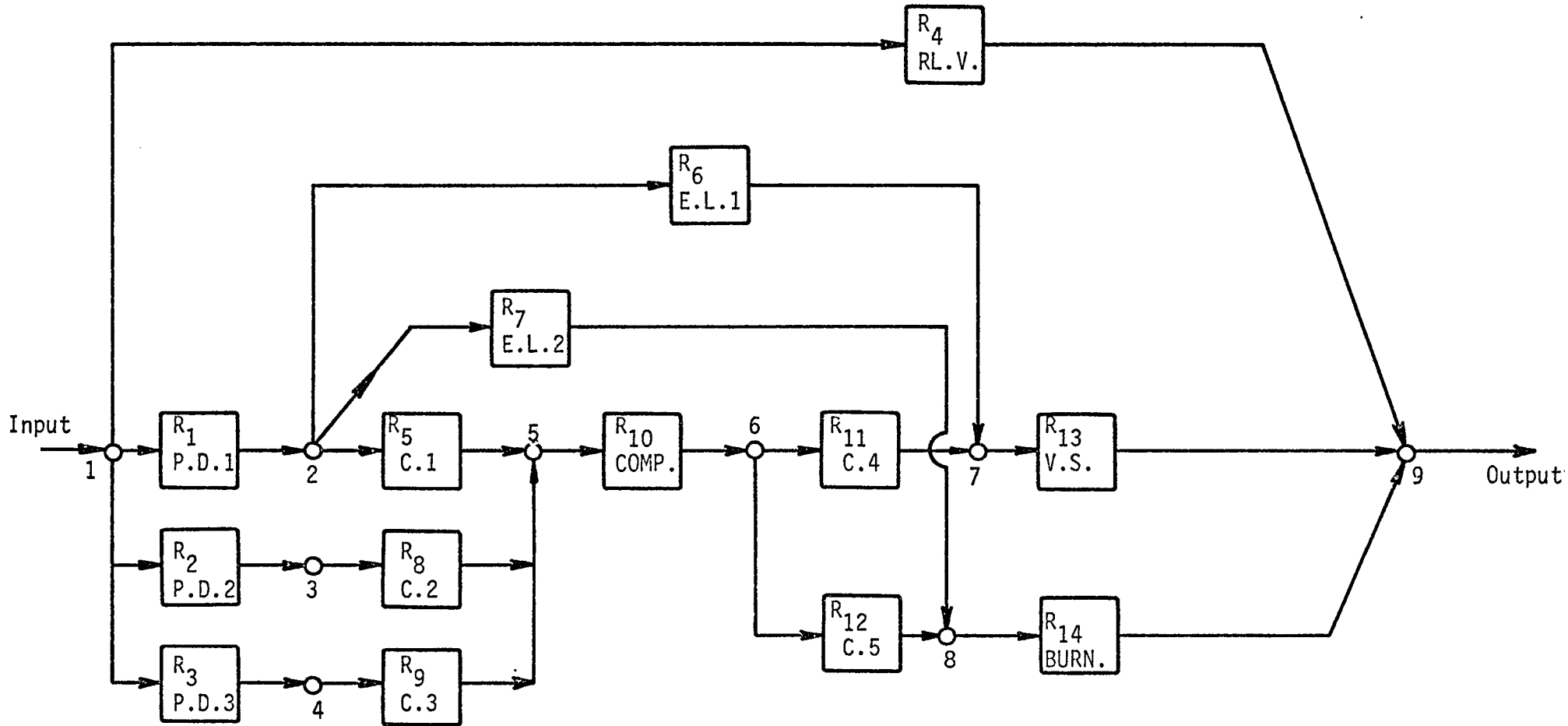


Figure 2.11

Module Representation of the Boiler DDC Safety System

MODULE	BRANCH			UNIT RELIABILITY	NUMBER OF UNITS
1	1	→	2	0.680000	1
2	1	→	3	0.750000	1
3	1	→	4	0.750000	1
4	1	→	9	0.750000	4
5	2	→	5	0.700000	1
6	2	→	7	0.680000	4
7	2	→	8	0.680000	3
8	3	→	5	0.700000	1
9	4	→	5	0.700000	1
10	5	→	6	0.700000	1
11	6	→	7	0.700000	2
12	6	→	8	0.700000	1
13	7	→	9	0.750000	1
14	8	→	9	0.680000	2

Table 2.11

Module Representation

PATH	TRACE NODES	TRACE MODULES
P1	1 → 9 →	4
P2	1 → 2 → 7 → 9 →	1 → 6 → 13
P3	1 → 2 → 5 → 6 → 7 → 9 →	1 → 5 → 10 → 11 → 13
P4	1 → 3 → 5 → 6 → 7 → 9 →	2 → 8 → 10 → 11 → 13
P5	1 → 4 → 5 → 6 → 7 → 9 →	3 → 9 → 10 → 11 → 13
P6	1 → 2 → 8 → 9 →	1 → 7 → 14
P7	1 → 2 → 5 → 6 → 8 → 9 →	1 → 5 → 10 → 12 → 14
P8	1 → 3 → 5 → 6 → 8 → 9 →	2 → 8 → 10 → 12 → 14
P9	1 → 4 → 5 → 6 → 8 → 9 →	3 → 9 → 10 → 12 → 14

Table 2.12

Path Detection

PATH	MODULE													
	1	2	3	4	5	6	7	8	9	10	11	12	13	14
P1	0	0	0	1	0	0	0	0	0	0	0	0	0	0
P2	1	0	0	0	0	1	0	0	0	0	0	0	1	0
P3	1	0	0	0	1	0	0	0	0	1	1	0	1	0
P4	0	1	0	0	0	0	0	1	0	1	1	0	1	0
P5	0	0	1	0	0	0	0	0	1	1	1	0	1	0
P6	1	0	0	0	0	0	1	0	0	0	0	0	0	1
P7	1	0	0	0	1	0	0	0	0	1	0	1	0	1
P8	0	1	0	0	0	0	0	1	0	1	0	1	0	1
P9	0	0	1	0	0	0	0	0	1	1	0	1	0	1

Table 2.13
Binary Representation of Paths

MODULE	MODULE RELIABILITY	MODULE SENSITIVITY $\partial R / \partial R_i$	UNIT SENSITIVITY $\partial R_i / \partial N_i$	SYSTEM SENSITIVITY $\partial R / \partial N_i$
1	0.680000	0.192148E-02	0.364619	0.700609E-03
2	0.750000	0.257737E-03	0.346574	0.893248E-04
3	0.750000	0.257737E-03	0.346574	0.893248E-04
4	0.996094	0.187413	0.541521E-02	0.101488E-02
5	0.700000	0.256121E-05	0.361192	0.925088E-06
6	0.989514	0.104235E-03	0.119478E-01	0.124538E-05
7	0.967232	0.330605E-03	0.373370E-01	0.123438E-04
8	0.700000	0.276147E-03	0.361192	0.997419E-04
9	0.700000	0.276147E-03	0.361192	0.997419E-04
10	0.700000	0.868859E-03	0.361192	0.313825E-03
11	0.910000	0.190398E-03	0.108357	0.206310E-04
12	0.700000	0.205885E-03	0.361192	0.743638E-04
13	0.750000	0.542155E-03	0.346574	0.187897E-03
14	0.897600	0.811069E-03	0.116678	0.946339E-04

Table 2.14
Sensitivities $\partial R / \partial N_i$

3. Although the module itself may be composed of multiple units subject to dependent failures, for example, a standby redundancy module, the module representation of the system has no dependent failures. In other words, failure of a module does not affect the operation of other modules in the system. This allows the path enumeration algorithm to be used for reliability calculation.

Extension to System MTBF Calculation

The most important parameter of the system reliability is the mean time between failures (MTBF). The system MTBF is defined as

$$\bar{M} = \int_0^{\infty} R(t) dt \quad (2.23)$$

The sensitivity of the MTBF to the reliability r_{ij} of the j th unit belonging to the i th module, and to the number of units in the i th module is given by

$$\frac{\partial \bar{M}}{\partial r_{ij}} = \int_0^{\infty} \frac{\partial R}{\partial R_i} \frac{\partial R_i}{\partial r_{ij}} dt \quad (2.24)$$

and

$$\frac{d\bar{M}}{dN_i} = \int_0^{\infty} \frac{\partial R}{\partial R_i} \frac{\partial R_i}{\partial N_i} dt \quad (2.25)$$

Expressions for $\frac{\partial R_i}{\partial r_{ij}}$ and $\frac{\partial R_i}{\partial N_i}$ are shown in Appendix A.

Therefore, by the addition of an integration routine to the reliability calculation program described earlier, the system MTBF and its sensitivities can be easily calculated. A computer program containing a numerical integration routine based on Simpson's rule has been developed.

The MTBF calculation routine developed here is more flexible and easier to apply to more complicated systems than the RELCOMP routine

of Fleming [22] which is only applicable to series-parallel configurations.

Approximations to System Reliability

Earlier, path and cut enumeration methods for system reliability calculation were discussed. The method adopted here was based on the path enumeration algorithm because an efficient path finding algorithm was readily available. The introduction of module structure allowed reliability evaluation of fairly complex networks.

If the number of paths is large (more than 10), the computation time increases rapidly. This is evident from Equation (2.11) which gives the total number of terms involved in the reliability expression as a function of total number of paths in the system. Same is true for reliability calculation based on the cut enumeration. To avoid this difficulty, some lower and upper bounds approximations to system reliability are presented below [4,28,34,37,46,52].

First we develop reliability approximations based on the path enumeration method. The series given by Equation (2.7) has the following properties [37,46,52]:

$$\begin{aligned}
 R &\leq R_{U_1} = \sum_{i=1}^M P_r \{P_i\} \\
 R &\geq R_{L_1} = \sum_{i=1}^M P_r \{P_i\} - \sum_{i=1}^M \sum_{j>i}^M P_r \{P_i \cap P_j\} \\
 R &\leq R_{U_2} = \sum_{i=1}^M P_r \{P_i\} - \sum_{i=1}^M \sum_{j>i}^M P_r \{P_i \cap P_j\} \\
 &\quad + \sum_{i=1}^M \sum_{j>i}^M \sum_{k>j}^M P_r \{P_i \cap P_j \cap P_k\} \\
 &\quad \vdots
 \end{aligned} \tag{2.26}$$

and

$$\begin{aligned} R_{U_1} &\geq R_{U_2} \geq \dots \\ R_{L_1} &\leq R_{L_2} \leq \dots \end{aligned} \quad (2.27)$$

Therefore, R_{U_1}, R_{U_2}, \dots can be used as successive upper bounds for R , and R_{L_1}, R_{L_2}, \dots can be used as lower bounds for R . This is the lower and upper bounds reliability approximation. These approximations are very close when element reliabilities are small. This is also called the low reliability region approximation.

In the analogous way, we can develop high reliability region approximation formulae from the cut enumeration method [28,37,46,52].

From Equation (2.9) we have

$$\begin{aligned} R &\geq R_{L_1} = 1 - \sum_{j=1}^M P_r \{ \bar{C}_j \} \\ R &\leq R_{U_1} = 1 - \sum_{j=1}^M P_r \{ \bar{C}_j \} + \sum_{j=1}^M \sum_{k>j}^M P_r \{ \bar{C}_j \bar{C}_k \} \\ &\vdots \end{aligned} \quad (2.28)$$

and

$$\begin{aligned} R_{U_1} &\geq R_{U_2} \geq \dots \\ R_{L_1} &\leq R_{L_2} \leq \dots \end{aligned} \quad (2.29)$$

From Equation (2.28) the number of terms in the lower and upper bound computations R_{L_1} and R_{U_1} are M and $M(M+1)/2$, respectively. This is compared to $2^M - 1$ terms obtained by expanding Equation (2.9).

In the system MTBF calculation both the approximations based on path enumeration and cut enumeration will play an important role because element reliabilities change from high reliability to low reliability region as time increases from 0 to ∞ .

Conclusions

The proposed algorithm is more powerful and efficient for deriving and evaluating the total system reliability, the sensitivities, and the MTBF of complex systems than other algorithms based on state enumeration method, since systems are often complex enough to have many units, but not so complicated as to have many paths in the module representation.

The procedures presented above have been programmed in FORTRAN IV and consist of eight basic subroutines. The nature of the MAIN program depends on the application and may call one or more of these subroutines. Appendix G contains a description of each subroutine, including a schematic diagram, and a listing of the corresponding FORTRAN program. In addition to these basic subroutines there are several others which are concerned with specific applications. These will be discussed in later chapters.

CHAPTER III
OPTIMAL DESIGN FOR RELIABILITY AND AVAILABILITY
OF COMPLEX SYSTEMS

Most of the earlier literature in the area of system reliability optimization consider only the series-parallel system [2,3,6,8,18-20,23, 24,30,31,33,36,38,40-44,48,51,54-56,61]. Such a system is shown in Figure 3.1 where each module has several identical components in parallel to provide redundancy, and the system fails if all the components in a module fail. The reason that this particular system has received so much attention is the separability of the reliability function. For the configuration shown in Figure 3.1 the system reliability is

$$\begin{aligned} R(\underline{n}, \underline{r}) &= \prod_{j=1}^N R_j(n_j, r_j) \\ &= \prod_{j=1}^N \left[1 - (1-r_j)^{n_j} \right] \end{aligned} \quad (3.1)$$

where r_j is the reliability of a component in the j th module and n_j is the total number of identical components in the j th module. Taking logarithms of both sides of Equation (3.1)

$$\begin{aligned} \ln R(\underline{n}, \underline{r}) &= \sum_{j=1}^N \ln R_j(n_j, r_j) \\ &= \sum_{j=1}^N \ln \left[1 - (1-r_j)^{n_j} \right] \end{aligned} \quad (3.2)$$

Since the natural logarithm is monotonic, maximization of the logarithm is equivalent to maximization of its argument. The form of Equation (3.2) for reliability maximization is more convenient to use since each term of the

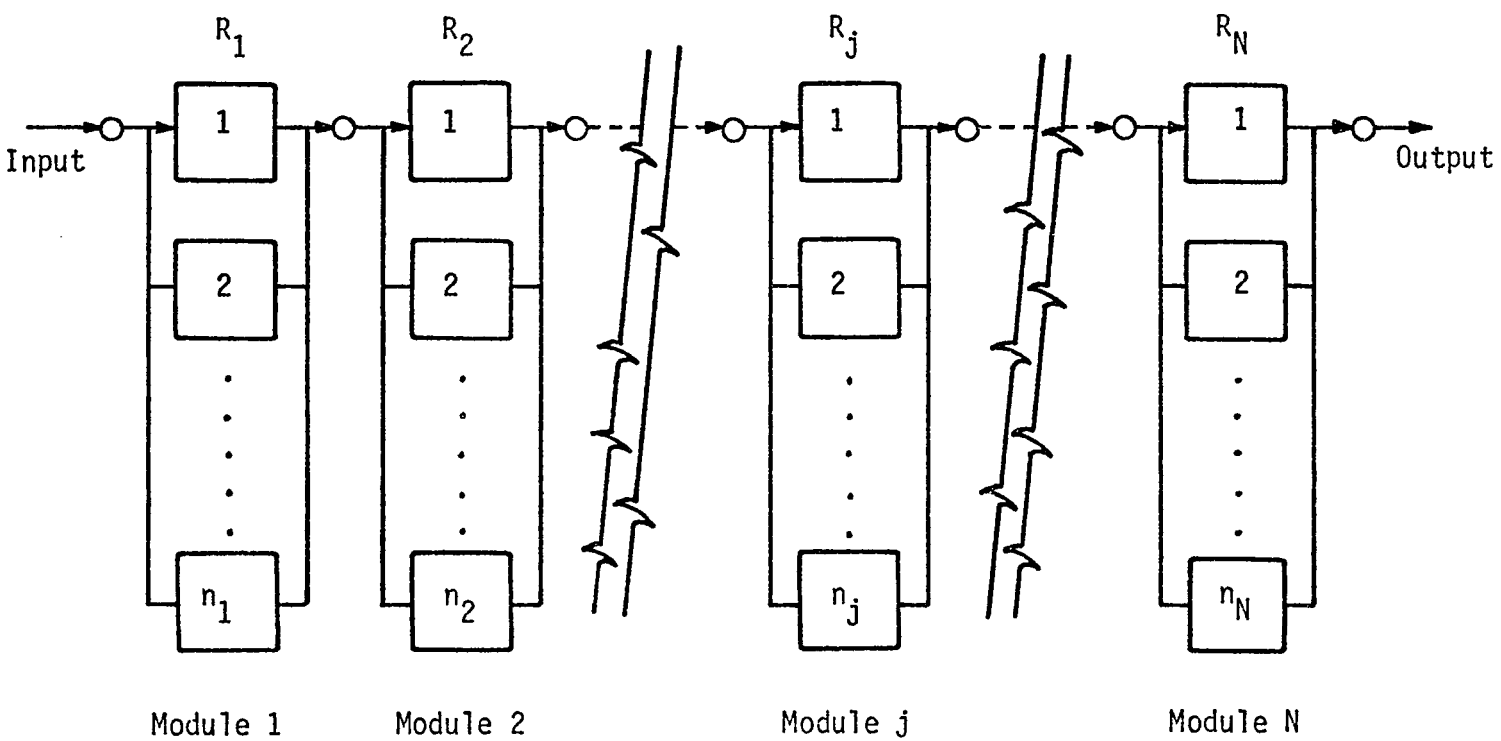


Figure 3.1
Series-Parallel System

sum depends on a single variable. A separable function can be analyzed as a multistage process to which the methods of dynamic programming [6,23,30,33,38,43,61] and discrete maximum principle [1955] are applicable. The problem may also be transformed and solved as an integer programming problem [24,41,44,54,56] or by using the Lagrangian multipliers technique [2,3,18,38,40]. The concept of dominating sequence (or a family of undominated allocations) has also been used to solve the problem [3,8,30,38,48].

Burton and Howard [11,12] consider the reliability optimization of a series-parallel-series system. They present a dynamic programming model. The notion of the generalized decomposition operator is used to develop a set of recursive relations. Terano et.al. [53] also consider a series-parallel-series system. They linearize the system at the nominal point and apply dynamic programming to find optimum redundancies. A non-series-parallel system is treated by Tillman et.al. [57]. They use the sequential unconstrained minimization technique to find optimum module reliabilities that maximize the system reliability.

In the previous chapter a general computer algorithm has been developed to calculate the reliability of a complex system having been given the reliability of each unit and the system configuration in the form of a reliability graph. The algorithm also calculates the sensitivity of the system reliability to individual system units.

In this chapter the sensitivity function is used for maximizing the reliability, availability or profit of a complex system (series-parallel-series or non-series-parallel) subject to linear or non-linear constraints. The optimum seeking algorithm which finds the optimum redundancies and preventive maintenance schedules is based on an integer gradient method of Reiter and

Rice (R-R) [50]. The proposed algorithm can be used to optimize a large class of reliability problems encountered in practice.

Three types of reliability optimization problems are formulated:

(1) Optimal Allocation of Redundancy

The problem is to find the optimal redundancy for each module (unit reliabilities are specified) so as to maximize the system reliability subject to linear or non-linear cost constraints and is stated as follows:

Maximize the system reliability

$$R(\underline{n}, \underline{r}) \quad (3.3)$$

where \underline{r} is specified.

Subject to linear or non-linear cost constraints

$$\sum_{i=1}^N g_{ij}(n_i, r_i) \leq G_j \quad j = 1, 2, \dots, m \quad (3.4)$$

and non-negative and integer constraints

$$\begin{aligned} n_i &\geq k_i & i &= 1, 2, \dots, N \\ n_i &: \text{integer} \end{aligned} \quad (3.5)$$

where k_i is the minimum number of units in the i th module that must operate for the module to function successfully.

This constitutes a non-linear integer programming problem.

(2) Optimal Allocation of Unit-Reliability or Unit Maintenance Interval

The problem is to find the optimal unit reliabilities or the optimum maintenance interval for each unit (module redundancies are specified) so as to maximize the system reliability subject to linear or non-linear cost constraints.

Once the optimal unit reliabilities are known the optimal maintenance interval can be evaluated by the relationship

$$r_i = e^{-T_i/(MTBF)_i} \quad (3.6)$$

where T_i is the maintenance interval of a unit belonging to the i th module, and $(MTBF)_i$ is the mean time between failure of a unit belonging to the i th module.

The optimization problem can be stated as follows

Maximize the system reliability

$$R(\underline{n}, \underline{r}) \quad (3.7)$$

where \underline{n} is specified.

Subject to linear or non-linear cost constraints

$$\sum_{i=1}^N g_{ij}(n_i, r_i) \leq G_j \quad j=1, 2, \dots, m \quad (3.8)$$

and non-negative constraints.

$$1 \geq r_i \geq 0 \quad i=1, 2, \dots, N \quad (3.9)$$

The problem can be extended to 1-out-of-n active redundancy module with dissimilar units.

The above problem constitutes a non-linear programming problem.

(3) Mixed Optimal Problem

A combination of problems (1) and (2) can be stated as follows:

Maximize the system reliability $R(\underline{n}, \underline{r})$ with respect to the number of redundant units (\underline{n}) and unit reliabilities (\underline{r}) subject to linear or non-linear cost constraints.

This is a non-linear mixed integer programming problem.

System Availability and System Profit

System availability can be improved by adding redundancy to the system and/or by performing preventive maintenance on the system according to some prescribed schedule [5,49].

In the absence of scheduled preventive maintenance, the system availability A is

$$A(\underline{n}) = \frac{MTBF}{MTBF + MDT} = \frac{\int_0^{\infty} R(\underline{n},t) dt}{\int_0^{\infty} R(\underline{n},t)dt + MDT} \quad (3.10)$$

where MTBF is the mean time to system failure, MDT is system mean downtime or mean repair time.

Now suppose preventive maintenance is performed on the system every "T" hours of continuing operation. If the system fails before "T" hours have elapsed, emergency maintenance is performed at that time. Preventive maintenance is then rescheduled. We assume that the system is as good as new after any type of maintenance, scheduled or emergency, and that the system either operates at full capacity or is down for maintenance. Let,

$f(\underline{n},t)$ = system failure time probability density function

T_E = mean time to perform emergency maintenance on the system

T_S = mean time to perform scheduled maintenance on the system.

Thus, mean time between system maintenance is given by:

$$\begin{aligned} MTBM &= T * R(\underline{n},t) + \int_0^T t f(\underline{n},t)dt \\ &= \int_0^T R(\underline{n},t)dt \end{aligned} \quad (3.11)$$

Also, assuming that T_E and T_S are independent of \underline{n} and T , the mean down time for the system is:

$$\begin{aligned} \text{MDT} &= T_E (1 - R(\underline{n}, T)) + T_S * R(\underline{n}, t) \\ &= T_E - (T_E - T_S) * R(\underline{n}, t) \end{aligned} \quad (3.12)$$

System availability is then defined as

$$A(\underline{n}, T) = \frac{\int_0^T R(\underline{n}, t) dt}{\int_0^T R(\underline{n}, T) dt + T_E - (T_E - T_S) R(\underline{n}, t)} \quad (3.13)$$

In many situations, system profit is closely related to system availability. To examine this relationship, we define the following terms:

T^+ = total life time of the system in hours

T_u = total time during which the system is operating at full capacity

$$T_u = A(\underline{n}, T) * T^+ \quad (3.14)$$

T_D = total time during which the system is down

$$T_D = (1 - A(\underline{n}, T)) T^+ \quad (3.15)$$

$$T^+ = T_u + T_D \quad (3.16)$$

C_{Fi} = fixed capital cost of a unit in the i th module

C_{TF} = total fixed capital investment

$$C_{TF} = \sum_{i=1}^N n_i C_{Fi} \quad (3.17)$$

C_M = system maintenance cost per hour of system downtime. This can be taken as some percentage, X , of fixed capital investment.

$$C_M = \left(\sum_{i=1}^N n_i C_{Fi} \right) \frac{X}{100} \quad (3.18)$$

C_{TM} = total maintenance costs for a system lifetime of T^+ hours

$$C_{TM} = C_M * T_D \quad (3.19)$$

C_I = net income per hour; the difference between the selling price of the product and all expenses other than the maintenance costs and capital costs.

C_{TI} = total net income during system lifetime of T^+ hours.

$$C_{TI} = C_I * T_u \quad (3.20)$$

We can now write a cash balance for the system lifetime period of T^+ hours.

Profit = net income - capital cost - maintenance cost

or

$$P = C_{TI} - C_{TF} - C_{TM}$$

or

$$P = C_I * T_u - \sum_{i=1}^N n_i C_{Fi} - C_M * T_D$$

or

$$P = C_I * A(\underline{n}, T) * T^+ - \sum_{i=1}^N n_i C_{Fi} - \left(\sum_{i=1}^N n_i C_{Fi} \right) * \frac{X}{100} (1 - A(\underline{n}, T)) T^+ \quad (3.21)$$

The equation (3.21) does not take into account the cost of capital, the time value of money, or any other of the widely used economic methods.

Two types of optimization problems based on system availability are now formulated:

(1) Availability Maximization

The problem can be stated as follows:

Maximize the system availability

$$A(\underline{n}, T) = \frac{\int_0^T R(\underline{n}, t) dt}{\int_0^T R(\underline{n}, t) dt + T_E - (T_E - T_S) R(\underline{n}, T)} \quad (3.22)$$

subject to the cost constraint

$$C_{TF} + C_{TM} \leq C_{A0}$$

or

$$\left(\sum_{i=1}^N n_i C_{Fi} \right) \left(1 + \frac{\chi}{100} (1 - A(\underline{n}, T) T^+) \right) \leq C_{A0} \quad (3.23)$$

and non-negative and integer constraints

$$\begin{aligned} n_i &\geq k_i & i &= 1, 2, \dots, N \\ n_i &: \text{integer} \\ T &\geq 0 \end{aligned} \quad (3.24)$$

Other constraints, if present, can also be included.

This is a non-linear mixed integer programming problem.

(2) Profit Maximization

This problem can be stated as follows:

Maximize the system profit over its lifetime of T^+ hours

$$P = C_I * A(\underline{n}, T) * T^+ - \left(\sum_{i=1}^N n_i C_{Fi} \right) \left(1 + (1 - A(\underline{n}, T)) \frac{\chi}{100} * T^+ \right) \quad (3.25)$$

subject to the cost constraint

$$C_{TF} \leq C_{P0}$$

or

$$\sum_{i=1}^N n_i C_{Fi} \leq C_{P0} \quad (3.26)$$

and non-negative and integer constraints

$$\begin{aligned} n_i &\geq k_i & i &= 1, 2, \dots, N \\ n_i &: \text{integer} \\ T &\geq 0 \end{aligned} \quad (3.27)$$

This is also a non-linear mixed integer programming problem.

Method of Solution

There are few existing methods that can solve non-linear integer and mixed integer programming problems. The two methods considered here are (1) a pseudo-Boolean programming method based on partial enumeration developed by Lawler and Bell [32]. The method is used to solve the non-linear integer programming problems formulated in Chapter IV. The details of the algorithm are discussed in Appendix F. (2) The other method proposed by Reiter and Rice (R-R) [50] is based on the modified integer gradients of the objective function. It can be extended to solve non-linear mixed integer programming problems. The method is effectively applicable to the problems formulated earlier since we can easily generate the gradients of the objective function or the sensitivities in our terminology.

The R-R algorithm is used to solve the system optimization problems formulated in this chapter. The solution procedure has three phases:

Phase 1: Setting up of an initial point \underline{X}_I using uniformly distributed random numbers.

Phase 2: Searching for a feasible solution \underline{X}^+ of the problem, starting with \underline{X}_I . A search by the method of weighted perpendicular is performed. If the feasible solution is not obtained in a given number of iterations, the search is terminated, and we go back to Phase 1 of the search algorithm.

Phase 3: Locating a point \underline{X}_p which maximizes the objective function in the feasible region round \underline{X}^+ . This involves calculating the modified gradients of the objective function which define the path on which we search for improved values of the objective function.

\underline{X}_I , \underline{X}^+ , and \underline{X}_p denote vectors of variables, real and/or integer.

Details of the search algorithm are discussed in Appendix D. The module reliability and sensitivity expressions are given in Appendix A, and the sensitivity expressions for system availability and system profit are shown in Appendix E.

The R-R solution procedure provides locally maximal points. Repeated use of the algorithm, with judicious improvement in the lower and upper bounds on the system variables based on the previous results, gives better results, and, hopefully, the global optimum. The method is illustrated by the following examples:

Illustrative Examples

Solutions are now obtained for some of the optimization problems formulated earlier in this chapter. Two different module reliability graphs are considered. The module reliability graph of Figure 3.2 shows a series-parallel-series system and that of Figure 3.3 represents a non-series-parallel system. Henceforth, these will be referred to as System 1 and System 2, respectively. For the sake of illustration it is assumed that

all modules have 1-out-of- n_i active redundancy in the following examples:

Example (1) Optimal allocation of redundancy to maximize the system reliability.

Tables 3.1 and 3.3 give the unit reliability, unit, cost, and maximum number of units (upper bound in n_i) for each module of System 1 and System 2, respectively. These upper bounds on n_i 's were obtained from an initial analysis of the problem. There is a linear cost constraint of the type

$$\sum_{i=1}^N C_i n_i \leq C \quad (3.28)$$

where C is the total allowable cost. C_i is the cost per unit of the i th module. The allowable cost C is given in Table 3.1 and Table 3.3 for System 1 and System 2, respectively.

The problem then is to maximize the system reliability subject to the cost constraint of Equation (3.28).

Discussion of Results: Some of the locally optimal results obtained are tabulated in Tables 3.2 and 3.4 and correspond to System 1 and System 2, respectively. The problem was also solved for comparison purposes using the partial enumeration algorithm of Lawler and Bell which finds a unique optimal solution or the global optimum. It was observed that the R-R algorithm not only found locally maximal points but the best solution obtained, Result 10 in Tables 3.2 and 3.4, was also the global optimum.

Example (2) Optimal allocation of redundancy and maintenance interval to maximize the system availability or the system profit.

Consider the module reliability graph of System 1. Assume that units in each module have exponential failure pdf. Table 3.5 gives the capital cost, failure rate data, and upper bounds on n_i for identical units of each module.

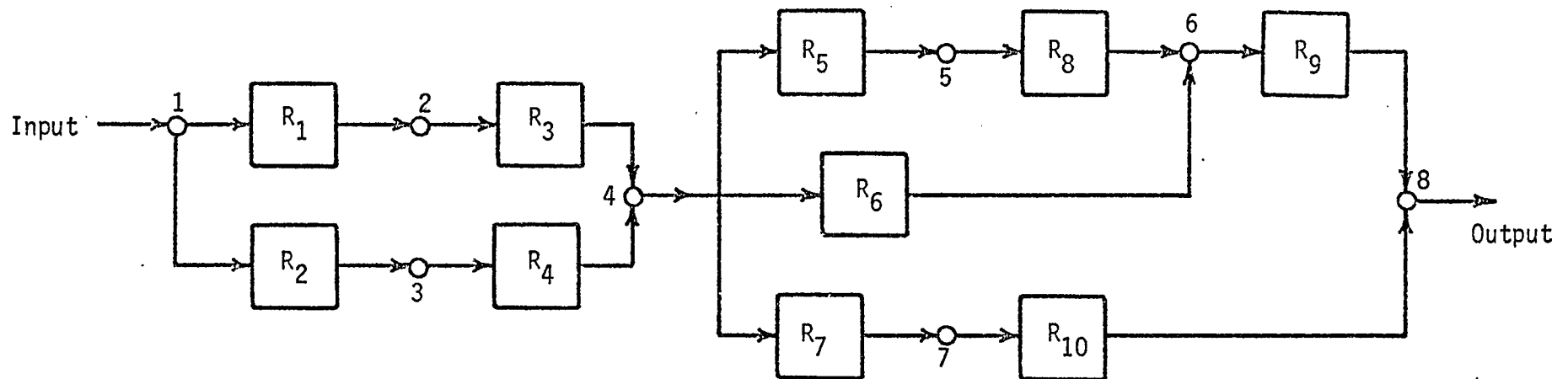


Figure 3.2.
Reliability Graph - System 1

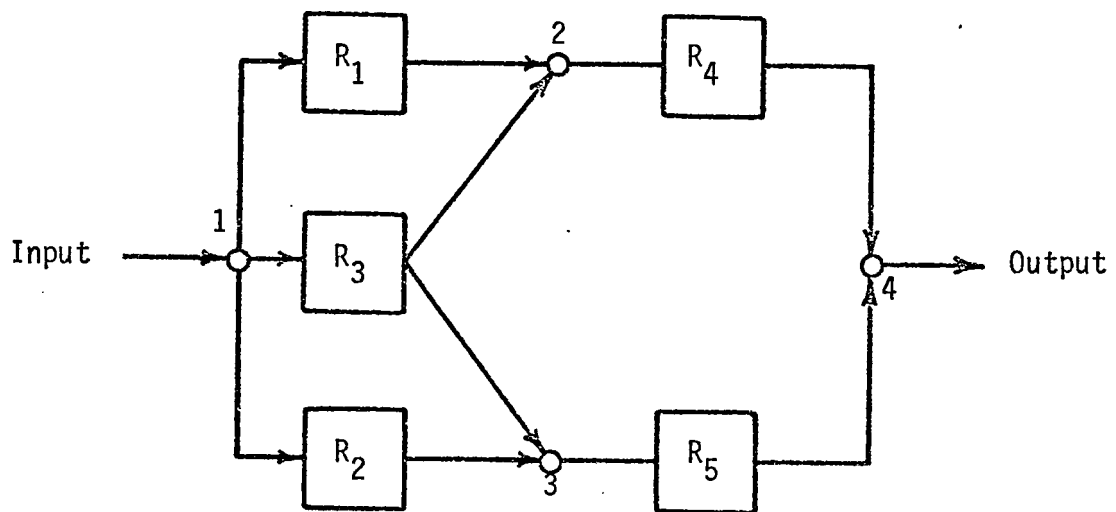


Figure 3.3
Reliability Graph - System 2

MODULE i	UNIT COST c_i	UNIT RELIABILITY r_i	UPPER BOUND $n_i(\max)$
1	75.00	0.75	4
2	90.00	0.90	4
3	80.00	0.80	4
4	87.00	0.87	4
5	72.00	0.72	4
6	88.00	0.88	4
7	75.00	0.75	4
8	89.00	0.89	4
9	75.00	0.75	4
10	80.00	0.80	4

Allowable Cost = 1300.00

Table 3.1
Reliability Optimization Data - System 1

MODULE	Optimum Number of Units in Each Module									
	Result 1	Result 2	Result 3	Result 4	Result 5	Result 6	Result 7	Result 8	Result 9	Result 10
1	2	2	2	2	1	2	1	1	1	1
2	2	2	1	2	2	1	2	2	2	2
3	2	2	2	1	1	2	1	1	2	1
4	1	1	2	2	2	2	2	2	2	2
5	1	1	1	1	2	1	2	2	1	1
6	1	2	1	1	1	2	1	2	1	1
7	2	1	2	2	1	1	2	1	2	3
8	1	1	1	1	2	1	1	1	1	1
9	2	3	2	3	3	3	2	3	2	2
10	2	1	2	1	1	1	2	1	2	2
COST UTILIZED	1286.0	1294.0	1283.0	1288.0	1299.0	1291.0	1290.0	1298.0	1298.0	1293.0
MAXIMUM RELIABILITY	0.975982	0.977958	0.978306	0.978900	0.978938	0.980288	0.981123	0.982116	0.982301	0.983708

Table 3.2
Reliability Optimization - System 1

MODULE i	UNIT COST C_i	UNIT RELIABILITY r_i	UPPER BOUND n_i
1	7000.	0.90	3
2	3000.	0.70	4
3	5000.	0.70	4
4	3000.	0.70	6
5	7000.	0.90	4

Allowable Cost = 50,000.00

Table 3.3
Reliability Optimization Data - System 2

MODULE	Optimum Number of Units in Each Module									
	Result 1	Result 2	Result 3	Result 4	Result 5	Result 6	Result 7	Result 8	Result 9	Result 10
1	2	1	1	1	1	1	1	1	2	2
2	2	3	3	3	1	2	2	3	1	3
3	1	2	3	1	3	2	3	2	2	1
4	3	1	4	5	6	4	5	3	5	5
5	2	3	1	2	1	2	1	2	1	1
COST UTILIZED	48,000	50,000	50,000	50,000	50,000	49,000	50,000	49,000	49,000	50,000
MAXIMUM RELIA- BILITY	0.998717	0.998792	0.998848	0.998856	0.998919	0.998966	0.999262	0.999340	0.999364	0.999364

Table 3.4

Reliability Optimization - System 2

Other data are:

$$C_I = \$10.00/\text{hour}$$

$$X = 0.02\%$$

$$T^+ = 80,000 \text{ hours}$$

$$T_E = 400 \text{ hours}$$

$$T_S = 100 \text{ hours}$$

$$C_{AO} = \$350,000$$

$$C_{PO} = \$200,000$$

Discussion of Results: The results of the ten randomly sampled points for the system availability maximization and the system profit maximization are tabulated in Tables 3.6 and 3.7, respectively. Also, the solutions obtained at each step in the optimization procedure are shown in Tables 3.8 and 3.9 for the Result 1 of each problem. The results show that in the case of system availability maximization the local optima are dependant on the starting point as far as n_i 's are concerned. This is due to the system availability being a monotonically increasing function with respect to the n_i 's, the gradient always points forward increasing n_i 's, and thus, the solution lies on or near the constraint boundary. The system profit maximization problem, on the other hand, gives fewer local optima and the solution does not necessarily lie on the constraint boundary. The method is thus seen to be more effective in the case of a problem with non-monotonically increasing objective function and/or the optimum solution being in the interior of the constraint boundary.

The computer time to find a single locally maximal point is about one minute on the IBM 360 and fifteen seconds on the UNIVAC 1108.

MODULE i	UPPER BOUND $n_{i(\max)}$	CAPITAL COST/UNIT (C_{Fi}) \$/UNIT	FAILURE RATE (λ_i) HOUR ⁻¹
1	6	5000.00	$3.0 * 10^{-4}$
2	6	15000.00	$1.0 * 10^{-4}$
3	6	10000.00	$2.2 * 10^{-4}$
4	6	15000.00	$1.2 * 10^{-4}$
5	6	5000.00	$3.2 * 10^{-4}$
6	6	15000.00	$1.1 * 10^{-4}$
7	6	5000.00	$3.0 * 10^{-4}$
8	6	15000.00	$1.0 * 10^{-4}$
9	12	5000.00	$3.0 * 10^{-4}$
10	6	10000.00	$2.2 * 10^{-4}$

Table 3.5

Availability or Profit Maximization Data - System 1

MODULE NUMBER	OPTIMUM NUMBER OF UNITS IN EACH MODULE									
	Result 1	Result 2	Result 3	Result 4	Result 5	Result 6	Result 7	Result 8	Result 9	Result 10
1	5	1	2	2	1	2	1	1	3	1
2	2	3	3	2	2	2	2	3	2	3
3	2	1	1	1	1	1	1	1	1	1
4	2	3	3	3	3	3	3	3	2	3
5	1	2	1	2	1	3	1	1	1	1
6	2	2	2	2	2	1	2	2	2	2
7	2	1	2	2	4	4	3	2	4	3
8	1	1	1	1	1	1	1	1	1	1
9	6	8	6	8	5	4	8	7	7	6
10	1	1	1	1	2	2	1	1	1	1
Fixed Capital Cost (\$)	205,000	215,000	210,000	210,000	205,000	200,000	205,000	210,000	200,000	210,000
Total Maintenance Cost (\$)	140,822	134,648	138,428	139,485	142,286	146,922	141,368	135,422	145,765	137,502
Optimum Maintenance Interval (Hours)	3,190	3,575	3,375	3,490	3,380	3,030	3,510	3,490	3,260	3,430
Optimized System Availability	0.957066	0.960858	0.958801	0.958487	0.956620	0.954087	0.956900	0.959696	0.954448	0.959077
System Profit (\$)	419,831	419,038	418,613	417,304	418,010	416,347	419,153	422,334	417,794	419,759

Table 3.6
Availability Optimization - System 1

MODULE NUMBER	OPTIMUM NUMBER OF UNITS IN EACH MODULE									
	Result 1	Result 2	Result 3	Result 4	Result 5	Result 6	Result 7	Result 8	Result 9	Result 10
1	3	3	1	3	3	3	1	1	3	3
2	1	1	2	1	1	1	2	2	1	1
3	2	2	1	2	2	2	1	1	2	2
4	1	1	2	1	1	1	2	2	1	1
5	1	1	1	1	1	1	1	2	1	2
6	1	1	2	1	1	1	1	1	1	1
7	2	2	1	2	2	2	3	2	2	1
8	1	1	1	1	1	1	1	1	1	1
9	3	3	5	3	3	4	3	4	4	4
10	1	1	1	1	1	1	2	1	1	1
Fixed Capital Cost (\$)	135,000	135,000	165,000	135,000	135,000	140,000	160,000	155,000	140,000	140,000
Total Maintenance Cost (\$)	155,112	155,138	143,377	155,116	155,130	153,017	148,688	150,276	153,045	152,779
Optimum Maintenance Interval (Hours)	2,145	2,170	2,840	2,150	2,170	2,240	2,560	2,580	2,290	2,190
Optimized System Profit (\$)	452,439	452,404	448,176	452,433	452,415	452,334	444,847	446,248	452,297	452,657
Availability	.928189	.928177	.945691	.928187	.928181	.931689	.941919	.939405	.931676	.931795

Table 3.7
Profit Optimization - System 1

MODULE NUMBER	1	2	3	4	5	6	7	8	9	10	MAINTENANCE INTERVAL (Hours)	COST CONSTRAINT (\$)	SYSTEM AVAILABILITY (\$)
Initial Point (1)	5	2	6	6	1	3	2	3	10	4	2131	+338,723	--
Feasible Point (2)	5	1	1	2	1	1	1	1	2	1	2484	- 11,563	0.911412
(3)	5	1	1	2	1	1	1	1	3	1	2481	- 24,423	0.922165
(4)	5	1	1	2	1	2	1	1	3	1	2478	- 14,032	0.931263
(5)	5	2	1	2	1	2	1	1	3	1	2478	- 6,748	0.939910
(6)	5	2	1	2	1	2	1	1	4	1	2480	- 13,996	0.945832
(7)	5	2	2	2	1	2	1	1	4	1	2490	- 4,856	0.948966
(8)	5	2	2	2	1	2	1	1	5	1	2503	- 5,588	0.952112
(9)	5	2	2	2	1	2	1	1	6	1	2539	- 2,652	0.953954
(10)	5	2	2	2	1	2	1	1	6	1	2639	- 4,518	0.954537
(11)	5	2	2	2	1	2	2	1	6	1	2690	- 271	0.955875
(12)	5	2	2	2	1	2	2	1	6	1	2790	- 1,734	0.956321
(13)	5	2	2	2	1	2	2	1	6	1	2890	- 2,824	0.956654
(14)	5	2	2	2	1	2	2	1	6	1	2990	- 3,574	0.956882
(15)	5	2	2	2	1	2	2	1	6	1	3090	- 4,016	0.957017
Locally Optimum Point (16)	5	2	2	2	1	2	2	1	6	1	3190	- 4,178	0.957066

Table 3.8
Steps to Reach Locally Optimal Solution - Result 1 of Table 3.6

MODULE NUMBER	1	2	3	4	5	6	7	8	9	10	MAINTENANCE INTERVAL (Hours)	COST CONSTRAINT (\$)	SYSTEM PROFIT
Initial Point (1)	5	2	6	6	1	5	3	5	10	6	1131	+285,000	--
Feasible Point (2)	3	1	2	1	1	1	1	1	8	2	1131	- 35,000	330,717
(3)	3	1	2	1	1	1	1	1	8	1	1211	- 45,000	363,167
(4)	3	1	1	1	1	1	1	1	8	1	1298	- 55,000	366,674
(5)	3	1	1	1	1	1	1	1	7	1	1380	- 60,000	386,472
(6)	3	1	1	1	1	1	1	1	6	1	1448	- 65,000	403,457
(7)	3	1	1	1	1	1	1	1	5	1	1507	- 70,000	418,106
(8)	3	1	1	1	1	1	1	1	4	1	1558	- 75,000	429,769
(9)	3	1	1	1	1	1	1	1	3	1	1606	- 80,000	435,455
(10)	2	1	1	1	1	1	1	1	3	1	1640	- 85,000	439,400
(11)	2	1	2	1	1	1	1	1	3	1	1771	- 75,000	445,012
(12)	3	1	2	1	1	1	1	1	3	1	1870	- 70,000	448,389
(13)	3	1	2	1	1	1	2	1	3	1	1945	- 65,000	451,396
(14)	3	1	2	1	1	1	2	1	3	1	2045	- 65,000	452,222
Locally Optimum Point (15)	3	1	2	1	1	1	2	1	3	1	2145	- 65,000	452,439

Table 3.9

Steps to Reach Locally Optimal Solution - Result 1 of Table 3.7

Conclusions

The efficient computer algorithm developed in Chapter II to calculate the reliability and the sensitivities of complex systems is used in conjunction with the R-R modified integer gradient method for system optimization. The proposed method is easily applicable to series-parallel-series or non-series-parallel systems and can handle multiple linear and/or non-linear constraints. Although this method does not insure a global optimum, it does find various near-optimum solutions. From a practical consideration, this could provide for a wider choice during the design phase.

The method presented above has been programmed in FORTRAN IV. In addition to the eight basic subroutines, there are seven others which are used during system optimization. Appendix G contains a description of each subroutine, including a listing of the corresponding FORTRAN program and a schematic diagram of the main program.

CHAPTER IV

OPTIMAL RELIABILITY DESIGN OF PROCESS SYSTEMS

In the design of electrical, electronic and safety systems, the flow of information is the most important factor, whereas in process systems, the flow of materials is also very important. In this chapter an optimal reliability design of process systems is proposed which takes into account the quantitative aspect of systems throughput or capacity. It is based on the k -out-of- n configuration instead of the conventional parallel 1-out-of- n configuration. The problem is so formulated that determining the optimum k -out-of- n configuration also determines the optimum capacity of units in each stage of the system.

Basic Idea

Consider a pump system which must deliver at least L gallons per second. If we install one pump of capacity L and cost C_L , the reliability of the system is that of the single pump. If we appropriately install N pumps of capacity L in parallel, the reliability will increase but so will the cost which will go up by a factor of N as shown in Figure 4.1. If the allowed cost is more than C_L but less than $2C_L$, the reliability cannot be increased by the conventional parallel redundancy configuration.

On the other hand, if we install three pumps each of capacity $L/2$, and assuming for the moment that the cost is proportional to capacity, this can be achieved at a cost of $1.5 C_L$ which is less than the allowed cost of $2C_L$. The system can now operate if at least any two out of the three pumps function as shown in Figure 4.2. This gives rise to the well-known k -out-of- n configuration, and the system has a higher reliability compared to a single pump of capacity L . Here it is assumed that

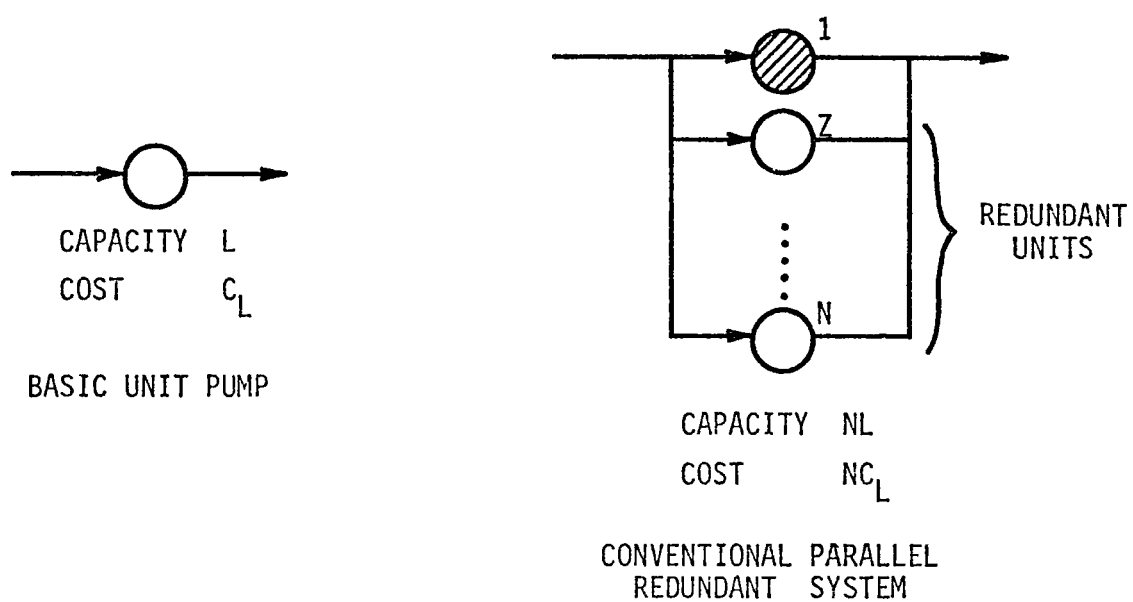


Figure 4.1
Conventional Parallel Configuration

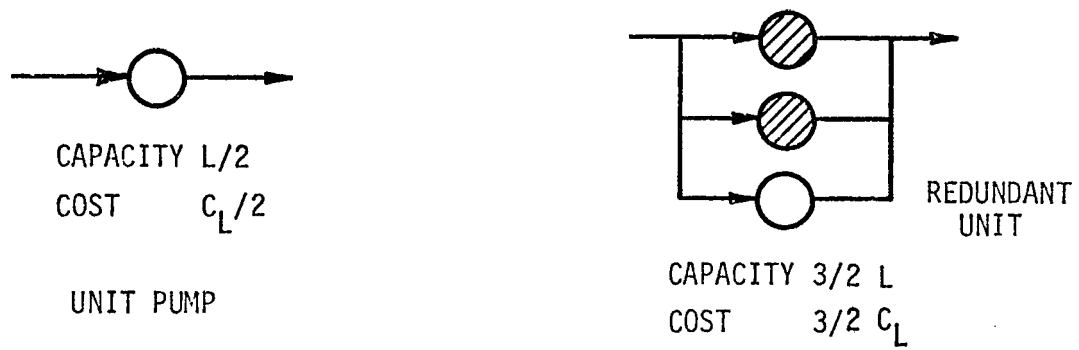


Figure 4.2
2-Out-Of-3 Configuration

- (1) the reliability of a unit is independent of its capacity; and,
- (2) the pump system has perfect valving.

There are many system configurations of this type which cost less than $2 C_L$. As calculated from Equation (4.1) the reliabilities of some of these configurations are listed in Table 4.1. One can clearly see the possibility of increasing the system reliability by this means. It is important to note that changes in k affect both the reliability of the system and the capacity of units in the system.

Problem Formulation

Although the basic idea can be extended to complex systems, the attention is first restricted to the optimal design of a single stage. Consider a system which operates if at least k out of n units function. This system configuration is called a k -out-of- n configuration. The two types of configurations considered are: (1) an active k -out-of- n configuration, and (2) a standby k -out-of- n configuration.

1. Active k -out-of- n Configuration

All the n identical units in parallel are operating independent of each other. The probability that at least k -out-of- n units are operating is given by (see Appendix A):

$$R_A(n, k; p) = \sum_{i=k}^n \binom{n}{i} p^i (1-p)^{n-i} \quad (4.1)$$

where p is the reliability of a single unit.

2. Standby k -out-of- n Configuration

If k units are functioning while the other $(n-k)$ units in standby have zero failure probability, the system is called a cold standby k -out-of- n configuration. Assuming that all units are identical and have exponential

		SYSTEM RELIABILITY (ACTIVE)				
Configuration		3-out-of-4	2-out-of-3	4-out-of-6	3-out-of-5	4-out-of-7
Cost		1.33 C_L	1.5 C_L	1.5 C_L	1.67 C_L	1.75 C_L
UNIT RELIABILITY	0.95	0.9860	0.9928	0.9978	0.9988	0.9998
	0.85	0.8905	0.9392	0.9527	0.9734	0.9879
	0.75	0.7383	0.8438	0.8306	0.8965	0.9294

Cost is proportional to capacity; valving is perfect;
reliability of a unit is independent of its capacity.

Table 4.1

System reliability comparison for several
k-out-of-n configurations

failure probability density function, the reliability expression is (see Appendix A):

$$R_C(n, k; p) = \sum_{j=0}^{n-k} p^k \frac{(-k \ln p)^j}{j!} \quad (4.2)$$

where $p = e^{-\lambda t}$ is the probability of success in service of identical units with constant hazard rate λ for an operating time t .

If the units in standby have the same failure probability as those in service, it is called a hot standby system. The reliability expression for a hot standby k -out-of- n configuration is the same as Equation (4.1) of the active k -out-of- n configuration; that is,

$$R_H(n, k; p) = R_A(n, k; p) \quad (4.3)$$

It is assumed that any one of the three configurations is allowed as a subsystem of the total system. The use of a particular configuration depends on the properties and the purpose of the subsystem.

System Cost

The assumption that the cost is proportional to the capacity is somewhat unrealistic. The cost-capacity relationship used here is given by

$$\text{cost} = K (\text{capacity})^s \quad (4.4)$$

where K and s are positive constants and $0 < s < 1$.

If k units are necessary to support the minimum required capacity L , then one unit must support the capacity L/k . Therefore, the total cost of the system which consists of such n units is:

$$\text{cost} = n \cdot K \cdot \left(\frac{L}{k} \right)^s \quad (4.5)$$

or

$$\text{cost} = n \cdot k_c \cdot k^{-S} \quad (4.6)$$

where k_c is the cost of a unit with capacity L .

System Optimization

The two types of problems considered are (1) reliability maximization, and, (2) cost minimization.

1. Reliability Maximization

If the amount $1 \cdot k_c$ can be spent to improve the reliability of the system, the cost constraint is

$$n \cdot k_c \cdot k^{-S} \leq 1 \cdot k_c \quad (4.7)$$

or

$$1 - n \cdot k^{-S} \geq 0 \quad (4.8)$$

where 1 is a positive real number greater than 1.

The optimization problem can then be stated as follows:

Select n and k to maximize the system reliability given in either Equation (4.1), (4.2) or (4.3) subject to the cost constraint given in Equation (4.8) and the inherent constraints

$$n - k \geq 0 \quad (4.9)$$

$$n \text{ and } k: \text{ integers not less than } 1 \quad (4.10)$$

This is a nonlinear integer programming problem with a nonlinear objective function and two constraints, a nonlinear and a linear.

2. Cost Minimization

When it is necessary to minimize the cost under the condition that the system reliability is not less than a preassigned value, say R^* , the optimization problem can then be stated as follows:

Select n and k to minimize the system cost given in Equation (4.6) subject to the constraint that the reliability given in either Equation (4.1), (4.2) or (4.3) is not less than a preassigned value R^* and also satisfies the inherent constraints given in Equations (4.9) and (4.10).

This is again a nonlinear integer programming problem with a nonlinear objective function and two constraints, a nonlinear and a linear.

In some cases in addition to the cost constraint in the reliability maximization problem or the reliability constraint in the cost minimization problem, nonlinear weight and/or volume constraints may be present. A weight constraint may be of the form

$$W = n \cdot k_w \cdot k^{-s_w} \leq W^* \quad (4.11)$$

and a volume constraint may be

$$V = n \cdot k_v \cdot k^{-s_v} \leq V^* \quad (4.12)$$

where k_w and k_v are constants corresponding to the standard weight and volume of a single unit of capacity L . s_w and s_v are non-negative constants less than one. W^* and V^* are maximum allowable weight and volume, respectively.

A general problem of this type is, therefore, a nonlinear integer programming problem with a nonlinear objective function, several nonlinear constraints, and a linear constraint.

Extension to Series Parallel and Complex Systems

Consider a series-parallel system with N stages as shown in Figure 4.3. The i th stage is assumed to be a k_i -out-of- n_i configuration in either active, cold standby, or hot standby condition.

The total reliability of the system is expressed by

$$R_T(\underline{n}, \underline{k}; \underline{p}) = \prod_{i=1}^N \left[R_A(n_i, k_i; p_i), R_C(n_i, k_i; p_i) \text{ or } R_H(n_i, k_i; p_i) \right] \quad (4.13)$$

where p_i is the unit reliability of the i th stage.

The total system cost is

$$\text{Total cost} = \sum_{i=1}^N n_i \cdot k_{ci} \cdot k_i^{-s_i} \quad (4.14)$$

Similarly, the total weight and volume are

$$\text{Total weight} = \sum_{i=1}^N n_i \cdot k_{wi} \cdot k_i^{-s_{wi}} \quad (4.15)$$

and

$$\text{Total volume} = \sum_{i=1}^N n_i \cdot k_{vi} \cdot k_i^{-s_{vi}} \quad (4.16)$$

When the system to be considered is a complex system, as shown in Figure 4.4, and has several paths between input and output, the total reliability is (see Chapter II):

$$\begin{aligned} R_T(\underline{n}, \underline{k}; \underline{p}) &= \sum_{i=1}^M \prod_{1 \in P_i} (R_{A1}, R_{C1}, \text{ or } R_{H1}) \\ &- \sum_{i=1}^M \sum_{j>i}^M \prod_{1 \in P_i \cup P_j} (R_{A1}, R_{C1}, \text{ or } R_{H1}) + \dots \\ &+ (-1)^{M-1} \prod_{1 \in \bigcup_{i=1}^M P_i} (R_{A1}, R_{C1}, \text{ or } R_{H1}) \end{aligned} \quad (4.17)$$

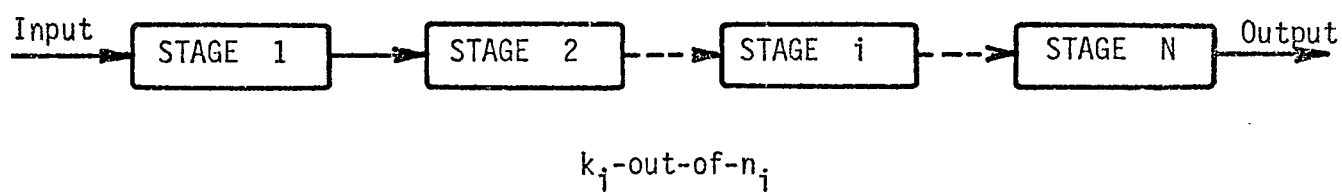


Figure 4.3
Series-Parallel System with N Stages

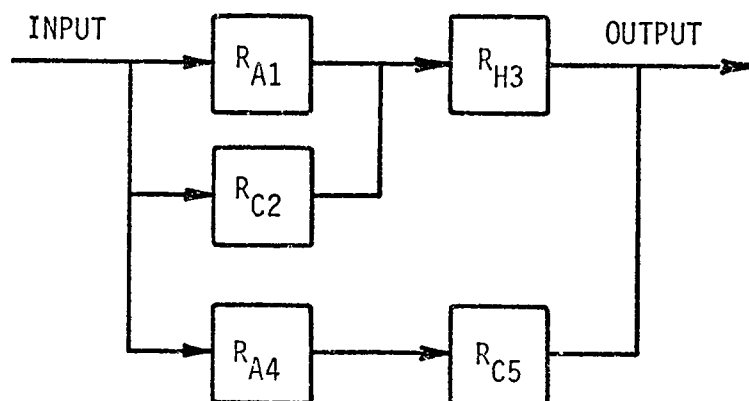


Figure 4.4
Example of a Complex System

where P_i is a path, M is the total number of paths, and U represents path unions.

In the case of Figure 4.4, we have

$$\begin{aligned}
 R_T(\underline{n}, \underline{k}; \underline{p}) = & R_{A1} \cdot R_{H3} + R_{C2} \cdot R_{H3} + R_{A4} \cdot R_{C5} \\
 - & (R_{A1} \cdot R_{C2} \cdot R_{H3} + R_{A1} \cdot R_{H3} \cdot R_{A4} \cdot R_{C5} + R_{C2} \cdot R_{H3} \cdot R_{A4} \cdot R_{C5}) \\
 & + R_{A1} \cdot R_{C2} \cdot R_{H3} \cdot R_{A4} \cdot R_{C5}
 \end{aligned} \quad (4.18)$$

Since total cost, weight, and volume for the complex system can also be described by Equations (4.14), (4.15) and (4.16), respectively, it is easy to construct a reliability maximization or a cost minimization problem for a complex system. The computer algorithm which can calculate the reliability expression from a given reliability graph has been described in Chapter II.

Solution Method

Two methods have been considered for solving the non-linear integer programming problems. One is the integer gradient method proposed by Reiter and Rice [50] and used to solve the system optimization problems formulated in Chapter III. The other is the pseudo-Boolean programming method based on partial enumeration developed by Lawler and Bell (L-B) [32].

The L-B algorithm is used to solve the system optimization problem formulated in this chapter. It is applicable to any problem that can be put into the form

Minimize

$$g_0(\underline{x}) \quad (4.19)$$

subject to

$$g_{i1}(\underline{x}) - g_{i2}(\underline{x}) \geq 0 \quad i = 1, 2, \dots, m \quad (4.20)$$

where

$$\underline{x} = (x_1, x_2, \dots, x_n) \quad (4.21)$$

and

$$x_j = 0 \text{ or } 1 \quad j = 1, 2, \dots, n \quad (4.22)$$

and, where the restriction is applied, each of the functions g_0, g_{i1}, g_{i2} ($i = 1, 2, \dots, m$) is monotone non-decreasing in each of the variables x_1, x_2, \dots, x_n .

Consider the single stage reliability maximization problem. It is noted from physical reasoning that the system reliability to be maximized, R_A, R_H , or R_C , is monotone increasing with respect to n and monotone decreasing with respect to k . Mathematical proof of this property is shown below for R_A or R_H . By use of the formula

$$\binom{n+1}{k} = \binom{n}{k-1} + \binom{n}{k} \quad k \leq n \quad (4.23)$$

we have

$$R_A(n+1, k; p) - R_A(n, k; p) = \binom{n}{k-1} p^k (1-p)^{n+1-k} \geq 0 \quad (4.24)$$

Also,

$$R_A(n, k+1; p) - R_A(n, k; p) = - \binom{n}{k} p^k (1-p)^{n-k} \leq 0 \quad (4.25)$$

The objective function to be minimized, $g_0 = -R_A, -R_H$, or $-R_C$, is thus monotone decreasing with respect to n and monotone nondecreasing with respect to k . By the simple transformation of variable, $n = n_{\max} - n'$

or $n' = n_{\max} - n$, where n_{\max} is an upper bound on n , we can make $g_0(n', k)$ monotone nondecreasing in both n' and k .

Non-negative integer variables can be transformed into binary variables by means of the substitution

$$n = 1 + x_{n1} + 2x_{n2} + 2^2x_{n3} + \dots + 2^{j-1}x_{nj} \quad (4.26)$$

$$k = 1 + x_{k1} + 2x_{k2} + 2^2x_{k3} + \dots + 2^{j-1}x_{kj} \quad (4.27)$$

where x_{ni} and x_{ki} are binary variables and j is chosen to be sufficiently large for 2^j to be an upper bound on the value of n .

Applying the substitution

$$x_{ni} = 1 - x'_{ni} \quad i = 1, 2, \dots, j \quad (4.28)$$

where x'_{ni} is a binary variable, we have

$$n = 1 + (1 - x'_{n1}) + 2(1 - x'_{n2}) + 2^2(1 - x'_{n3}) + \dots + 2^{j-1}(1 - x'_{nj})$$

or

$$n = 2^j - x'_{n1} - 2x'_{n2} - 2^2x'_{n3} - \dots - 2^{j-1}x'_{nj} \quad (4.29)$$

Using Equations (4.27) and (4.29), the objective function $g_0(n, k)$ can be expressed in terms of x'_{ni} 's and x_{ki} 's, and it is obviously monotone decreasing in each of the variables x'_{ni} and x_{ki} ($i = 1, 2, \dots, j$).

The cost constraint of Equation (4.8) can then be rewritten as

$$1 - \frac{2^j - x'_{n1} - 2x'_{n2} - 2^2x'_{n3} - \dots - 2^{j-1}x'_{nj}}{(1 + x_{k1} + 2x_{k2} + 2^2x_{k3} + \dots + 2^{j-1}x_{kj})^s} \geq 0 \quad (4.30)$$

Comparing with Equation (4.20) we have

$$g_{11}(\underline{x}) = 1 - \frac{2^j - x'_{n1} - 2x'_{n2} - 2^2x'_{n3} - \dots - 2^{j-1}x'_{nj}}{(1 + x_{k1} + 2x_{k2} + 2^2x_{k3} + \dots + 2^{j-1}x_{kj})^s} \quad (4.31)$$

and

$$g_{12}(\underline{x}) = 0 \quad (4.32)$$

The volume and weight constraints can be converted into similar forms.

The linear constraint of Equation (4.9) can be written as

$$\begin{aligned} & (2^j - x'_{n1} - 2x'_{n2} - \dots - 2^{j-1}x'_{nj}) \\ & - (1 + x_{k1} + 2x_{k2} + \dots + 2^{j-1}x_{kj}) \geq 0 \end{aligned} \quad (4.33)$$

Again comparing with Equation (4.20), we have

$$g_{21}(\underline{x}) = 0 \quad (4.34)$$

and

$$\begin{aligned} g_{22}(\underline{x}) = & (x'_{n1} + 2x'_{n2} + 2^2x'_{n3} + \dots + 2^{j-1}x'_{nj} - 2^j) \\ & + (1 + x_{k1} + 2x_{k2} + 2^2x_{k3} + \dots + 2^{j-1}x_{kj}) \end{aligned} \quad (4.35)$$

Thus, the problem has been converted to a standard form to which the L-B algorithm is applicable. It is easily seen that the cost minimization problem can also be converted to the standard form. Similarly, complex multistage problems can be transformed and solved by the L-B algorithm.

The important features of the L-B algorithm are:

- (1) the true optimum solution found;
- (2) the extreme ease of programming;
- (3) the almost complete absence of housekeeping operations (no lists, no pushdown storage routines);
- (4) the very small amount of storage required; and,
- (5) the wide applicability to fairly general nonlinear integer programming problems.

A major limitation of the L-B algorithm is the computational infeasibility when the number of binary variables is large, approximately more than 30. Details of the L-B algorithm are explained in Appendix F.

Illustrative Examples

1. Single Stage Problem

The single stage reliability maximization problem is solved for optimum n and k using the values

$$l = 1.8 \quad \text{and} \quad s = 0.85 \quad (4.36)$$

An upper bound on n is estimated by Equation (4.8), that is,

$$n \leq l k^s \underset{0 < k \leq n}{\leq} l k_{\max}^s = l n^s = 1.8 n^{0.85} \quad (4.37)$$

This gives

$$n \leq 50 \quad (4.38)$$

The results with p as a parameter are shown in Table 4.2.

Unit Reliability	ACTIVE OR HOT STANDBY			COLD STANDBY		
	OPTIMUM n k		Maximum Reliability	OPTIMUM n k		Maximum Reliability
0.95	7	5	0.9962	7	5	0.9977
0.85	3	2	0.9392	3	2	0.9573
0.75	3	2	0.8437	3	2	0.8861

$l = 1.8$

$s = 0.85$

Table 4.2

Optimum solutions for single-stage problem

2. Two Stage Problem

Consider a two-stage reliability maximization problem as shown in Figure 4.5. The first stage is a cold standby k_1 -out-of- n_1 configuration and the second stage is an active k_2 -out-of- n_2 configuration.

The problem is to find the optimum values of n_1 , k_1 , n_2 and k_2 that maximize the total system reliability expressed by

$$R_T(n_1, n_2, k_1, k_2; p_1, p_2) = \left[\sum_{j=0}^{n_1-k_1} \frac{p_1^{k_1}}{j!} (-k_1 \ln p_1)^j \right] \left[\sum_{l=k_2}^{n_2} \binom{n_2}{l} p_2^l (1-p_2)^{n_2-l} \right] \quad (4.39)$$

subject to

$$n_1 k_{c1} k_1^{-s_1} + n_2 k_{c2} k_2^{-s_2} \leq C^* \quad (4.40)$$

$$n_1 - k_1 \geq 0 \quad (4.41)$$

$$n_2 - k_2 \geq 0 \quad (4.42)$$

$$n_1, n_2, k_1, k_2: \text{ integers not less than one} \quad (4.43)$$

C^* : the total allowable cost

In the case of multistage problems, there is no way to estimate appropriate upper bounds on the n 's, so we write an extra constraint

$$n_1, n_2 \leq 16 \quad (4.44)$$

The results with p_1 and p_2 as parameters are shown in Table 4.3. For the extreme case of $s_1 = s_2 = 0$, the results correspond to the optimal

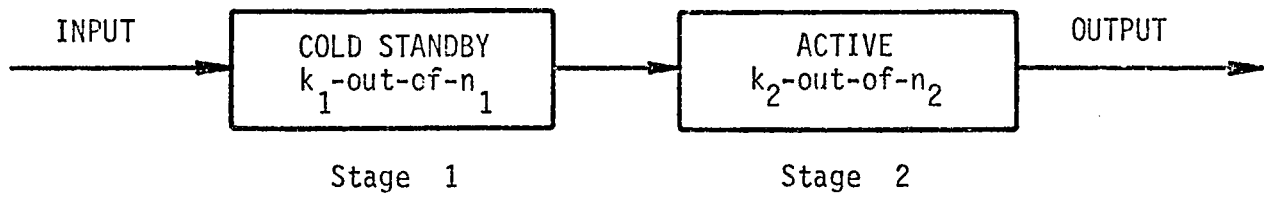


Figure 4.5
Two-Stage System

STAGE-1 (COLD STANDBY)					STAGE-2 (ACTIVE)					Cost	C*	R _T
P ₁	k _{C1}	s ₁	n ₁	k ₁	P ₂	k _{C2}	s ₂	n ₂	k ₂			
0.90	200	0.80	3	2	0.80	100	0.80	2	1	545	570	0.9415
0.90	200	0	1	1	0.80	100	0	3	1	500		0.8928
0.80	200	0.80	2	1	0.90	100	0.80	4	3	566		0.9273
0.80	200	0	2	1	0.90	100	0	1	1	500		0.8807
0.90	200	0.75	2	1	0.80	100	0.75	10	5	699	700	0.9885
0.90	200	0	2	1	0.80	100	0	3	1	700		0.9869
0.80	200	0.75	4	2	0.90	100	0.75	5	3	695		0.9809
0.80	200	0	2	1	0.90	100	0	3	1	700		0.9775

Table 4.3

Optimum solutions for two-stage problem

solutions for the conventional 1-out-of-n series-parallel configuration and are shown for comparison. The same results are obtained for a loosely guessed upper bounds on the n's.

$$n_1, n_2 \leq 32 \quad (4.45)$$

Therefore, the results shown in Table 4.3 are believed to be the optimal solutions without the extra constraint (4.44) or (4.45).

Conclusions

A new formulation of the optimal reliability design of process systems has been proposed, and is based on the k-out-of-n configuration. It is more suitable for reliability optimization of process systems than that based on the conventional 1-out-of-n parallel redundancy configuration.

An efficient computer program based on the Lawler and Bell [32] pseudo-Boolean algorithm has been developed to solve the above optimization problem. The program was written in FORTRAN IV. Appendix G contains a description of the program including a schematic diagram and a listing of the FORTRAN program.

CHAPTER V

CONCLUSIONS AND RECOMMENDATIONS

This dissertation consists of two main parts. First, an efficient general computer algorithm based on the path enumeration method has been developed to calculate reliability and sensitivity functions of complex systems. An important feature of the method is the module representation of the reliability graph. A modified integer gradient method is used for system optimization, which in general is a non-linear mixed integer programming problem subject to multiple linear or non-linear constraints. Although the method does not ensure a global optimum, it does find various near-optimum solutions. From a practical consideration, this could provide for a wider choice during the design phase. A number of examples are included to illustrate the method.

Second, a new formulation of the optimal reliability design of process systems is proposed. It takes into account the quantitative aspects of systems throughout and is based on the k-out-of-n configuration instead of the conventional parallel redundancy configuration. The problem is so formulated that determining the optimum configuration also determines the optimum capacity of units to be used at each stage of the system. A computer program based on the pseudo-Boolean algorithm of Lawler and Bell [32] is used to solve this non-linear integer programming problem. Several examples are included to illustrate the method.

With regard to further research in this area, the following two related problems are suggested:

1. Development of an algorithm to calculate exact reliability of systems which cannot be handled by path finding algorithm. This will be the case for systems with approximately more than 10 paths. The proposed method will first use the network reduction technique to reduce the reliability network to a non-series-parallel system and then apply the path enumeration technique to calculate the system reliability.

2. In the type of reliability network considered here it is assumed that modules are composed of different type of units. In practice, however, a particular type of unit may appear at many different locations in the system and it is usually desired to find the optimum unit reliability of each type to maximize the system objective function. Such a system is shown in Figure 5.1. The development of reliability and sensitivity expressions for this system and finally its optimization will be an extension of the techniques presented in this dissertation.

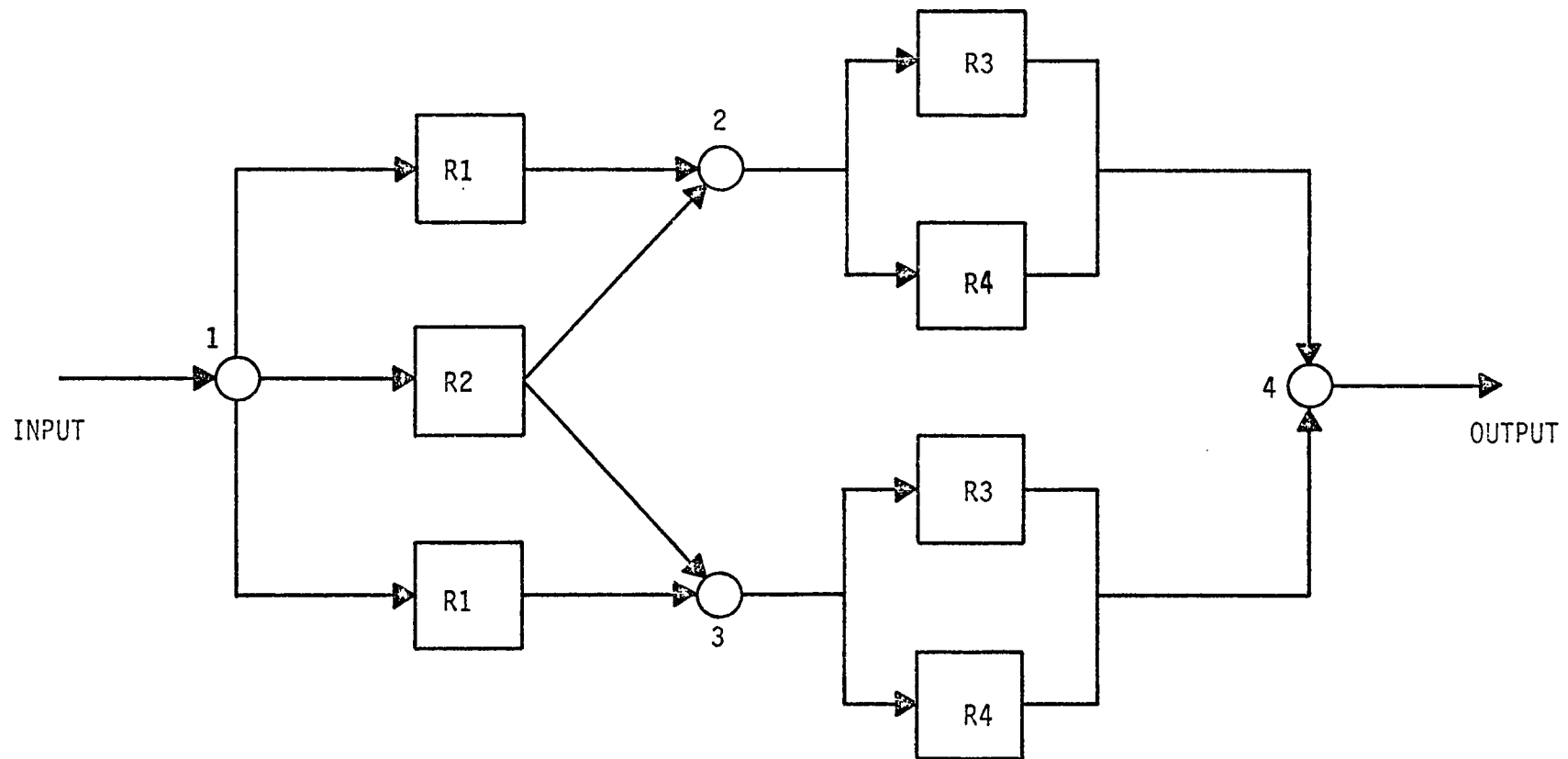


Figure 5.1

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NOMENCLATURE

A	- system availability
C	- upper limit on system cost
C^*	- upper limit on system cost
C_{A0}	- upper limit on cost constraint - system availability maximization
C_{Fi}	- capital cost of a unit in the i th module
C_i	- cost per unit of the i th module
C_I	- net income per hour
C_j	- j th minimal cut
\bar{C}_j	- complement of the event C_j
C_M	- maintenance cost per hour of system downtime
C_{P0}	- upper limit on cost constraint - system profit maximization
C_{TF}	- total fixed capital investment
C_{TI}	- total net income during system lifetime of T^+ hours
C_{TM}	- total maintenance cost during system lifetime of T^+ hours
f	- failure density function
F	- system unreliability
g_0	- objective function to be minimized
g_{i1}	- 1st term in the i th constraint, see Equation (4-20)
g_{i2}	- 2nd term in the i th constraint, see Equation (4-20)
g_{ij}	- contribution of i th module to i th constraint, see Equation (3-4)
G_j	- upper limit on j th constraint, see Equation (3-4)
K	- positive constant, see Equation (4-4)
k_c	- constant, corresponds to the cost of a single unit of capacity L , see Equation (4-6)
k_i	- minimum number of units in the i th module that might function for the module to operate successfully

k_v	- constant, corresponds to the volume of a single unit of capacity L, see Equation (4-12)
k_w	- constant, corresponds to the weight of a single unit of capacity L, see Equation (4-11)
L	- system capacity
\bar{M}	- system mean time between failure
MDT	- system mean downtime
MTBF	- system mean time between failure
$(MTBF)_i$	- mean time between failure of a unit of the i th module
MTBM	- system mean time between maintenance
N	- total number of modules in the system
n_i	- number of units in the i th module
N_i	- number of units in the i th module
P	- system profit
P_i	- i th minimal path
p_i	- reliability of a unit belonging to the i th stage, see Equation (4-13)
R	- system reliability
R_A	- reliability of active redundancy module
R_C	- reliability of cold standby redundancy module
R_H	- reliability of hot standby redundancy module
R_i	- i th module reliability
r_i	- reliability of a unit of the i th module
r_{ij}	- reliability of j th unit of the i th module
R_L	- lower bound approximation for R
R_U	- upper bound approximation for R
s, s_v, s_w	- non-negative constants less than one
S_{R_i}	- sensitivity of system reliability to the i th module reliability

S_{rij}	- sensitivity of system reliability to the reliability of the j th unit of the i th module
T	- system maintenance interval
t	- variable of integration
T^+	- system lifetime
T_D	- total time during which the system is down
T_E	- mean time for emergency maintenance on the system
T_i	- maintenance interval of a unit of the i th module
T_S	- mean time for scheduled maintenance on the system
T_u	- total time during which the system is operating
V^*	- upper limit on system volume
W^*	- upper limit on system weight
X	- percent of fixed capital investment
\underline{x}	- vector with elements $x_j = 0$ or 1
$\underline{x}_I, x_p, x^+$	- denote vectors of variables, real and/or integer
x_j	- j th binary variable
x_{kj}, x_{nj}	- j th binary variables belonging to the non-negative integer variables k and n , respectively, see Equations (4-26) and (4-27)
x'_{nj}	- given by $(1 - x_{nj})$, see Equation (4-28)
Z	- total number of terms involved in Equation (2-10)
λ	- exponential failure pdf parameter, also called the conditional failure rate
λ_i	- of the i th module
λ_i'	- of the i th module in standby
λ_i^0	- of the i th module in service
z	- hazard rate

APPENDIX A

MODULE RELIABILITY AND SENSITIVITY EXPRESSIONS

APPENDIX A

MODULE RELIABILITY AND SENSITIVITY EXPRESSIONS

Module Reliability

Active Redundancy Module - All n_i units are actively connected in parallel. If we let r_i be the probability that each component in module i functions for time t , then the probability that at least k_i out of n_i identical components function for time t is given by the binomial distribution [15]

$$R_{Ai}(n_i; k_i; r_i) = \sum_{l=k_i}^{n_i} \binom{n_i}{l} r_i^l (1-r_i)^{n_i-l} \quad (A-1)$$

The case $k_i=1$ leads to the familiar expression for the reliability of a parallel connection of n_i units

$$R_{Ai}(n_i, 1; r_i) = 1 - (1 - r_i)^{n_i} \quad (A-2)$$

When any one component is sufficient for module operation but the various components are different, the reliability expression is [52]:

$$R_{Ai}(n_i, 1; r_{ij}) = 1 - \prod_{j=1}^{n_i} (1 - r_{ij}) \quad (A-3)$$

where r_{ij} is reliability of the j th unit of module i .

When exponential failure probability density function (pdf) is assumed with parameter λ_i for identical components or λ_{ij} for different components, we have

$$\begin{aligned} r_i &= e^{-\lambda_i t} \\ r_{ij} &= e^{-\lambda_{ij} t} \end{aligned} \quad (A-4)$$

Standby Redundancy Module - We distinguish between three standby redundancy modules.

1. Hot Standby Redundancy Module

The units in hot standby have the same failure pdf as those in service. The module reliability expressions are identical to those in active redundancy module, Equations (A-1) to (A-4), that is,

$$R_{Hi}(n_i, k_i; r_i) = R_{Ai}(n_i, k_i; r_i) \quad (A-5)$$

2. Cold Standby Redundancy Module

A component or unit whose failure probability in standby is zero is called a cold standby. If the exponential failure pdf is assumed, the probability that at least k_i out of n_i identical components will function for time t is [7]:

$$R_{Ci}(n_i, k_i; t) = e^{-k_i \lambda_i t} \sum_{j=0}^{(n_i - k_i)} \frac{(k_i \lambda_i t)^j}{j!} \quad (A-6)$$

A specific case of cold standby arises when $k_i=1$.

$$R_{Ci}(n_i, 1; t) = e^{-\lambda_i t} \sum_{j=0}^{(n_i - 1)} \frac{(\lambda_i t)^j}{j!} \quad (A-7)$$

Although R_{Ci} cannot be obtained in a closed form for a general failure pdf, a recursive computational procedure can be used in the case $k_i=1$ to obtain an approximation to R_{Ci} . Let $g_i(t)$ be the failure pdf for the components in module i , and let $P_i(n_i, t)$ be the probability that n_i components are sufficient to operate module i for time t . Suppose that the components are used in the order $n_i, n_i-1, n_i-2, \dots, 2, 1$, then we have

$$P_i(1,t) \equiv 1 - \int_0^t g_i(\tau) d\tau \quad (A-8)$$

$$P_i(m,t) \equiv P_i(1,t) + \int_0^t P_i(m-1,t-\tau) g_i(\tau) d\tau \quad (A-9)$$

$$2 \leq m \leq n_i$$

The recursive Equation (A-9) gives the probability that m components are sufficient for module i for time t . It is expressed as the probability that the m th component alone will function until time t plus the probability that the m th component will fail at τ and the remaining $m-1$ components will function from τ until t . A discrete approximation to Equation (A-9) can be used to compute recursively for each m , an approximation to $P_i(m,t)$, which provides a point on the R_{Ci} function.

3. Warm Standby Redundancy Module

If the exponential failure pdf is assumed with parameter λ_i^0 in service and λ_i^1 in standby for units in i th module, Polovko [47] gives the following expression for the module reliability of a 1-out-of- n_i system

$$R_{Wi}(n_i, 1; t) = e^{-\lambda_i^0 t} \left[1 + \sum_{l=1}^{(n_i-1)} \frac{A_l}{l!} \left(1 - e^{-\lambda_i^1 t} \right)^l \right] \quad (A-10)$$

where

$$A_l = \sum_{j=0}^{l-1} \left(j + \frac{\lambda_i^0}{\lambda_i^1} \right) \quad (A-11)$$

For a general failure pdf $g_i(t)$ for the components in module i , Polovko [47] gives a recursive relationship for a 1-out-of- n_i system

$$R_{Wi}(n_i, 1; t) = R_{Wi}(n_i - 1, 1; t) + \int_0^t p_i^1(\tau) p_i^0(t - \tau) g_i^{(n_i - 1)}(\tau) d\tau \quad (A-12)$$

where $p_i^1(\tau)$ is the probability of failure free operation of the standby component up to the instant of its connection at time τ ; $p_i^0(t - \tau)$ is the probability of failure free operation of the standby component from the instant of its connection at time τ up to time t ; $g_i^{(n_i - 1)}(\tau)$ is the failure pdf for a system with $(n_i - 1)$ components.

Module Sensitivity

This is evaluated for a system in which individual units in a module have exponential failure pdf with parameter λ_i .

Active Redundancy Module -

$$\frac{\partial R_{Ai}(n_i, k_i; r_i)}{\partial n_i} = R_{Ai}(n_i, k_i; r_i) - R_{Ai}(n_i - 1, k_i; r_i) \quad (A-13)$$

Using the formula

$$\begin{pmatrix} n_i \\ l \end{pmatrix} = \begin{pmatrix} n_i - 1 \\ l - 1 \end{pmatrix} + \begin{pmatrix} n_i - 1 \\ l \end{pmatrix} \quad l \geq 1$$

we have from Equation (A-1)

$$R_{Ai}(n_i, k_i; r_i) = \sum_{l=k_i}^{n_i} \left[\begin{pmatrix} n_i - 1 \\ l - 1 \end{pmatrix} + \begin{pmatrix} n_i - 1 \\ l \end{pmatrix} \right] r_i^l (1 - r_i)^{n_i - l}$$

Expanding the summation, we obtain

$$R_{Ai}(n_i, k_i; r_i) = R_{Ai}(n_i-1, k_i; r_i) + \binom{n_i-1}{k_i-1} r_i^{k_i} (1-r_i)^{n_i-k_i} \quad (A-14)$$

Substituting in Equation (A-13)

$$\frac{\partial R_{Ai}(n_i, k_i; r_i)}{\partial n_i} = \binom{n_i-1}{k_i-1} r_i^{k_i} (1-r_i)^{n_i-k_i} \quad (A-15)$$

when $k_i=1$, we have from Equation (A-15)

$$\frac{\partial R_{Ai}(n_i, 1; r_i)}{\partial n_i} = r_i(1-r_i)^{n_i-1} \quad (A-16)$$

The above result also follows from Equation (A-2).

From Equation (A-3) we have

$$\frac{\partial R_{Ai}(n_i, 1; r_{ij})}{\partial r_{ij}} = \sum_{\substack{j=1 \\ j \neq n_i}}^{n_i} (1-r_{ij}) \quad (A-17)$$

When $r_i = e^{-\lambda_i t}$, we have from Equation (A-1)

$$\frac{\partial R_{Ai}(n_i, k_i; t)}{\partial t} = \lambda_i \sum_{l=k_i}^{n_i} \binom{n_i}{l} \left\{ (n_i-1)(e^{-\lambda_i t})^{l+1} (1-e^{-\lambda_i t})^{n_i-l-1} - 1(e^{-\lambda_i t})^l (1-e^{-\lambda_i t})^{n_i-l} \right\} \quad (A-18)$$

Expanding the summations, we obtain

$$\begin{aligned}
& \sum_{l=k_i}^{n_i} \binom{n_i}{l} (n_i-1) (e^{-\lambda_i t})^{l+1} (1-e^{-\lambda_i t})^{n_i-l-1} \\
& = n_i (e^{-\lambda_i t}) R_{Ai}(n_i-1, k_i; t)
\end{aligned} \tag{A-19}$$

And

$$\begin{aligned}
& \sum_{l=k_i}^{n_i} \binom{n_i}{l} l (e^{-\lambda_i t})^l (1-e^{-\lambda_i t})^{n_i-l} \\
& = n_i (e^{-\lambda_i t}) R_{Ai}(n_i-1, k_i, t) + k_i \binom{n_i}{k_i} (e^{-\lambda_i t})^{k_i} (1-e^{-\lambda_i t})^{n_i-k_i}
\end{aligned} \tag{A-20}$$

Substituting Equations (A-19) and (A-20) in Equation (A-18)

$$\frac{\partial R_{Ai}(n_i, k_i; t)}{\partial t} = -k_i \lambda_i \binom{n_i}{k_i} \left(e^{-\lambda_i t} \right)^{k_i} \left(1-e^{-\lambda_i t} \right)^{n_i-k_i} \tag{A-21}$$

Comparing Equations (A-15) and (A-21) we have

$$\frac{\partial R_{Ai}(n_i, k_i; t)}{\partial t} = -n_i \lambda_i \frac{\partial R_{Ai}(n_i, k_i; r_i)}{\partial n_i} \tag{A-22}$$

Standby Redundancy Module - Two types of standby redundancy modules are considered

1. Hot Standby Redundancy Module

The sensitivity expressions are identical to those developed for active redundancy module.

2. Cold Standby Redundancy

$$\frac{\partial R_{Ci}(n_i, k_i; t)}{\partial n_i} = R_{Ci}(n_i, k_i; t) - R_{Ci}(n_i - 1, k_i; t) \quad (A-23)$$

From Equation (A-6)

$$\frac{\partial R_{Ci}(n_i, k_i; t)}{\partial n_i} = e^{-k_i \lambda_i t} \frac{(k_i \lambda_i t)^{n_i - k_i}}{(n_i - k_i)!} \quad (A-24)$$

When $k_i = 1$, we have from above

$$\frac{\partial R_{Ci}(n_i, 1; t)}{\partial n_i} = e^{-\lambda_i t} \frac{(\lambda_i t)^{n_i - 1}}{(n_i - 1)!} \quad (A-25)$$

Using Equation (A-6)

$$\frac{\partial R_{Ci}(n_i, k_i; t)}{\partial t} = -k_i \lambda_i e^{-k_i \lambda_i t} \frac{(k_i \lambda_i t)^{n_i - k_i}}{(n_i - k_i)!} \quad (A-26)$$

Comparing Equations (A-24) and (A-26)

$$\frac{\partial R_{Ci}(n_i, k_i; t)}{\partial t} = -k_i \lambda_i \frac{\partial R_{Ci}(n_i, k_i; t)}{\partial n_i} \quad (A-27)$$

APPENDIX B

PATH FINDING ALGORITHM

APPENDIX B

PATH FINDING ALGORITHM

The algorithm is a modified version of the path finding algorithm developed by Henley and Williams [27]. The main features of the algorithm are presented here. For a more detailed analysis please refer to the original work of Henley and Williams [27].

The reliability flow graph consists of a set of independent modules. These are inter-connected by an array of points called nodes and directed lines between nodes called edges or branches. Each module has an input node and one or more output nodes. The nodes are numbered such that the arrows point from a node with a lower address to a node with a higher address. Thus, the input node of the system has the lowest address and the output node of the system has the highest address.

The algorithm to generate all the paths between the input and the output nodes of the system is best explained with the help of an example. Consider the reliability flow graph shown in Figure B-1. All the edges together with the corresponding modules are first ordered in a list such that the edges with the lowest starting address appear highest in the list. This is shown in Table B.1. The following steps then generate all the paths connecting the input node and the output node of the system.

1. List all the edges which originate at the input node of the system. These are the edges which appear at the top of the list shown in Table B.1. These will be referred to as the input edges. The modules corresponding to the edges are listed in another column.

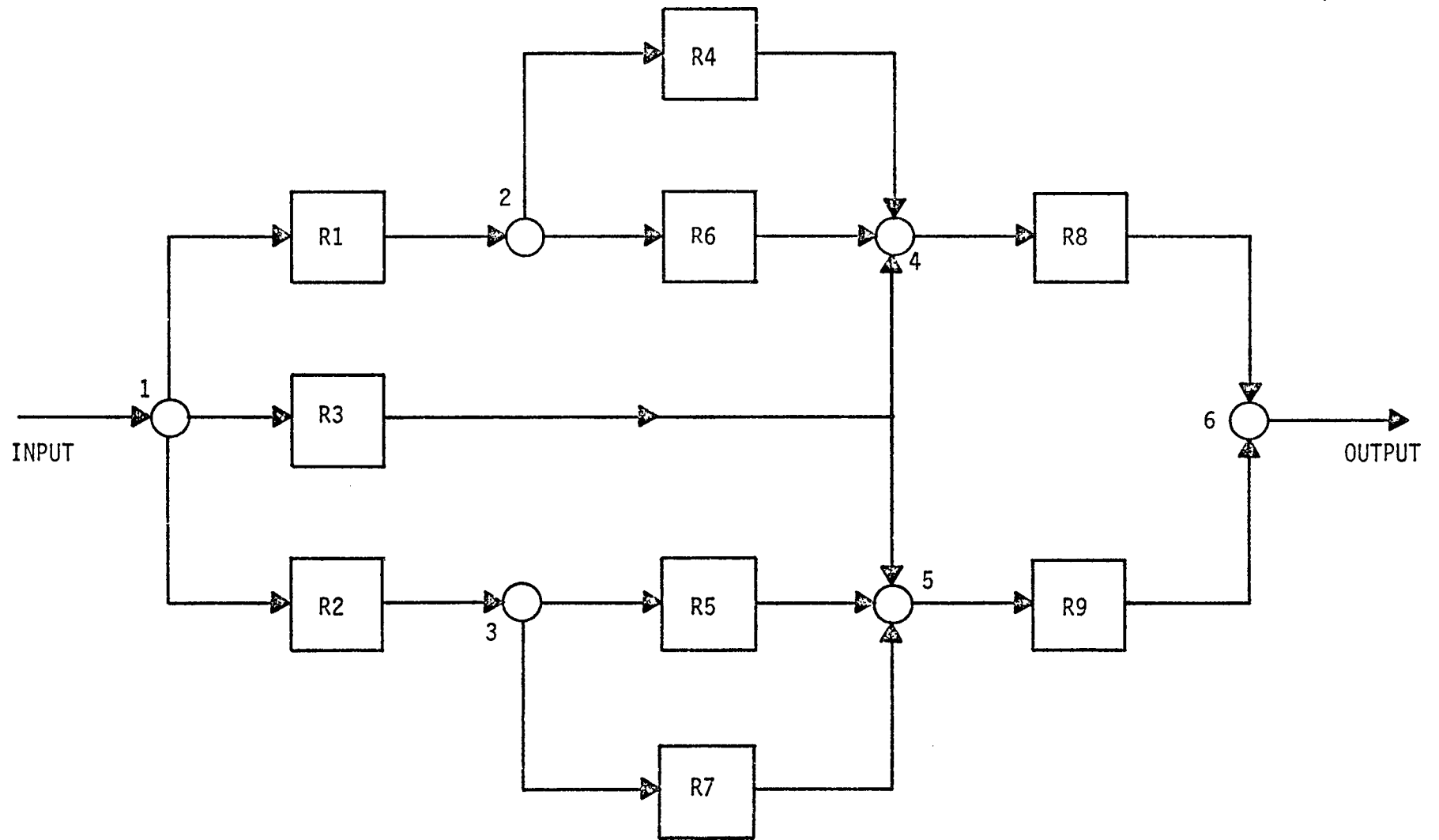


Figure B.1
Non-series-parallel system

<u>EDGE</u>	<u>MODULE</u>
1-2	R ₁
1-4	R ₃
1-5	R ₃
1-3	R ₂
2-4	R ₄
2-4	R ₆
3-5	R ₅
3-5	R ₇
4-6	R ₈
5-6	R ₉

Table B.1
System Configuration

2. Pick the first edge from the Table B.1 which follows the input edges. This will be referred to as the output edge. Append the output edge to those input edges whose ending node match its starting node. The corresponding modules are also appended and listed in another column. The new combinations formed are listed in the order they are generated. Next pick the edge following the output edge considered earlier. This will be the new output edge. Again append this output edge to those input edges and combinations formed earlier whose ending nodes match its starting node. The new combinations are listed in the order they are formed. This procedure continues until the last edge listed in Table B.1 has been used to generate new combinations. These operations are shown in Table B.2. The combinations whose last node is the same as the output node of the system, are the system paths. As Table B.2 illustrates, some of these paths have identical node configurations but are distinguishable by their different module configurations.

In the above discussion it is assumed that edges are numbered such that they point from a node with a lower address to a node with a higher address. These are called forward edges. The edges with nodes pointing from a higher address to a lower address are referred to as reverse edges. For a discussion of the path finding algorithm in the presence of reverse edges, please refer to the work of Williams [60].

	<u>Node Configuration</u>	<u>Module Configuration</u>
Edges starting at the input node "1"	1 - 2	R_1
	1 - 4	R_3
	1 - 5	R_3
	1 - 3	R_2
	1 - 2 - 4	$R_1 - R_4$
	1 - 2 - 4	$R_1 - R_6$
	1 - 3 - 5	$R_2 - R_5$
	1 - 3 - 5	$R_2 - R_7$
	1 - 4 - 6	$R_3 - R_8$
	1 - 2 - 4 - 6	$R_1 - R_4 - R_8$
Combinations ending at the output node "6": giving the system paths	1 - 2 - 4 - 6	$R_1 - R_6 - R_8$
	1 - 5 - 6	$R_3 - R_9$
	1 - 3 - 5 - 6	$R_2 - R_5 - R_9$
	1 - 3 - 5 - 6	$R_2 - R_7 - R_9$

Table B.2

Combination of Edges

APPENDIX C

PATH UNIONS

APPENDIX C

PATH UNIONS

The procedures to generate all the necessary path unions are best illustrated with the help of an example. Consider the reliability network shown in Figure C.1. The path finding algorithm traces all the paths connecting the input and the output nodes of the system. These paths in their node and module configurations are listed in Table C.1.

The following steps then generate the path unions:

1. Transform each path into the binary form; i.e. the locations containing 1's (0's) indicate the modules present (absent) in the path.
2. Generate the path unions by using a bitwise logical 'OR' operation in order of ascending path number. The 2-path unions are generated from the single path unions (or paths) $P-1 \text{ OR } P-2, P-1 \text{ OR } P-3, \dots, P-1 \text{ OR } P-M, P-2 \text{ OR } P-3, \dots, P-2 \text{ OR } P-M, \dots, P-(M-1) \text{ OR } P-M$, where M represents the total number of paths in the system. For the above example, 2-path unions are from (5) to (10) in Table C.2. Then proceed to form all the 3-path unions from the 2-path unions obtained earlier. These are $P-1 \text{ OR } P-2 \text{ OR } P-3, \dots, P-1 \text{ OR } P-2 \text{ OR } P-M, \dots, P-(M-2) \text{ OR } P-(M-1) \text{ OR } P-M$. 3-path unions are from (11) to (14) in Table C.2. Similar procedures continue until we form the M -path union, $P-1 \text{ OR } P-2 \text{ OR } \dots \text{ OR } P-(M-1) \text{ OR } P-M$. For the example, number (15) shows the 4-path union. There is a sign bit attached to each path and path union and it changes sign with the generation of each higher order path union. Thus, it is positive for paths, negative for 2-path unions, positive for 3-path unions, etc., and $(-1)^{M-1}$ for the M -path union.
3. Any path union generated as outlined in Step 2 that has 1's in every location in its binary representation and whose end path in the path

union is not the Mth path is rejected. This is so because it generates identical path unions with equal number of positive and negative sign bits. In the example, the 3-path union numbered (11) i.e. P-1 OR P-2 OR P-3, generates the 4-path union numbered (15), i.e. P-1 OR P-2 OR P-3 OR P-4. The two are identical and have opposite signs and hence cancel each other. These are excluded from the list of path and path unions. The reduction in number of terms may be larger for some other systems. The system shown in Figure 2.9 has six paths. The path enumeration technique predicts $2^6 - 1 = 63$ terms in the reliability expression. Introducing the above modification reduces the number of terms to 47. The savings in computer time is significant.

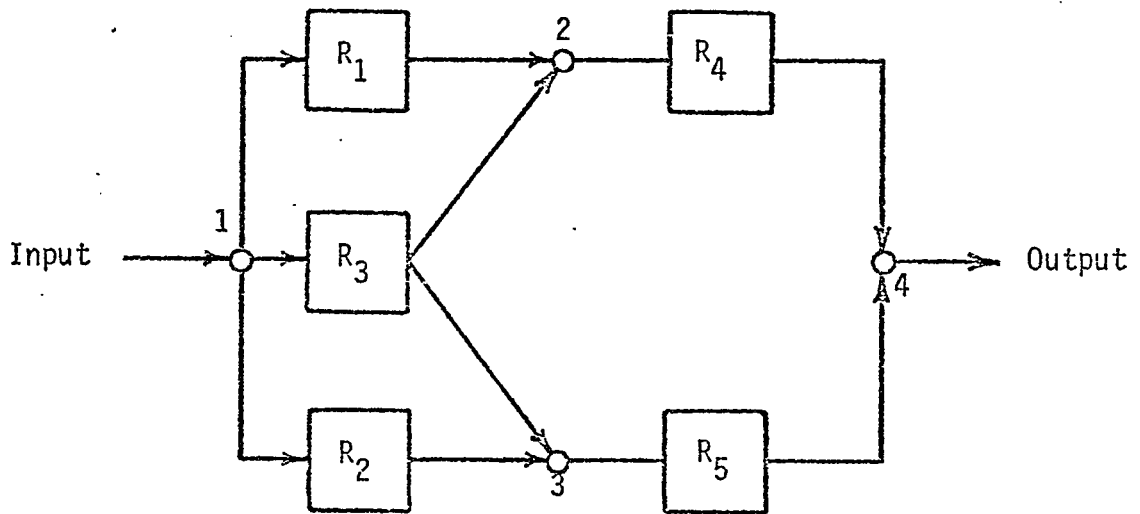


Figure C.1

Reliability Graph

Path	Node Configuration	Module Configuration
P-1	1 - 2 - 4	R1 - R4
P-2	1 - 2 - 4	R3 - R4
P-3	1 - 3 - 4	R2 - R5
P-4	1 - 3 - 4	R3 - R5

Table C.1

Path Detection

No.	Path or Path Union	Sign	Module				
			R1	R2	R3	R4	R5
1	P-1	+	1	0	0	1	0
2	P-2	+	0	0	1	1	0
3	P-3	+	0	1	0	0	1
4	P-4	+	0	0	1	0	1
5	P-1 OR P-2	-	1	0	1	1	0
6	P-1 OR P-3	-	1	1	0	1	1
7	P-1 OR P-4	-	1	0	1	1	1
8	P-2 OR P-3	-	0	1	1	1	1
9	P-2 OR P-4	-	0	0	1	1	1
10	P-3 OR P-4	-	0	1	1	0	1
*11	P-1 OR P-2 OR P-3	+	1	1	1	1	1
12	P-1 OR P-2 OR P-4	+	1	0	1	1	1
13	P-1 OR P-3 OR P-4	+	1	1	1	1	1
14	P-2 OR P-3 OR P-4	+	0	1	1	1	1
*15	P-1 OR P-2 OR P-3 OR P-4	-	1	1	1	1	1

*Path unions (11) and (15) will not appear in the final list.

Table C.2

Path and Path Unions

APPENDIX D

MODIFIED INTEGER GRADIENT METHOD

APPENDIX D

MODIFIED INTEGER GRADIENT METHOD

The general nonlinear integer programming problem can be stated as follows:

Maximize the nonlinear objective function

$$f(\underline{X}) \quad (D-1)$$

subject to several functional nonlinear (or linear) constraints

$$g_i(\underline{X}) \leq 0 \quad i = 1, 2, \dots, m \quad (D-2)$$

non-negative constraints

$$\underline{X} \geq \underline{0} \quad (D-3)$$

and finally the integer constraints

$$\underline{X} : \text{integral} \quad (D-4)$$

where \underline{X} is an n vector.

Optimal Seeking Algorithm

The procedure described below is mainly due to Reiter and Rice [50]. The search algorithm has three phases.

Phase (1): Selection of an Initial Point

The simplest method of selecting initial point is a statistical method using uniform random numbers. Assume that independent upper limits on \underline{X} also exist.

$$0 \leq x_j \leq u_j \quad j = 1, 2, \dots, n \quad (D-5)$$

Equation (D-5) together with Equation (D-4) define a finite set of points enclosed by a hyperrectangle in n -space. So we sample with a uniform distribution from the set of integers in the interval $[0, u_j]$ $j = 1, 2, \dots, n$. This results in a uniform sampling of lattice points in the hyperrectangle.

If the selected initial point \underline{X}_I is feasible, we go directly to Phase (3), otherwise, we go to Phase (2).

Phase (2): Search for a Feasible Solution \underline{X}^+

A search by the method of "weighted perpendiculars" (WP) may be the most general one. The quantity S_i

$$S_i = g_i(\underline{X}_I) \quad i = 1, 2, \dots, m \quad (D-6)$$

is the amount by which the i th constraint is violated at \underline{X}_I . We define

$$\underline{Z} = \frac{- \left[\sum_{i: S_i > 0} S_i [\nabla_i(g_i(\underline{X}))]_{\underline{X}=\underline{X}_I} \right]}{\sum_{i: S_i > 0} S_i} \quad (D-7)$$

where ∇_i denotes the vector of partial derivatives of the i th constraint with respect to x_i 's. The vector \underline{Z} is thus a positive linear combination of these perpendiculars, and, therefore, lies in the convex cone spanned by them, as shown in Figure D-1.

$$\text{Let } Z_j^\# = \min_{Z_j \neq 0} |Z_j| \quad (D-8)$$

$$\text{Define } \underline{Z}^+ = (Z_1^+, Z_2^+, \dots, Z_n^+)^T$$

$$Z_j^+ = \begin{cases} [(Z_j/Z_j^\#) + 0.5], & \text{if } Z_j \geq 0 \\ [(Z_j/Z_j^\#) - 0.5], & \text{if } Z_j < 0 \end{cases} \quad j = 1, 2, \dots, n \quad (D-9)$$

where $[w]$ denotes the largest integer not exceeding w .

The equation

$$\underline{Y}_b = \underline{X}_I + b \underline{Z}^+ \quad (D-10)$$

is the equation of a line passing through the point \underline{X}_I and has possible direction towards feasible region, as shown in Figure D-1. Taking successively the values $b = 1, 2, \dots$ we test the corresponding \underline{Y}_b , necessarily a lattice point, for feasibility; we eventually either locate a feasible point or violate one or more of the lower or upper limits. If one or more of these limits is violated, we hold the corresponding coordinate of \underline{Y}_b constant at those limits and continue to search by increasing b . Ultimately, we find a feasible point \underline{X}^+ or all coordinates go outside of their limit. In the second case we abandon \underline{X}_I and sample another starting point from the hyperrectangle as explained in Phase (1). In the case of linear constraints, another method called "permuted coordinates" (PC) may be better [50].

Phase (3): Search for the Locally Optimum Point \underline{X}_p

Having located a feasible lattice point \underline{X}^+ we start Phase (3), where we end up with a locally maximal point.

Let $\underline{h}(\underline{X})$ denote the gradient of the objective function

$$\text{i.e. } \underline{h}(\underline{X}) = \nabla f(\underline{X}) \quad (D-11)$$

$$\text{Define } \underline{Y}_d = \underline{X}^+ + d \cdot \underline{h}(\underline{X}^+) \quad (D-12)$$

Since, $\underline{h}(\underline{X}^+)$ is not in general integer valued, we apply the procedure used in the WP method to obtain a normalized integer gradient vector \underline{h}^+ , i.e.

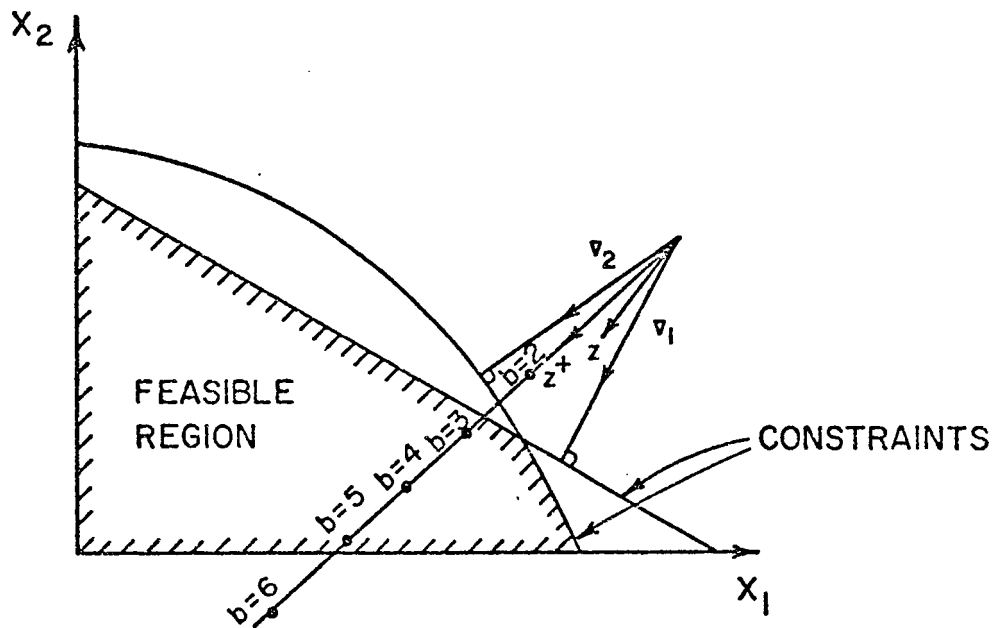


Figure D-1 Method of Weighted Perpendiculars (WP)

$$h_j^\# = \min_{h_j \neq 0} |h_j| \quad (D-13)$$

$$h_j^+ = \begin{cases} [(h_j/h_j^\#) + 0.5], & \text{if } h_j \geq 0 \\ [(h_j/h_j^\#) - 0.5], & \text{if } h_j \leq 0 \end{cases} \quad j = 1, 2, \dots, n \quad (D-14)$$

The vector \underline{h}^+ is the modified gradient of the objective function $f(\underline{X})$ and

$$\underline{Y}_d = \underline{X}^+ + d \cdot \underline{h}^+ \quad (D-15)$$

defines the path on which we search for improved values of the objective function. We find that value of d for which $f(\underline{Y}_d)$ is a maximum over feasible \underline{Y}_d on this path. This may be done by letting $d = 0, 1, 2, \dots$ successively and finding the first value of d for which either \underline{Y}_{d+1} is infeasible or $f(\underline{Y}_{d+1}) \leq f(\underline{Y}_d)$. If the best value of d denoted by d^* , is not zero, we set $\underline{X}^+ = \underline{Y}_{d^*}$ and repeat the procedure. If $d^* = 0$, we set the element of $\underline{h}(\underline{X}^+)$ of the smallest non-zero absolute value equal to zero to obtain $\underline{h}^1(\underline{X}^+)$ and proceed as before with \underline{h}^1 in place of \underline{h} . If \underline{h}^1 leads to no improvement, set the element of \underline{h}^1 of smallest non-zero absolute value equal to zero to obtain $\underline{h}^2(\underline{X}^+)$ and proceed with \underline{h}^2 in place of \underline{h} . This procedure is continued until \underline{h}^v is obtained with only one non-zero element. If using $\underline{h}^v(\underline{X}^+)$ results in no improvement, we try successively \underline{h}^{v+1} to \underline{h}^{v+k} where $k+1$ is the number of non-zero elements of \underline{h} and \underline{h}^{v+j} is a vector with 1 in the coordinate corresponding to the partial derivative in \underline{h} which ranks $(j+1)^{th}$ in absolute value, for $j \leq k$. If \underline{h}^{v+k} results in no improvement then \underline{X}^+ is a locally maximal point. If any \underline{h}^j leads to improvement, we set \underline{X}^+ equal to the \underline{Y}_{d^*} thus obtained and repeat the procedure.

Extension to Mixed Integer Problem

The solution procedure described above can be modified to apply to mixed integer problems.

To handle a mixed problem with $n' < n$ variables required to be integer valued, the variables are renumbered so that the first n' are the integer valued ones. Equations (D-8) and (D-9) for Z^+ and R^+ are replaced by

$$Z_j^{\#} = \min. \{ |Z_j| : Z_j \neq 0, 1 \leq j \leq n' \} \quad (D-16)$$

$$Z_j^+ = \begin{cases} [(Z_j/Z_j^{\#}) + 0.5], & \text{if } Z_j \geq 0 \text{ and } 1 \leq j \leq n' \\ [(Z_j/Z_j^{\#}) - 0.5], & \text{if } Z_j \leq 0 \text{ and } 1 \leq j \leq n' \\ (Z_j/Z_j^{\#}) & , \quad \text{if } n' \leq j \leq n \end{cases} \quad (D-17)$$

As long as the gradient $\underline{h}^j(\underline{x}^+)$ has some non-zero values among its first n' coordinates, the procedures described earlier, which restrict attention to integer values of the path parameters, apply directly. If all of the first n' coordinates of $\underline{h}^j(\underline{x}^+)$ are zero, it is unnecessary to normalize and adjust to integers. A grid is superimposed on the modified gradient path and procedures described in Phase (3) are continued.

In the case of functions and constraints which are not differentiable, paths of positive ascent calculated by first differences can be used in place of the gradient.

It should be noted that the method is not well suited to dealing with equality constraints.

APPENDIX E

SENSITIVITY EXPRESSIONS FOR SYSTEM AVAILABILITY AND SYSTEM PROFIT

APPENDIX E
SENSITIVITY EXPRESSIONS FOR SYSTEM
AVAILABILITY AND SYSTEM PROFIT

$$A(\underline{n}, T) = \frac{\int_0^T R(\underline{n}, t) dt}{\int_0^T R(\underline{n}, t) dt + T_E - (T_E - T_S) R(\underline{n}, T)} \quad (E-1)$$

From Equation (E-1)

$$\begin{aligned} \frac{\partial A(\underline{n}, T)}{\partial n_i} = & \frac{\left[\frac{\partial}{\partial n_i} \int_0^T R(\underline{n}, t) dt \right] \left[T_E - (T_E - T_S) R(\underline{n}, T) \right] + \int_0^T R(\underline{n}, t) dt \left[(T_E - T_S) \frac{\partial R(\underline{n}, t)}{\partial n_i} \right]}{\left[\int_0^T R(\underline{n}, t) dt + T_E - (T_E - T_S) R(\underline{n}, T) \right]^2} \\ & i = 1, 2, \dots N \end{aligned} \quad (E-2)$$

$$\frac{\partial A(\underline{n}, T)}{\partial T} = \frac{T_E \cdot R(\underline{n}, T) - (T_E - T_S) R^2(\underline{n}, T) + (T_E - T_S) \left[\int_0^T R(\underline{n}, t) dt \right] \frac{\partial R(\underline{n}, T)}{\partial T}}{\left[\int_0^T R(\underline{n}, t) dt + T_E - (T_E - T_S) R(\underline{n}, T) \right]^2} \quad (E-3)$$

$$P(\underline{n}, T) = C_I \cdot T^+ \cdot A(\underline{n}, T) - \left[\sum_{i=1}^N n_i \cdot C_{Fi} \right] \left[1 + (1 - A(\underline{n}, T)) \frac{X}{100} \cdot T^+ \right] \quad (E-4)$$

$$\begin{aligned} \frac{\partial P(\underline{n}, T)}{\partial n_i} &= \left[C_I + \left(\sum_{i=1}^N n_i \cdot C_{Fi} \right) \frac{\chi}{100} \right] T^+ \cdot \frac{\partial A(\underline{n}, T)}{\partial n_i} \\ &\quad - C_{Fi} \left[1 + (1-A(\underline{n}, T)) \frac{\chi}{100} \cdot T^+ \right] \\ i &= 1, 2, \dots, N \end{aligned} \quad (E-5)$$

$$\frac{\partial P(\underline{n}, T)}{\partial T} = \left[C_I + \left(\sum_{i=1}^N n_i \cdot C_{Fi} \right) \frac{\chi}{100} \right] T^+ \cdot \frac{\partial A(\underline{n}, T)}{\partial T} \quad (E-6)$$

Solution of Equations (E-1) to (E-6) involves the evaluation of the following terms,

1. $R(\underline{n}, T)$ - The system reliability $R(\underline{n}, T)$ is calculated according to the general computer algorithm developed in Chapter II.

$$\begin{aligned} 2. \quad \frac{\partial R(\underline{n}, T)}{\partial T} &= \sum_{i=1}^N \frac{\partial R(\underline{n}, T)}{\partial R_i(n_i, T)} \cdot \frac{\partial R_i(n_i, T)}{\partial T} \\ &= \sum_{i=1}^N \left(R \Big|_{R_i=1} - R \Big|_{R_i=0} \right) \frac{\partial R_i(n_i, T)}{\partial T} \end{aligned} \quad (E-7)$$

where $\frac{\partial R_i(n_i, T)}{\partial T}$ is given by either Equations (A-21) or (A-26).

$$\begin{aligned} 3. \quad \frac{\partial R(\underline{n}, T)}{\partial n_i} &= \frac{\partial R(\underline{n}, T)}{\partial R_i(n_i, T)} \cdot \frac{\partial R_i(n_i, T)}{\partial n_i} \\ &= \left(R \Big|_{R_i=1} - R \Big|_{R_i=0} \right) \frac{\partial R_i(n_i, T)}{\partial n_i} \end{aligned} \quad (E-8)$$

where $\frac{\partial R_i(n_i, T)}{\partial n_i}$ is given by either Equations (A-15) or (A-24).

4. $\int_0^T R(\underline{n}, t) dt$ - This term is evaluated by numerical integration using

Simpson's formula.

$$\begin{aligned}
 5. \quad \frac{\partial}{\partial n_i} \int_0^T R(\underline{n}, t) dt &= \int_0^T \frac{\partial R(\underline{n}, t)}{\partial n_i} dt \\
 &= \int_0^T \left(\frac{\partial R(\underline{n}, t)}{\partial R_i(n_i, t)} \right) \left(\frac{\partial R_i(n_i, t)}{\partial n_i} \right) dt \\
 &= \int_0^T \left(R \Big|_{R_i=1} - R \Big|_{R_i=0} \right) \left(\frac{\partial R_i(n_i, t)}{\partial n_i} \right) dt
 \end{aligned}$$

where $\frac{\partial R_i(n_i, t)}{\partial t}$ is given by either Equations (A-15) or (A-24) at any time t . Again the integral is evaluated by using Simpson's formula.

APPENDIX F

LAWLER AND BELL ALGORITHM

APPENDIX F

LAWLER AND BELL ALGORITHM

Lawler and Bell [32] describe a simple, easily programmed algorithm for solving discrete optimization problems with monotone objective functions and arbitrary constraints.

A brief review of the Lawler-Bell method is provided in this appendix. The type of problems that can be solved by this method may be put in the following form:

Minimize

$$g_0(\underline{x})$$

subject to m constraints of the form

$$g_{i1}(\underline{x}) - g_{i2}(\underline{x}) \leq 0 \quad i=1,2,\dots,m$$

where

$$\underline{x} = (x_1, x_2, x_3, \dots, x_n) \tag{F-1}$$

and

$$x_j = 0 \text{ or } 1, \quad j=1,2,\dots,n$$

and where the restriction is applied that each of the functions $g_0, g_{i1}, g_{i2}, (i=1,2,\dots,m)$ is monotone non-decreasing in each of the variables x_1, x_2, \dots, x_n . With some ingenuity, many problems can be put in this form.

Some preliminaries discussed below are essential for the description of the algorithm.

Vector \underline{x} is "binary" in the sense that each x_j is either 0 or 1; $\underline{x} \leq \underline{y}$ if and only if $x_j \leq y_j$ for $j=1,2,\dots,n$. This is the vector partial

ordering. There is also the lexicographic or numerical ordering of these vectors that is obtained by identifying with each vector \underline{x} , the integer value

$$N(\underline{x}) = x_1 2^{n-1} + x_2 2^{n-2} + \dots + x_n 2^0 \quad (\text{F-2})$$

Numerical ordering is a refinement of the vector partial ordering; i.e. $\underline{x} \leq \underline{y}$ implies $N(\underline{x}) \leq N(\underline{y})$; however, $N(\underline{x}) \leq N(\underline{y})$ does not imply $\underline{x} \leq \underline{y}$.

Suppose all binary n -vectors are listed in numerical order; i.e.

$(0, \dots, 0, 0, 0),$
 $(0, \dots, 0, 0, 1),$
 $(0, \dots, 0, 1, 0),$
 $(0, \dots, 0, 1, 1),$
 $(0, \dots, 1, 0, 0),$
 etc.

Immediately following an arbitrary vector \underline{x} , there may (or may not) be a number of vectors \underline{x}' with the property that $\underline{x} \leq \underline{x}'$. Roughly speaking these are vectors that differ from \underline{x} only in that they have 1's in place of one or more of the "right most" 0's of \underline{x} . For example, immediately following $\underline{x} = (0, 1, 0, 0)$ are $(0, 1, 0, 1)$, $(0, 1, 1, 0)$, $(0, 1, 1, 1)$, each of which is greater than \underline{x} in the vector partial ordering.

Let \underline{x}^* denote the first vector following \underline{x} in the numerical ordering that has the property that $\underline{x} \not\leq \underline{x}^*$. For any given \underline{x} , the vector \underline{x}^* is very easily calculated on the computer as follows:

Treat \underline{x} as a binary number:

1. Subtract 1 from \underline{x} to obtain $\underline{x} - 1$,
2. Logically OR \underline{x} and $\underline{x} - 1$ to obtain $\underline{x}^* - 1$,

3. Add 1 to obtain \underline{x}^* .

Note that $\underline{x}^* - 1$ is greater than each of $\underline{x}, \underline{x} + 1, \dots, \underline{x}^* - 2$ in the vector partial ordering.

The Lawler-Bell method is basically a search method, which starts with $\underline{x} = (0, 0, \dots, 0, 0)$ and examines the 2^n solution vectors in the numerical ordering described above. Further, the labor of examination is considerably cut down by following certain rules. As the examination proceeds through the list of vectors, a record is kept of the least costly solution found to date. If $\hat{\underline{x}}$ denotes this solution having cost $g_0(\hat{\underline{x}})$ and \underline{x} is the vector being examined, then the following steps indicate the conditions under which certain vectors in the numerical ordering can be skipped over:

1. If $g_0(\underline{x}) \geq g_0(\hat{\underline{x}})$, skip to \underline{x}^* and repeat the operation; otherwise proceed to Step 2.

Justification: Because g_0 is monotone non-decreasing, none of the vectors $\underline{x} + 1, \underline{x} + 2, \dots, \underline{x}^* - 1$ can be less costly than $\hat{\underline{x}}$.

2. If for any i ($i=1, 2, \dots, m$), $g_{i1}(\underline{x}^* - 1) - g_{i2}(\underline{x}) < 0$, skip to \underline{x}^* and go to Step 1; otherwise proceed to Step 3.

Justification: With respect to vectors in the interval $(\underline{x}, \underline{x}^* - 1)$, $\underline{x}^* - 1$ maximizes g_{i1} and \underline{x} maximizes $-g_{i2}$. Hence if it is the case that $g_{i1}(\underline{x}^* - 1) - g_{i2}(\underline{x}) < 0$, there can be no vector \underline{x}' in the interval such that $g_{i1}(\underline{x}') - g_{i2}(\underline{x}') \geq 0$.

3. If $g_{i1}(\underline{x}) - g_{i2}(\underline{x}) \geq 0$, ($i=1, 2, \dots, m$), replace $\hat{\underline{x}}$ by \underline{x} and skip to \underline{x}^* ; otherwise change \underline{x} to $\underline{x} + 1$. In either case further execution is transferred to Step 1.

Justification: Because g_0 is monotone non-decreasing, none of the vectors $\underline{x} + 1, \underline{x} + 2, \dots, \underline{x}^* - 1$ can be less costly than \underline{x} .

Following the above steps, all the vectors are examined and scanning continues until a vector having maximum numerical order. Viz., $(1,1,\dots,1)$ is found. In case one has skipped to a vector having numerical order higher than $(1,1,\dots,1)$, designate this state by "overflow" and terminate the procedure. The least "costly" vector recorded provides the optimum solution.

Problems involving non-negative integer variables can be transformed into problems involving binary variables by means of the substitution

$$x_j = x_{j1} + 2x_{j2} + 2^2x_{j3} + \dots + 2^{K-1}x_{jK} \quad (F-3)$$

where $x_{ji} = 0$ or 1 ($i = 1,2,\dots,K$). K is chosen to be sufficiently large for $2^K - 1$ to be an upper bound on the value of x_j .

The algorithm was seen to be most efficient if those variables that were roughly speaking "least significant" were assigned positions at the "right" end of the solution vector. The ordering of vector \underline{x} is shown below

$$\underline{x} = (x_{1K}, x_{2K}, \dots, x_{nK}, x_{1(K-1)}, \dots, x_{n2}, x_{11}, x_{21}, \dots, x_{n1}) \quad (F-4)$$

APPENDIX G
COMPUTER PROGRAMS

APPENDIX G

COMPUTER PROGRAMS

This appendix contains the documentation for two computational schemes which solve the system optimization problems in the text.

Each scheme is composed of subroutines. Most of the subroutines do not have calling arguments--the greater part of the communication between subroutines being handled through COMMON.

Computational Scheme I

Computational Scheme I is based on the modified integer gradient method. The details of the method are given in Appendix D. A schematic diagram of the main program is shown in Figure G.2. The computation is composed of fifteen subroutines. Depending on the objective function and the constraints, the appropriate statements are modified in the various subroutines. A computer listing is provided for each of the three system optimization problems discussed in Chapter III of the text. Subroutines PATHS, PATHP, TRACE, and PATHC are identical for the three cases. A functional description of the various subroutines and a list of the major program variables is given below.

SUBROUTINES

- DATAIN - Reads and echo checks data.
- PATHS - Argument (INPUT). Finds all paths in the reliability flow graph which start at a specified INPUT node. See Appendix B for details.
- PATHP - Argument (OUTPUT). Itemizes all paths starting at an INPUT node (defined by a call to subroutine PATHS) and ending at the OUTPUT node specified.

- TRACE - Argument (I). Retrieves the nodes and the modules associated with the Ith entry in the P-array and prints them in sequence. Also converts each path in the binary form.
- PATHC - Finds the binary representation of all the path combinations with appropriate signs. See Appendix C for details.
- RANDU - Arguments (IX, IY, IFL). Generates uniformity distributed number between zero and one.
- MODULE - Computes the module reliability and the sensitivity of the module reliability to the module redundancy for all the system modules. Only active redundancy modules or cold standby redundancy modules are considered. See Appendix A for details.
- RELIAB - Computes the system reliability from the module reliabilities found earlier.
- SENSE - Determines the sensitivity of the system reliability to the module redundancies.
- AVLBTY - Determines the system availability, system profit, and the sensitivity of the system availability and the system profit to the module redundancies and to the maintenance interval. The integrals are evaluated by using Simpson's formula. See Appendix E for details.
- CONSTR - Computes the system constraints.
- FESIBL - Finds the feasible solution by the method of weighted perpendiculars. See Appendix D for details.
- SORT - Sorts the system sensitivity vector in ascending order.
- OPTIMZ - Finds a locally optimal solution from the feasible solution. The computation is based on the modified integer gradient method discussed in Appendix D.
- OPTIM - Checks to see if the movement in the gradient direction gives an improved feasible solution.

PROGRAM VARIABLES

An explanation of the major program variables is given below.

- INNOD - Input node.
- OUTNOD - Output node.
- FOR - Number of forward edges.
- FREV - Row number of the first reverse edge appearing in the P-array.

- LEND - Row number of the last entry in the P-array.
- NEDGES - Total number of edges or branches in the reliability graph.
- NEDGS1 - NEDGES + 1.
- NNODES - Total number of nodes.
- NPATH - Total number of paths between the input and the output node.
- PATHL - Row dimension of the PATH-array. Length of vector PSIGN.
- PLENTH - Row dimension of the P-array.
- PPL - Dimension of PP, PI4, NUNIT, NMAX, NMIN, MUNIT, RUNIT, WUNIT, LAMDA, G, AGRADN, RGRADN, PGRADN, RIGRDN, MTGRDN, COSTF, and NSORT vectors. Also the column dimension of PATH-array.
- REV - Number of reverse edges.
- STATUS - Flag which is set to zero if an attempt has been made to overflow the PATH-array or the P-array.
- P - Array of dimension P (PLENTH,4) which contains the list structure shown in Figure G-1.
- PP - Auxiliary vector.
- PI4 - Auxiliary vector.
- PATH - Array of dimension (PATHL, PPL). Stores each path and path combinations in binary form.
- NMOD - Total number of modules.
- NMOD1 - NMOD + 1.
- MPATH - Total number of paths and path combinations.
- PSIGN - Vector. Stores the number + 1 or -1 associated with each path and path combination.
- NUNIT - Vector. Stores the module redundancies (number of redundant units in each module).
- NMAX - Vector. Contains the upper limit of the module redundancies.
- NMIN - Vector. Contains the lower limit of the module redundancies.
- MUNIT - Vector. Stores the module redundancies necessary for successful operation of each module.

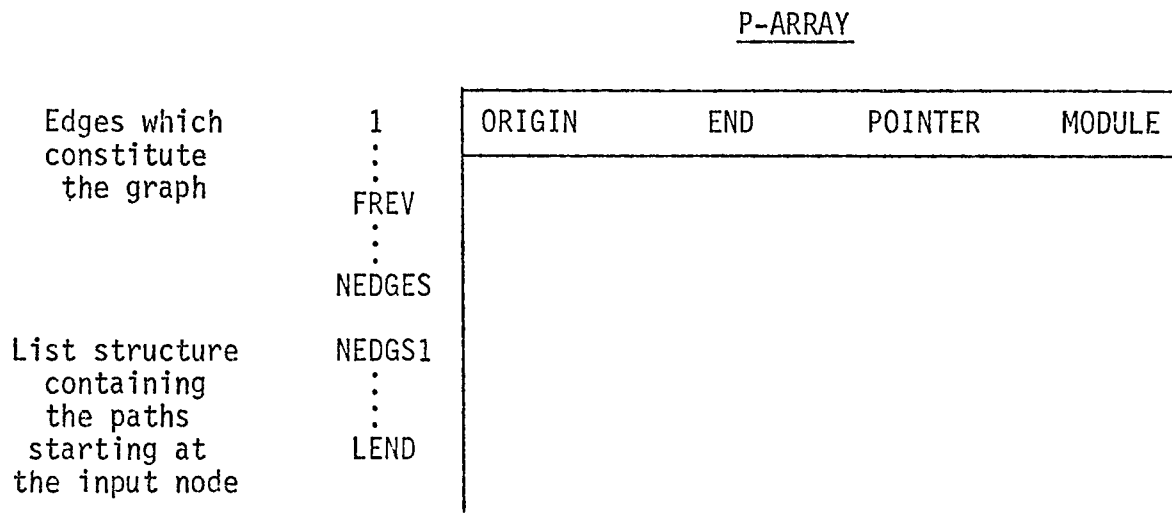


Figure G.1

List Structure of P-Array

KODE - Vector. Stores the binary numbers zero or one. A "zero" identifies an active redundancy module. A "one" identifies a cold standby redundancy module.

ASYS - System availability.

RSYS - System reliability.

PSYS - System profit.

RUNIT - Vector. Stores unit reliabilities.

WUNIT - Vector. Stores unit weights.

LAMDA - Vector. Stores unit failure rates.

G - Vector. Contains module reliabilities.

AGRADN - Vector. Contains sensitivity of system availability to the module redundancies.

AGRADT - Sensitivity of system availability to the maintenance interval.

PGRADN - Vector. Contains sensitivity of system profit to the module redundancies.

PGRADT - Sensitivity of system profit to the maintenance interval.

RGRADN - Vector. Contains sensitivity of system reliability to the module redundancies.

RGRADT - Sensitivity of system reliability to the maintenance interval.

RIGRDN - Vector. Contains sensitivity of module reliability to its own redundancy.

MTGRDN - Vector. Stores sensitivity of MTBM to module redundancies.

COSTF - Vector. Stores unit costs.

CONSTR - Vector 10 locations in length. Stores values of the constraints.

MTBM - Mean time between maintenance.

MDT - Mean down time.

OBJFN - System objective function.

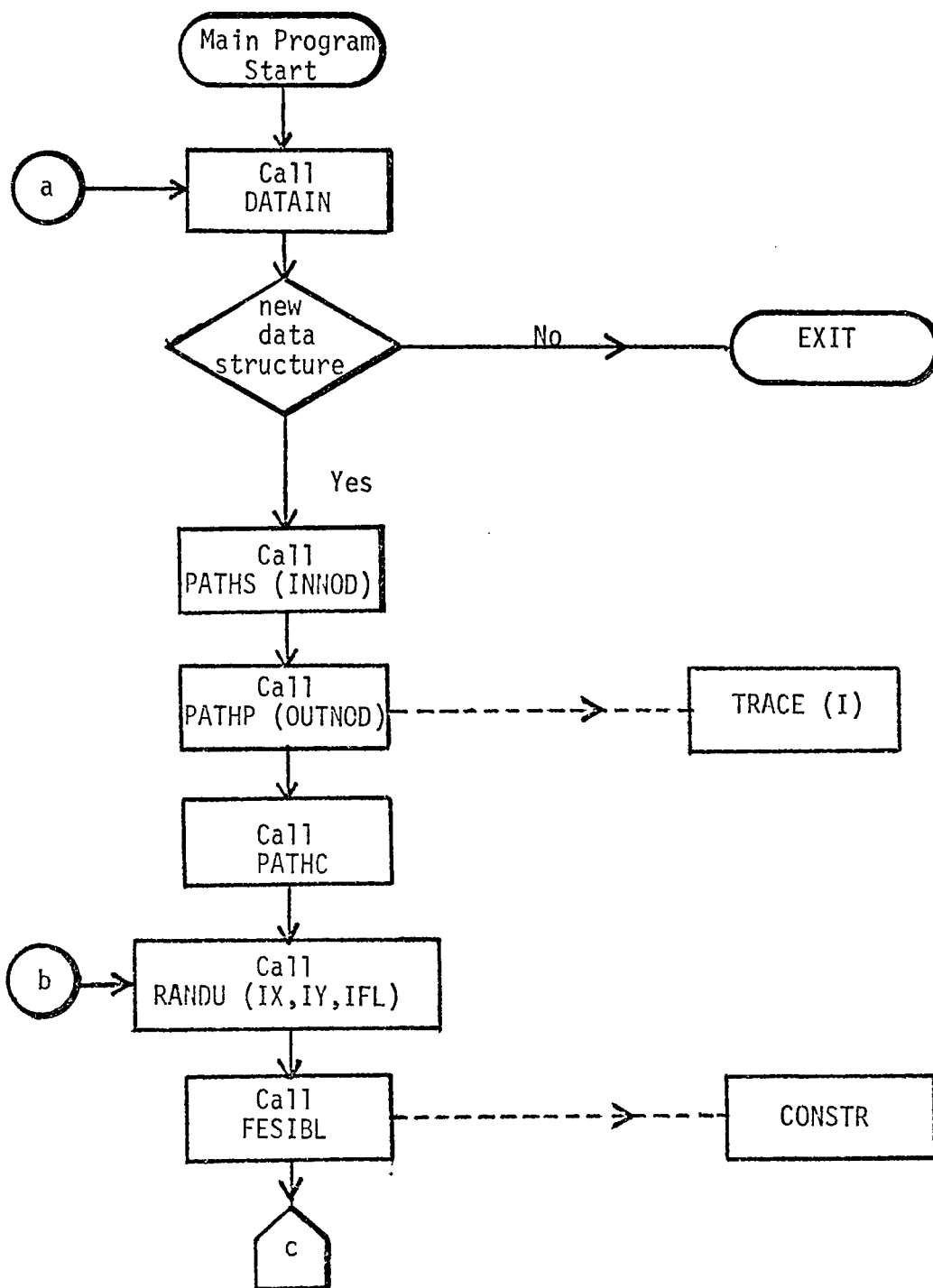
RLBTY - System reliability.

CMCOST - Upper limit on the value of cost constraint.

WEIGHT - Upper limit on the value of weight constraint.

INCOM - Net income per hour of system operation.

- COSTMF - Maintenance cost as a per cent of system cost per hour of system down time.
- T - Variable of integration.
- TEM - Emergency maintenance interval.
- TSC - Scheduled maintenance interval.
- TLF - System life time.
- TM - Maintenance interval.
- TMAX - Upper limit on the value of TM.
- TMIN - Lower limit on the value of TM.
- INTVL - Interval of integration used in the Simpson's formula.
- TSTEP - Constant to express maintenance interval in terms of (TM/TSTEP).
- FCOST - System fixed capital cost.
- MCOST - System maintenance cost.
- KOUNT - Number of random sample points.
- NONZRO - Number of nonzero elements in the system sensitivity vector.
- NSORT - Vector. Stores the module numbers in ascending order of magnitude of the system sensitivity to module redundancies.
- FLAG - It is either zero or one. When it is "one" only system availability and system profit are evaluated by subroutine AVLBTY. When it is "zero" sensitivity of system availability and system profit to module redundancies and maintenance interval are also evaluated.
- INDEX - A value equal to NMOD1 indicates that a feasible point could not be located.
- MOVE - As used in subroutine FESIBL - a value greater than 20 indicates that FESIBL point could not be located. As used in subroutines OPTIM and OPTIMZ - a value equal to zero indicates that an improved feasible solution is not possible for a particular gradient vector.



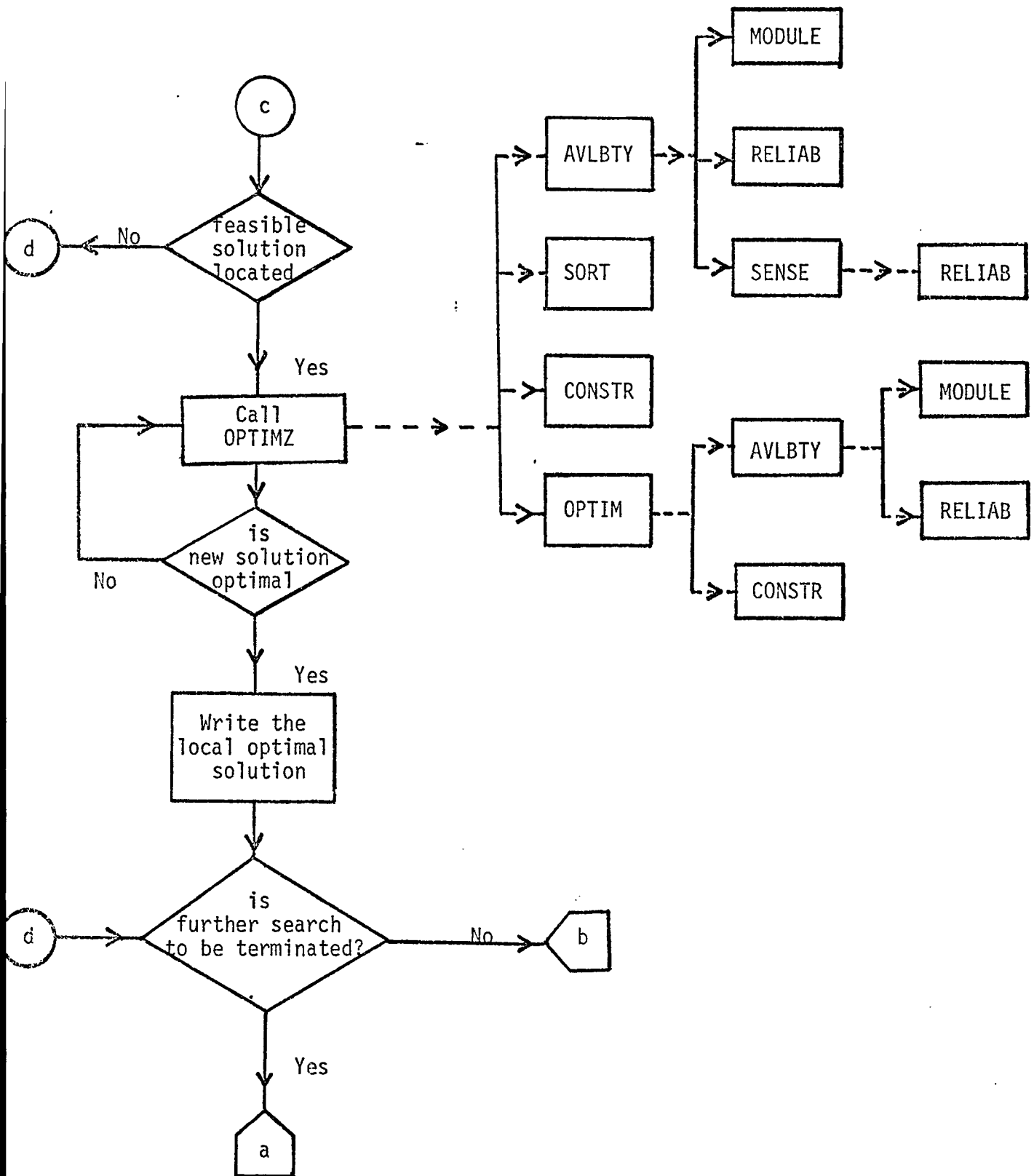


Figure G.2
Schematic Diagram of Computation I

---- indicates communication between subroutines.

Computational Scheme II

Computational Scheme II is based on the partial enumeration method of Lawler and Bell. Details of the method are given in the Appendix F and was used to solve the system reliability optimization problems discussed in Chapters III and IV of the text.

A schematic diagram of the main program is shown in Figure G-3. The computation is composed of seven subroutines. These include DATAIN, PATHS, PATHP, TRACE, PATHC, MODULE and RELIAB. A functional description of each has been presented with the first computational scheme discussed earlier in this appendix. Subroutines PATHS, PATHP, PATHC, TRACE are identical to the ones listed earlier.

PROGRAM VARIABLES

A list of major program variables and their explanations are given below:

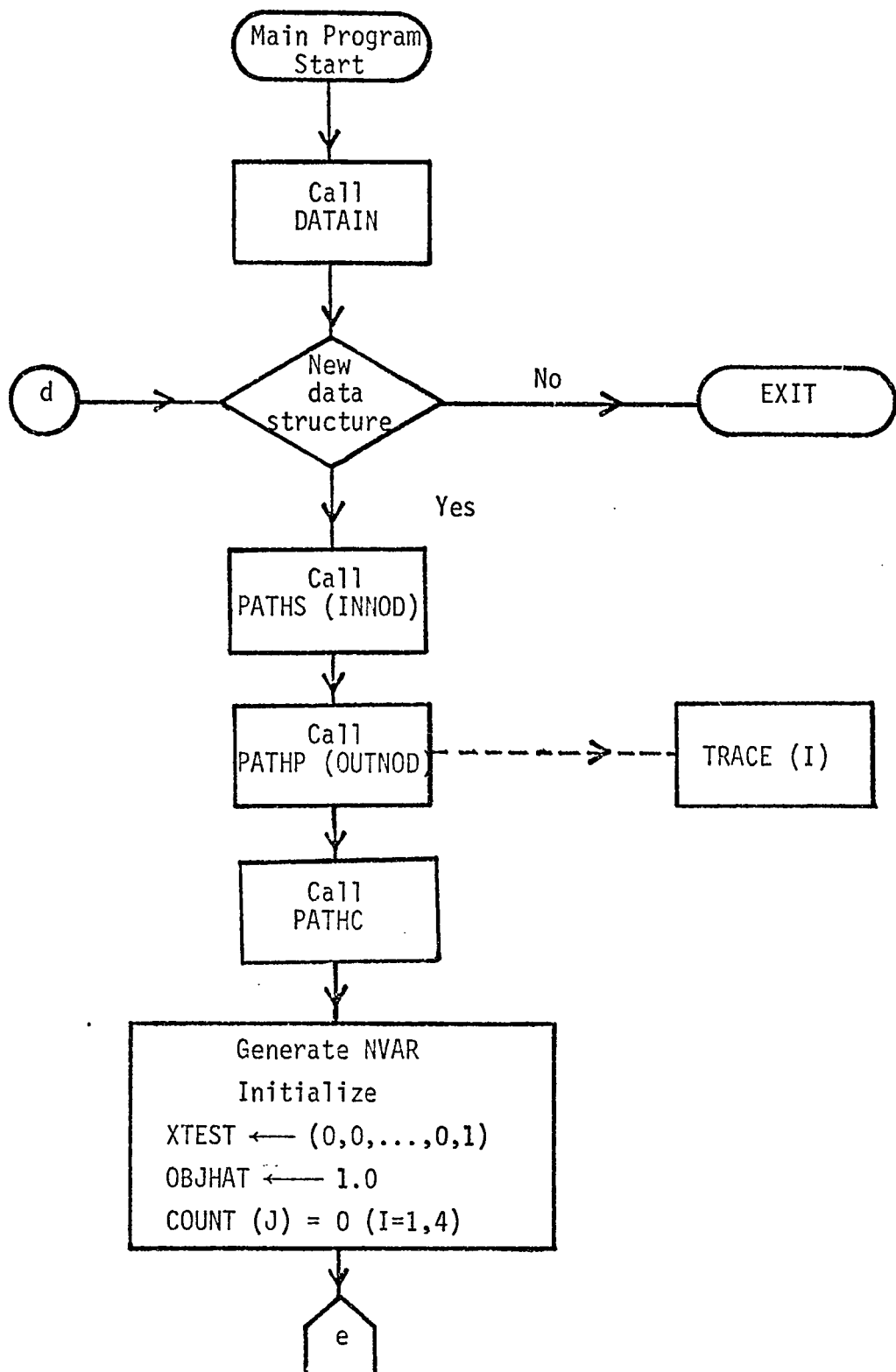
- RKODE(I) - Binary number zero or one. A "zero" identifies an active redundancy Ith module. A "one" identifies a standby redundancy Ith module.
- NKODE(I) - Number of binary variables associated with the number of units in the Ith module.
- MKODE(I) - Number of binary variables associated with the number of units necessary for successful operation of the Ith module.
- NTEST(I) - Number of units in the Ith module. Generated from the binary variables NXTEST (I,J).
- NSTAR1(I)- Number of units in the Ith module. Generated from the binary variables NXSTR1(I,J).
- MTEST(I) - Number of units in the Ith module that must operate for its successful operation. Generated from the binary variables MXTEST(I,J).
- MSTAR1(I)- Number of units in the Ith module that must operate for its successful operation. Generated from binary variable MXSTR1(I,J).

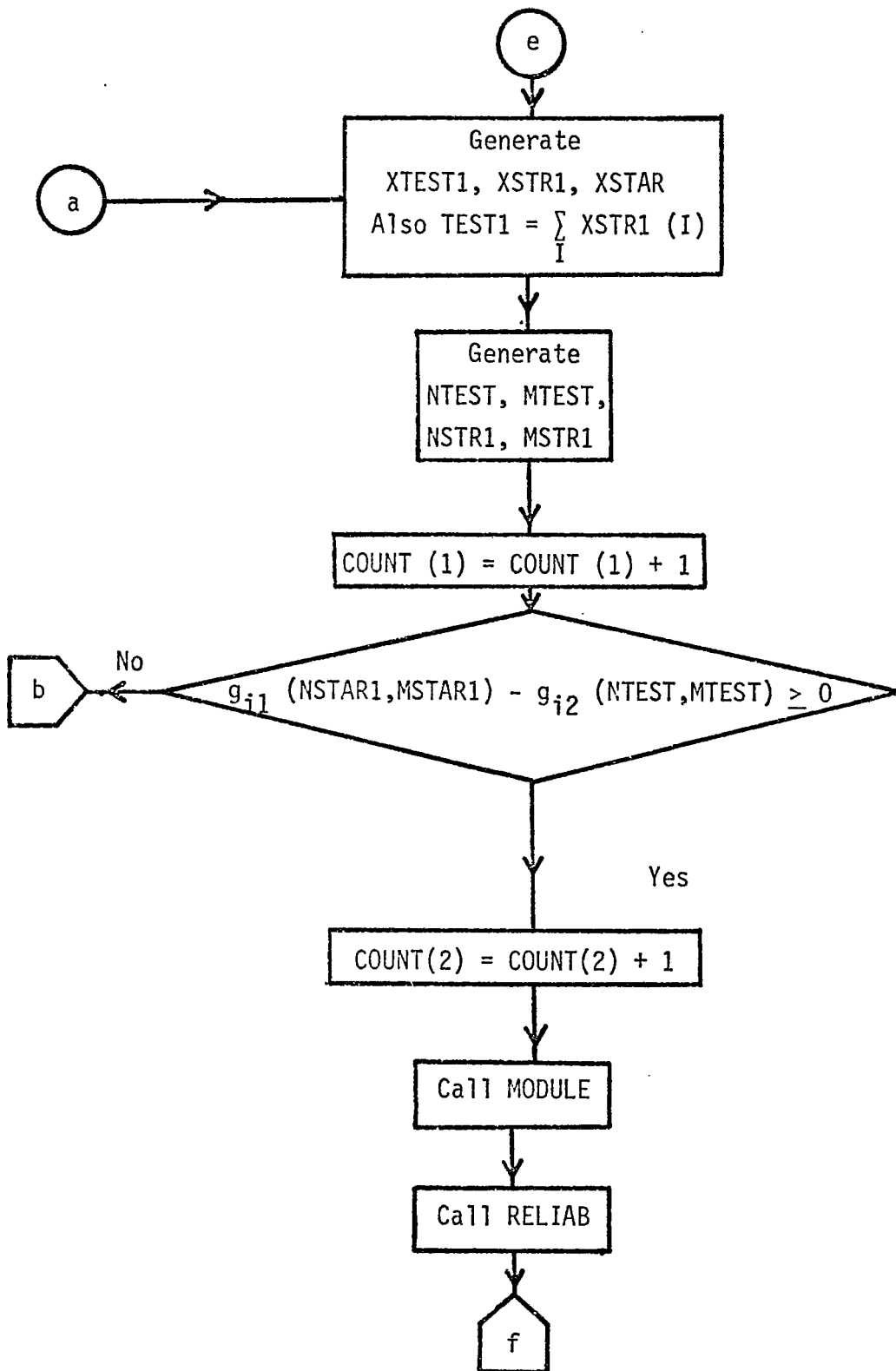
- NXTEST(I,J) - jth binary variable associated with the variable NTEST(I).
It is obtained from the vector XTEST.
- MXTEST(I,J) - jth binary variable associated with the variable MTEST(I).
It is obtained from the vector XTEST.
- NXSTR1(I,J) - jth binary variable associated with the variable NSTAR1(I).
It is obtained from the vector XSTR1.
- MXSTR1(I,J) - jth binary variable associated with the variable MSTAR1(I).
It is obtained from the vector XSTR1.
- XTEST - Vector containing NVAR binary system variables.
- XTEST1 - Vector obtained by the binary subtraction of (XTEST - XONE).
- XONE - Vector with the binary number 1 in the first position and
zero in all other positions.
- XSTR1 - Vector obtained by performing logical "OR" operation on
XTEST and XTEST1.
- XSTAR - Vector obtained by binary addition of (XSTR1 + XONE).
- XHAT - Vector. Stores the previous best value of vector XTEST.
- NHAT - Vector. Stores the previous best value of vector NTEST.
- MHAT - Vector. Stores the previous best value of vector MTEST.
- R(I) - Unit reliability of the Ith module.
- G(I) - Ith module reliability.
- GHAT(I) - Previous best value of G(I).
- KINDEX(I) - Constant " k_c " associated with the unit cost of the Ith
module.
- SINDEX(I) - Constant "s" associated with the cost of the Ith module.
- CONSTR(I) - Value of the Ith system constraint.
- COST - Upper limit on the total system cost.
- COSTHT - Total system cost corresponding to the system variables
NHAT and MHAT.
- RLBTY - System reliability.
- OBJFN - Given by $(1. - \text{RLBTY})$.
- OBJHAT - Previous best value of OBJFN.

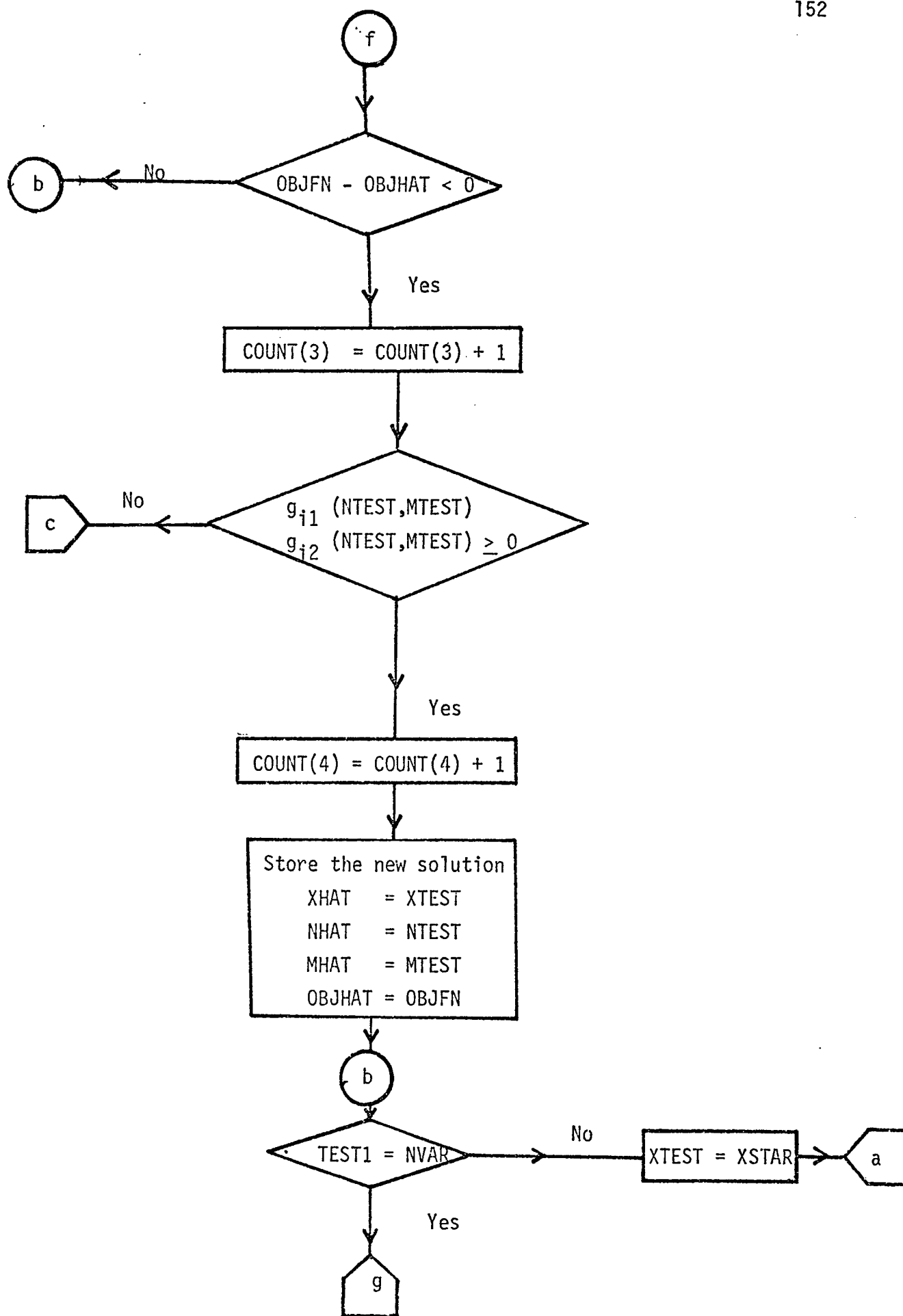
NVAR - Total number of binary system variables.

TEST1 and TEST2 - Variables to see if all the solution vectors have been examined.

COUNT - Vector of dimension four. Used for counting the number of operations at various stages of the computation.







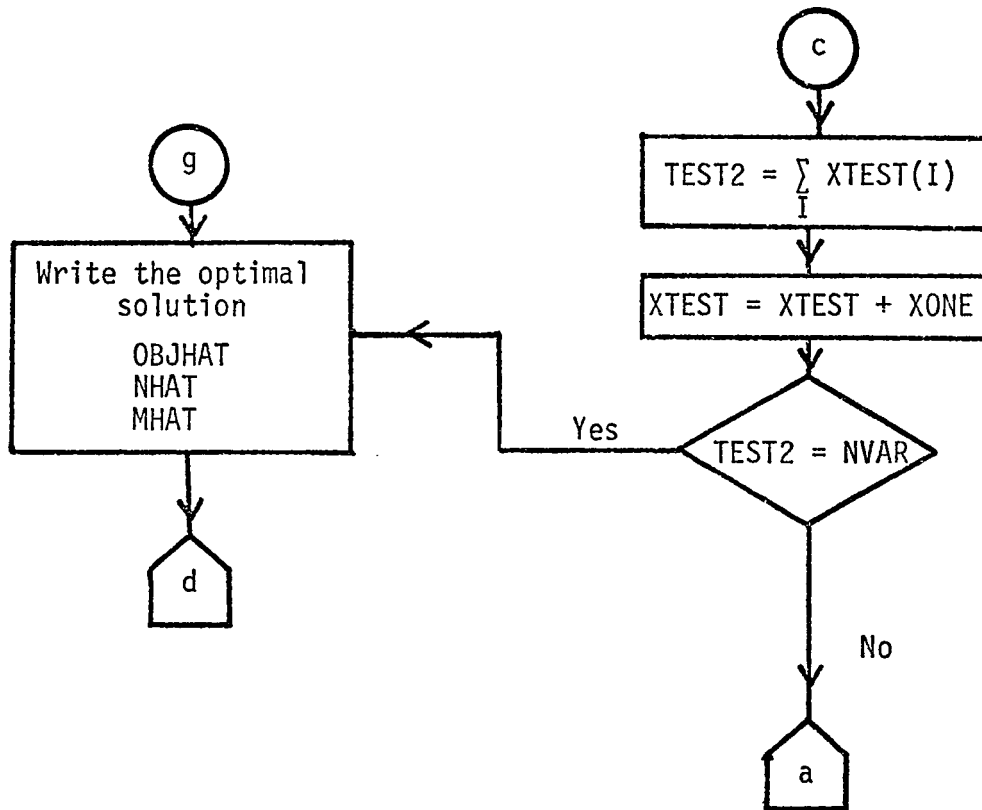


Figure G.3

Schematic diagram of Computation II

COMPUTER LISTING FOR COMPUTATION SCHEME I

SYSTEM RELIABILITY OPTIMIZATION

1:	C	MAIN PROGRAM	MAIN	10
2:	C	SYSTEM RELIABILITY OPTIMIZATION	MAIN	20
3:	C	MODIFIED INTEGER GRADIENT METHOD	MAIN	30
4:		IMPLICIT INTEGER*4 (A-Z)	MAIN	40
5:		COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MAIN	50
6:		1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MAIN	60
7:		2 PATH(300,50),NMOD,MPATH,PSIGN(300)	MAIN	70
8:		INTEGER*2 P,PP,PATH	MAIN	80
9:		COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MAIN	90
10:		1 NMIN(50),FLAG,INDEX,MOVE,ONZRO	MAIN	100
11:		COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	MAIN	110
12:		1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	MAIN	120
13:		REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	MAIN	130
14:		1 RLBTY,WEIGHT	MAIN	140
15:		REAL*4 YFL	MAIN	150
16:	C		MAIN	160
17:		PATHL = 300	MAIN	170
18:		PLENTH = 1000	MAIN	180
19:		PPL = 50	MAIN	190
20:		10 CALL DATAIN	MAIN	200
21:		READ(5,1000) INNOD,OUTNOD	MAIN	210
22:		CALL PATHS(INNOD)	MAIN	220
23:		IF(STATUS) 20,10,20	MAIN	230
24:		20 CALL PATHP(OUTNOD)	MAIN	240
25:		WRITE(6,1100)	MAIN	250
26:		DO 30 L=1,NPATH	MAIN	260
27:		WRITE(6,1200) L,(PATH(L,M),M=1,NMOD)	MAIN	270
28:		30 CONTINUE	MAIN	280
29:		CALL PATHC	MAIN	290
30:		WRITE(6,1001) NPATH,MPATH	MAIN	300
31:		WRITE(6,2600)	MAIN	310
32:		DO 35 L=1,MPATH	MAIN	320
33:		WRITE(6,2700) L,PSIGN(L),(PATH(L,M),M=1,NMOD)	MAIN	330
34:		35 CONTINUE	MAIN	340
35:		KOUNT = 1	MAIN	350

36:	ISEED = 984376521	MAIN 360
37:	IX = ISEED	MAIN 370
38:	DO 45 I=1,NMOD	MAIN 380
39:	45 NUNIT(I) = 1	MAIN 390
40:	GO TO 55	MAIN 400
41:	C GENERATING A UNIFORMLY DISTRIBUTED RANDOM STARTING POINT	MAIN 410
42:	40 DO 50 I=1,NMOD	MAIN 420
43:	CALL RANDU(IX,IY,YFL)	MAIN 430
44:	NUNIT(I) = YFL*(NMAX(I)- NMIN(I)+1) + NMIN(I)	MAIN 440
45:	IX = IY	MAIN 450
46:	50 CONTINUE	MAIN 460
47:	55 WRITE (6,1300) KOUNT	MAIN 470
48:	WRITE(6,2100)	MAIN 480
49:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 490
50:	C LOCATING A FEASIBLE POINT	MAIN 500
51:	WRITE(6,2200)	MAIN 510
52:	CALL FESIBL	MAIN 520
53:	WRITE(6,1410) MOVE,INDEX	MAIN 530
54:	IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD1) GO TO 90	MAIN 540
55:	GO TO 60	MAIN 550
56:	60 WRITE(6,2400)	MAIN 560
57:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 570
58:	WRITE(6,1600) CONST(1),CONST(2)	MAIN 580
59:	C SEARCHING FOR LOCAL OPTIMUM	MAIN 590
60:	WRITE(6,2300)	MAIN 600
61:	65 CALL OPTIMZ	MAIN 610
62:	IF(MOVE .EQ. 0) GO TO 70	MAIN 620
63:	GO TO 65	MAIN 630
64:	70 WRITE(6,2500)	MAIN 640
65:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 650
66:	CALL MODULE	MAIN 660
67:	CALL RELIAB	MAIN 670
68:	CALL CONSTR	MAIN 680
69:	WRITE(6,1500) RLBTY	MAIN 690
70:	WRITE(6,1600) CONST(1),CONST(2)	MAIN 700
71:	90 KOUNT = KOUNT + 1	MAIN 710

72:	IF (KOUNT .GT. 25) GO TO 10	MAIN 720
73:	GO TO 40	MAIN 730
74:	1000 FORMAT (2I3)	MAIN 740
75:	1001 FORMAT(1H0,/,30X,'NPATH = ',I3,/,30X,'MPATH = ',I5)	MAIN 750
76:	1100 FORMAT (1H0,/,T20,'PATHS IN BINARY CODE',/)	MAIN 760
77:	1200 FORMAT (1H0,T16,'PATH',I3,' = ',50I2)	MAIN 770
78:	1300 FORMAT(1H1,40X,'KOUNT = ',I3,//)	MAIN 780
79:	1400 FORMAT(1H0,/,10X,'NUNIT(1) = ',20I5)	MAIN 790
80:	1410 FORMAT (1H0,/,30X,'MOVE = ',I4,10X,'INDEX = ',I4)	MAIN 800
81:	1500 FORMAT (1H0,/,30X,'SYSTEM RELIABILITY = ',G13.6)	MAIN 810
82:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) = ',G13.6,/,30X,'CONSTRAINT(2) = ',	MAIN 820
83:	1 G13.6)	MAIN 830
84:	2100 FORMAT(1H0,50X,'THE INITIAL POINT IS')	MAIN 840
85:	2200 FURMAT(1H0,50X,'SEARCHING FOR THE FEASIBLE POINT')	MAIN 850
86:	2300 FORMAT(1H0,50X,'SEARCHING FOR THE LOCAL OPTIMUM POINT')	MAIN 860
87:	2400 FORMAT(1H0,////,50X,'THE FEASIBLE POINT IS')	MAIN 870
88:	2500 FORMAT(1H0,////,50X,'THE OPTIMAL POINT IS')	MAIN 880
89:	2600 FORMAT(1H0,/,T20,'PATHS AND PATH COMBINATIONS IN BINARY CODE',/)	MAIN 890
90:	2700 FORMAT(1H0,T16,'PATHC',I3,' = ',I2,4X,50I2)	MAIN 900
91:	END	MAIN 910

1:	SUBROUTINE DATAIN	DATA	10
2:	C THIS SUBROUTINE READS AND ECHOCHECKS THE DATA	DATA	20
3:	IMPLICIT INTEGER*4 (A-Z)	DATA	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	DATA	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	DATA	50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	DATA	60
7:	INTEGER*2 P,PP,PATH	DATA	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	DATA	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	DATA	90
10:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	DATA	100
11:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	DATA	110
12:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	DATA	120
13:	1 RLBTY,WEIGHT	DATA	130
14:	DIMENSION TITLE(30)	DATA	140
15:	READ(5,999,END=100)(TITLE(I),I=1,30)	DATA	150
16:	READ(5,1000) NNODES,NMOD,CMCOST,WEIGHT	DATA	160
17:	DO 10 I=1,NMOD	DATA	170
18:	10 READ(5,1007) KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),WUNIT(I),	DATA	180
19:	1 RUNIT(I)	DATA	190
20:	C	DATA	200
21:	FOR=0	DATA	210
22:	REV=0	DATA	220
23:	DO 40 I=1,1000	DATA	230
24:	READ(5,1001)P(I,1),P(I,2),P(I,4)	DATA	240
25:	IF(P(I,1)-P(I,2)) 20,50,30	DATA	250
26:	20 FOR=FOR + 1	DATA	260
27:	GO TO 40	DATA	270
28:	30 REV=REV + 1	DATA	280
29:	40 CONTINUE	DATA	290
30:	50 NEDGES=FOR + REV	DATA	300
31:	NEDGS1=NEDGES + 1	DATA	310
32:	FREV=FOR + 1	DATA	320
33:	WRITE(6,1002) (TITLE(I),I=1,30)	DATA	330
34:	WRITE(6,1003)	DATA	340
35:	WRITE(6,1004)	DATA	350

36:	WRITE(6,1003)	DATA 360
37:	DO 60 I=1,NEDGES	DATA 370
38:	WRITE(6,1005) I,P(I,1),P(I,2),P(I,4)	DATA 380
39:	WRITE(6,1003)	DATA 390
40:	60 CONTINUE	DATA 400
41:	WRITE(6,1008)	DATA 410
42:	WRITE(6,1011)	DATA 420
43:	WRITE(6,1009)	DATA 430
44:	WRITE(6,1011)	DATA 440
45:	DO 70 I=1,NMOD	DATA 450
46:	WRITE(6,1010) I,KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),	DATA 460
47:	1 WUNIT(I),RUNIT(I)	DATA 470
48:	WRITE(6,1011)	DATA 480
49:	70 CONTINUE	DATA 490
50:	WRITE(6,1006) CMCOST,WEIGHT	DATA 500
51:	999 FORMAT(15A4)	DATA 510
52:	1000 FORMAT (2I10,2F10.0)	DATA 520
53:	1001 FORMAT (3I10)	DATA 530
54:	1002 FORMAT(1H1,/,T10,30A4,/,T50,'INVENTORY OF BRANCHES',/)	DATA 540
55:	1003 FORMAT (1H+,T38,44(' _'))	DATA 550
56:	1004 FORMAT (1H ,T38,' BRANCH ', ' ORIGIN ', ' END ',	DATA 560
57:	1 ' MODULE ')	DATA 570
58:	1005 FORMAT (1H ,T38,4(' ',4X,I2,4X),' ')	DATA 580
59:	1006 FORMAT (1H0,T10,'CMCOST=',G13.6,/,T10,'WEIGHT=',G13.6)	DATA 590
60:	1007 FORMAT(4I10,3F10.0)	DATA 600
61:	1008 FORMAT(1H0,T50,'INVENTORY OF MODULES',/)	DATA 610
62:	1009 FORMAT(1H ,T20,' MODULE ', ' KODE ', ' MAX-UNIT ',	DATA 620
63:	1 ' MIN-UNIT ', 'M OUT OF N ', 'UNIT COSTF ', 'UNIT WEGHT ',	DATA 630
64:	1 ' RELIABILITY ')	DATA 640
65:	1010 FORMAT(1H ,T20,5(' ',4X,I2,4X),2(' ',F10.3),' ',G13.6,' ')	DATA 650
66:	1011 FORMAT(1H+,T20,91(' _'))	DATA 660
67:	RETURN	DATA 670
68:	100 CALL EXIT	DATA 680
69:	END	DATA 690

1:	SUBROUTINE PATHS(INPUT)	PATH 10
2:	C THIS SUBROUTINE ACCEPTS AN INPUT NODE AND FINDS ALL PATHS	PATH 20
3:	C STARTING AT THIS NODE	PATH 30
4:	IMPLICIT INTEGER*4 (A-Z)	PATH 40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	PATH 50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	PATH 60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	PATH 70
8:	INTEGER*2 P,PP,PATH	PATH 80
9:	C	PATH 90
10:	DO 5 I=1,NNODES	PATH 100
11:	5 PP(I)=0	PATH 110
12:	PP(INPUT)=1	PATH 120
13:	7 A1=NEDGS1	PATH 130
14:	CYCLE=0	PATH 140
15:	STATUS=1	PATH 150
16:	MARK=0	PATH 160
17:	NSTART=A1	PATH 170
18:	10 CYCLE=CYCLE + 1	PATH 180
19:	DO 100 I=1,NEDGES	PATH 190
20:	IF((I.NE.1).AND.(I.NE.FREV)) GO TO 20	PATH 200
21:	IF(MARK.LT.0) GO TO 120	PATH 210
22:	ISTART=NSTART	PATH 220
23:	NSTART=A1	PATH 230
24:	MARK=-1	PATH 240
25:	20 FIN=A1 - 1	PATH 250
26:	IF((CYCLE.GT.1).OR.(I.GT.FOR)) GO TO 30	PATH 260
27:	J=P(I,1)	PATH 270
28:	IF(PP(J).EQ.0) GO TO 30	PATH 280
29:	P(A1,1)=P(I,1)	PATH 290
30:	P(A1,2)=P(I,2)	PATH 300
31:	P(A1,3)=A1	PATH 310
32:	P(A1,4)=P(I,4)	PATH 320
33:	A1=A1 + 1	PATH 330
34:	MARK=0	PATH 340
35:	30 A2=ISTART	PATH 350

36:	40	IF(A2.GT.FIN) GO TO 100	PATH 360
37:		IF(P(A2,2).NE.P(I,1)) GO TO 90	PATH 370
38:		IF(P(I,2)-P(A2,1)) 90,90,50	PATH 380
39:	50	IF((CYCLE.EQ.1).AND.(I.LE.FOR)) GO TO 80	PATH 390
40:		A3=P(A2,3)	PATH 400
41:	60	IF(P(A3,2).EQ.P(I,2)) GO TO 90	PATH 410
42:		IF(A3.EQ.P(A3,3)) GO TO 80	PATH 420
43:		A3=P(A3,3)	PATH 430
44:		GO TO 60	PATH 440
45:	80	MARK=0	PATH 450
46:		P(A1,1)=P(A2,1)	PATH 460
47:		P(A1,2)=P(I,2)	PATH 470
48:		P(A1,3)=A2	PATH 480
49:		P(A1,4)=P(I,4)	PATH 490
50:		A1=A1 + 1	PATH 500
51:		IF(A1.GT.PLENTH) GO TO 110	PATH 510
52:	90	A2=A2 + 1	PATH 520
53:		GO TO 40	PATH 530
54:	100	CONTINUE	PATH 540
55:		IF(REV) 120,120,10	PATH 550
56:	110	WRITE(6,1000)	PATH 560
57:		STATUS=0	PATH 570
58:	120	LEND=A1 - 1	PATH 580
59:		RETURN	PATH 590
60:	1000	FORMAT(1H1,35('*'),/, ' DIAGNOSTIC - LONGER P ARRAY REQUIRED',	PATH 600
61:		1/, ' ',35('*'))	PATH 610
62:		END	PATH 620

1:	SUBROUTINE PATHP(OUTPUT)	PATH	10
2:	C THIS SUBROUTINE ITEMIZES ALL THE PATHS STARTING AT INPUT (DEFINED	PATH	20
3:	C BY A CALL TO PATHS) AND ENDING AT THE OUTPUT NODE SPECIFIED.	PATH	30
4:	IMPLICIT INTEGER*4 (A-Z)	PATH	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	PATH	50
6:	1 NPATH,PATHL,PLENGTH,PPL,REV,STATUS,P(1000,4),PP(50),	PATH	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	PATH	70
8:	INTEGER*2 P,PP,PATH	PATH	80
9:	C	PATH	90
10:	WRITE(6,1001) P(NEDGS1,1),OUTPUT	PATH	100
11:	WRITE(6,1002)	PATH	110
12:	WRITE(6,1003)	PATH	120
13:	WRITE(6,1002)	PATH	130
14:	NPATH=0	PATH	140
15:	DO 350 I=NEDGS1,LEND	PATH	150
16:	IF((P(I,2).EQ.OUTPUT)) GO TO 340	PATH	160
17:	GO TO 350	PATH	170
18:	340 CALL TRACE(I)	PATH	180
19:	WRITE(6,1002)	PATH	190
20:	350 CONTINUE	PATH	200
21:	RETURN	PATH	210
22:	1001 FORMAT('1',T35,'INVENTORY OF PATHS',//,T21,'THE FOLLOWING PATHS EX	PATH	220
23:	11ST BETWEEN NODES',I3,' AND',I3,//)	PATH	230
24:	1002 FORMAT(1H+,T16,62(' '))	PATH	240
25:	1003 FORMAT(' ',T16,' ROW OF P TRACE --> NODES TRACE -->	PATH	250
26:	1 MODULES ')	PATH	260
27:	END	PATH	270

1:	SUBROUTINE TRACE(I)	TRAC	10
2:	C THIS SUBROUTINE PRINTS THE ELEMENTS ASSOCIATED WITH THE	TRAC	20
3:	C ITH ENTRY IN THE P-ARRAY.	TRAC	30
4:	IMPLICIT INTEGER*4 (A-Z)	TRAC	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	TRAC	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	TRAC	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	TRAC	70
8:	INTEGER*2 P,PP,PATH	TRAC	80
9:	INTEGER*2 PI4(50)	TRAC	90
10:	C	TRAC	100
11:	J=I	TRAC	110
12:	K=PPL	TRAC	120
13:	10 PP(K)=P(J,2)	TRAC	130
14:	PI4(K)=P(J,4)	TRAC	140
15:	K=K - 1	TRAC	150
16:	IF(P(J,3).EQ.J) GO TO 20	TRAC	160
17:	J=P(J,3)	TRAC	170
18:	GO TO 10	TRAC	180
19:	20 PP(K)=P(J,1)	TRAC	190
20:	FL=K	TRAC	200
21:	30 NN=PPL - FL + 1	TRAC	210
22:	IF(NN.EQ.0) GO TO 40	TRAC	220
23:	IF(NN.GT.4) NN=4	TRAC	230
24:	LL=FL + NN - 1	TRAC	240
25:	WRITE(6,1000)(PP(J),J=FL,LL)	TRAC	250
26:	IF(FL.EQ.K) GO TO 35	TRAC	260
27:	WRITE(6,1001) (PI4(J),J=FL,LL)	TRAC	270
28:	GO TO 45	TRAC	280
29:	35 FL1=FL+1	TRAC	290
30:	WRITE(6,1002) I	TRAC	300
31:	WRITE(6,1001) (PI4(J),J=FL1,LL)	TRAC	310
32:	45 FL=LL+1	TRAC	320
33:	GO TO 30	TRAC	330
34:	40 NPATH=NPATH+1	TRAC	340
35:	DO 50 J=1,NMOD	TRAC	350

36:	50	PATH(NPATH,J) = 0	TRAC 360
37:		K1 = K + 1	TRAC 370
38:		DO 60 J=K1,PPL	TRAC 380
39:		MM = PI4(J)	TRAC 390
40:	60	PATH(NPATH,MM) = 1	TRAC 400
41:		RETURN	TRAC 410
42:	1000	FORMAT(' ',T16,' ',T27,' ',T52,' ',T28,4(I3,'-->'))	TRAC 420
43:	1001	FORMAT('+',T77,' ',T53,4(I3,'-->'))	TRAC 430
44:	1002	FORMAT('+',T20,I4)	TRAC 440
45:		END	TRAC 450

1:	SUBROUTINE PATHC	PATH	10
2:	C THIS SUBROUTINE FINDS ALL THE PATH COMBINATIONS IN	PATH	20
3:	C THEIR BINARY FORM.	PATH	30
4:	IMPLICIT INTEGER*4 (A-Z)	PATH	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	PATH	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	PATH	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	PATH	70
8:	INTEGER*2 P,PP,PATH	PATH	80
9:	INTEGER*2 PTAIL(300)	PATH	90
10:	C	PATH	100
11:	DO 10 I=1,NPATH	PATH	110
12:	PTAIL(I) = I	PATH	120
13:	PSIGN(I) = +1	PATH	130
14:	10 CONTINUE	PATH	140
15:	N1 = 1	PATH	150
16:	N2 = NPATH	PATH	160
17:	50 IF (PTAIL(N1) .EQ. NPATH) GO TO 100	PATH	170
18:	K1 = PTAIL(N1) + 1	PATH	180
19:	DO 20 I=K1,NPATH	PATH	190
20:	N2 = N2 + 1	PATH	200
21:	IF (N2 .GT. PATHL) GO TO 200	PATH	210
22:	PTAIL(N2) = I	PATH	220
23:	PSIGN(N2) = -PSIGN(N1)	PATH	230
24:	INDEX = 0	PATH	240
25:	DO 30 J=1,NMOD	PATH	250
26:	PATH(N2,J) = 0	PATH	260
27:	IF(PATH(N1,J).EQ.1 .OR. PATH(I,J).EQ.1) GO TO 40	PATH	270
28:	GO TO 30	PATH	280
29:	40 PATH(N2,J) = 1	PATH	290
30:	INDEX = INDEX + 1	PATH	300
31:	30 CONTINUE	PATH	310
32:	IF(I.EQ.NPATH) GO TO 20	PATH	320
33:	IF (INDEX .EQ. NMOD) N2 = N2-1	PATH	330
34:	20 CONTINUE	PATH	340
35:	100 N1 = N1 + 1	PATH	350

36:	IF (N1 .EQ. N2) GO TO 500	PATH 360
37:	GO TO 50	PATH 370
38:	200 WRITE(6,1000)	PATH 380
39:	STATUS = 0	PATH 390
40:	500 MPATH = N2	PATH 400
41:	RETURN	PATH 410
42:	1000 FORMAT(1H1,39('*'),/, 'DIAGONOSTIC-LONGER PATH ARRAY REQUIRED',/,	PATH 420
43:	1 ' ',39('*'))	PATH 430
44:	END	PATH 440

1:	SUBROUTINE MODULE	MODU 10
2:	C THIS SUBROUTINE COMPUTES THE MODULE RELIABILITY AND	MODU 20
3:	C SENSITIVITY OF MODULE RELIABILITY TO THE MODULE REDUNDANCY.	MODU 30
4:	IMPLICIT INTEGER*4 (A-Z)	MODU 40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MODU 50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MODU 60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	MODU 70
8:	INTEGER*2 P,PP,PATH	MODU 80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MODU 90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	MODU 100
11:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	MODU 110
12:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	MODU 120
13:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	MODU 130
14:	1 RLBTY,WEIGHT	MODU 140
15:	REAL*4 DUMMY1,DUMMY2,DUMMY3	MODU 150
16:	C	MODU 160
17:	DO 200 K=1,NMOD	MODU 170
18:	IF(KODE(K) .EQ. 0) GO TO 100	MODU 180
19:	DUMMY1 = RUNIT(K) ** MUNIT(K)	MODU 190
20:	DUMMY3 = - MUNIT(K) * ALOG(RUNIT(K))	MODU 200
21:	G(K) = 1.0	MODU 210
22:	DUMMY2 = 1.0	MODU 220
23:	NM = NUNIT(K) - MUNIT(K)	MODU 230
24:	IF(NM .EQ. 0) GO TO 20	MODU 240
25:	DO 10 I=1,NM	MODU 250
26:	DUMMY2 = DUMMY2 * DUMMY3/I	MODU 260
27:	G(K) = G(K) + DUMMY2	MODU 270
28:	10 CONTINUE	MODU 280
29:	20 G(K) = G(K) * DUMMY1	MODU 290
30:	RIGRDN(K) = DUMMY1 * DUMMY2 * DUMMY3/(NM+1)	MODU 300
31:	GO TO 200	MODU 310
32:	100 DUMMY1 = RUNIT(K)	MODU 320
33:	DUMMY2 = DUMMY1**NUNIT(K)	MODU 330
34:	G(K) = DUMMY2	MODU 340
35:	NM = NUNIT(K) - MUNIT(K)	MODU 350

36:	IF(NM .EQ. 0) GO TO 120	MODU 360
37:	DO 110 I=1,NM	MODU 370
38:	DUMMY2 = DUMMY2 * ((1.-DUMMY1)/DUMMY1) * (NUNIT(K)-I+1)/I	MODU 380
39:	G(K) = G(K) + DUMMY2	MODU 390
40:	110 CONTINUE	MODU 400
41:	120 RIGRDN(K) = DUMMY2 * (1.-DUMMY1) * MUNIT(K)/(NM+1)	MODU 410
42:	200 CONTINUE	MODU 420
43:	RETURN	MODU 430
44:	END	MODU 440

1:	SUBROUTINE RELIAB	RELI	10
2:	C THIS SUBROUTINE CALCULATES THE SYSTEM RELIABILITY FROM THE	RELI	20
3:	C MODULE RELIABILITIES.	RELI	30
4:	IMPLICIT INTEGER*4 (A-Z)	RELI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	RELI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	RELI	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	RELI	70
8:	INTEGER*2 P,PP,PATH	RELI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	RELI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	RELI	100
11:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	RELI	110
12:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	RELI	120
13:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	RELI	130
14:	1 RLBTY,WEIGHT	RELI	140
15:	REAL*4 GPATH	RELI	150
16:	C	RELI	160
17:	RLBTY = 0.0	RELI	170
18:	DO 10 I=1,MPATH	RELI	180
19:	GPATH = PSIGN(I)	RELI	190
20:	DO 20 J=1,NMOD	RELI	200
21:	IF (PATH(I,J) .EQ. 0) GO TO 20	RELI	210
22:	GPATH = GPATH * G(J)	RELI	220
23:	20 CONTINUE	RELI	230
24:	RLBTY = RLBTY + GPATH	RELI	240
25:	10 CONTINUE	RELI	250
26:	RETURN	RELI	260
27:	END	RELI	270

1:	SUBROUTINE SENSE	SENS	10
2:	C THIS SUBROUTINE DETERMINES THE SENSITIVITY OF THE SYSTEM	SENS	20
3:	C RELIABILITY TO MODULE REDUNDANCIES.	SENS	30
4:	IMPLICIT INTEGER*4 (A-Z)	SENS	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SENS	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SENS	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	SENS	70
8:	INTEGER*2 P,PP,PATH	SENS	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SENS	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SENS	100
11:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	SENS	110
12:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	SENS	120
13:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	SENS	130
14:	1 RLBTY,WEIGHT	SENS	140
15:	REAL*4 DUMMY(50)	SENS	150
16:	C	SENS	160
17:	DO 10 I=1,NMOD	SENS	170
18:	10 DUMMY(I) = G(I)	SENS	180
19:	DO 100 I=1,NMOD	SENS	190
20:	DO 90 J=1,NMOD	SENS	200
21:	90 G(J) = DUMMY(J)	SENS	210
22:	G(I) = 1.0	SENS	220
23:	CALL RELIAB	SENS	230
24:	RGRADN(I) = RLBTY	SENS	240
25:	G(I) = 0.0	SENS	250
26:	CALL RELIAB	SENS	260
27:	RGRADN(I) = (RGRADN(I)-RLBTY) * RIGRDN(I)	SENS	270
28:	100 CONTINUE	SENS	280
29:	DO 150 I=1,NMOD	SENS	290
30:	150 G(I) = DUMMY(I)	SENS	300
31:	RETURN	SENS	310
32:	END	SENS	320

1:	SUBROUTINE CONSTR	CONS	10
2:	C THIS SUBROUTINE COMPUTES THE SYSTEM CONSTRAINTS.	CONS	20
3:	IMPLICIT INTEGER*4 (A-Z)	CONS	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	CONS	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	CONS	50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	CONS	60
7:	INTEGER*2 P,PP,PATH	CONS	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	CONS	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	CONS	90
10:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	CONS	100
11:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	CONS	110
12:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	CONS	120
13:	1 RLBTY,WEIGHT	CONS	130
14:	C	CONS	140
15:	CONST(1) = 0.0	CONS	150
16:	CONST(2) = 0.0	CONS	160
17:	DO 10 I=1,NMOD	CONS	170
18:	CONST(1) = CONST(1) + NUNIT(I)*COSTF(I)	CONS	180
19:	CONST(2) = CONST(2) + NUNIT(I)*WUNIT(I)	CONS	190
20:	10 CONTINUE	CONS	200
21:	CONST(1) = CONST(1) - CMCOST	CONS	210
22:	CONST(2) = CONST(2) - WEIGHT	CONS	220
23:	RETURN	CONS	230
24:	END	CONS	240

1:	SUBROUTINE FESIBL	FESI	10
2:	C THIS SUBROUTINE FINDS THE FEASIBLE SOLUTION BY THE METHOD	FESI	20
3:	C OF WEIGHTED PERPENDICULARS.	FESI	30
4:	IMPLICIT INTEGER*4 (A-Z)	FESI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	FESI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	FESI	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	FESI	70
8:	INTEGER*2 P,PP,PATH	FESI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	FESI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	FESI	100
11:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	FESI	110
12:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	FESI	120
13:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	FESI	130
14:	1 RLBTY,WEIGHT	FESI	140
15:	REAL*4 Z(50),ZMIN,S,DUMMY1,ZMAX	FESI	150
16:	C	FESI	160
17:	MOVE = 0	FESI	170
18:	INDEX = 0	FESI	180
19:	5 CALL CONSTR	FESI	190
20:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	FESI	200
21:	WRITE(6,1600) CONST(1),CONST(2)	FESI	210
22:	IF(CONST(1) .LE. 0.0 .AND. CONST(2) .LE. 0.0) GO TO 1000	FESI	220
23:	IF(CONST(1) .LE. 0.0) GO TO 30	FESI	230
24:	S = CONST(1)	FESI	240
25:	DUMMY1 = 0.0	FESI	250
26:	DO 10 I=1,NMOD	FESI	260
27:	10 DUMMY1 = DUMMY1 + NUNIT(I)*COSTF(I)	FESI	270
28:	DO 20 I=1,NMOD	FESI	280
29:	20 Z(I) = -CONST(1) * COSTF(I)	FESI	290
30:	30 IF(CONST(2) .LE. 0.0) GO TO 50	FESI	300
31:	S = S + CONST(2)	FESI	310
32:	DO 40 I=1,NMOD	FESI	320
33:	40 Z(I) = Z(I) - CONST(2) * WUNIT(I)	FESI	330
34:	50 DO 60 I=1,NMOD	FESI	340
35:	60 Z(I) = +Z(I)/S	FESI	350

36:	ZMAX = ABS(Z(1))	FESI 360
37:	DO 70 J=2,NMOD	FESI 370
38:	IF (ZMAX .LT. ABS(Z(J))) ZMAX = ABS(Z(J))	FESI 380
39:	70 CONTINUE	FESI 390
40:	DO 80 I=1,NMOD	FESI 400
41:	IF(Z(I)) 90,80,100	FESI 410
42:	90 Z(I) = (Z(I)/ZMAX) * 2. - 0.5	FESI 420
43:	Z(I) = AINT(Z(I))	FESI 430
44:	GO TO 80	FESI 440
45:	100 Z(I) = (Z(I)/ZMAX) * 2. + 0.5	FESI 450
46:	Z(I) = AINT(Z(I))	FESI 460
47:	80 CONTINUE	FESI 470
48:	INDEX = 0	FESI 480
49:	MOVE = MOVE + 1	FESI 490
50:	DO 120 I=1,NMOD	FESI 500
51:	NUNIT(I) = NUNIT(I) + Z(I)	FESI 510
52:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) INDEX=INDEX+1	FESI 520
53:	IF(NUNIT(I) .LT. NMIN(I)) NUNIT(I) = NMIN(I)	FESI 530
54:	IF(NUNIT(I) .GT. NMAX(I)) NUNIT(I) = NMAX(I)	FESI 540
55:	120 CONTINUE	FESI 550
56:	IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD) GO TO 999	FESI 560
57:	GO TO 5	FESI 570
58:	999 WRITE(6,2500)	FESI 580
59:	1000 RETURN	FESI 590
60:	1400 FORMAT(1H0,/,10X,'NUNIT(I) =',20I5)	FESI 600
61:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	FESI 610
62:	1 G13.6)	FESI 620
63:	2500 FORMAT(1H0,////,30X,'FESIBL POINT CAN NOT BE LOCATED',////)	FESI 630
64:	END	FESI 640

1:	SUBROUTINE SORT	SORT 10
2:	C THIS SUBROUTINE ARRANGES THE GRADIENT VECTOR IN ASCENDING ORDER.	SORT 20
3:	IMPLICIT INTEGER*4 (A-Z)	SORT 30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SORT 40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SORT 50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	SORT 60
7:	INTEGER*2 P,PP,PATH	SORT 70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SORT 80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SORT 90
10:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	SORT 100
11:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	SORT 110
12:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	SORT 120
13:	1 RLBTY,WEIGHT	SORT 130
14:	REAL*4 GRDMIN,DUMMY(50)	SORT 140
15:	C	SORT 150
16:	DO 5 I=1,NMOD	SORT 160
17:	DUMMY(I) = RGRADN(I)	SORT 170
18:	5 CONTINUE	SORT 180
19:	NONZRO = 0	SORT 190
20:	DO 10 I=1,NMOD	SORT 200
21:	IF(ABS(DUMMY(I))) .LT. 1.E-15) GO TO 10	SORT 210
22:	NONZRO = NONZRO + 1	SORT 220
23:	10 CONTINUE	SORT 230
24:	DO 20 I=1,NONZRO	SORT 240
25:	GRDMIN = 1.E30	SORT 250
26:	DO 30 J=1,NMOD	SORT 260
27:	IF(ABS(DUMMY(J))) .LT. 1.E-15) GO TO 30	SORT 270
28:	I1 = I - 1	SORT 280
29:	IF(I1 .EQ.0) GO TO 50	SORT 290
30:	DO 40 K=1,I1	SORT 300
31:	IF(J .EQ. NSORT(K)) GO TO 30	SORT 310
32:	40 CONTINUE	SORT 320
33:	50 IF(GRDMIN .LE. ABS(DUMMY (J))) GO TO 30	SORT 330
34:	GRDMIN = ABS(DUMMY (J))	SORT 340
35:	NSORT(I) = J	SORT 350

36: 30 CONTINUE
37: 20 CONTINUE
38: RETURN
39: END

SORT 360
SORT 370
SORT 380
SORT 390

1:	SUBROUTINE OPTIM	OPTI	10
2:	C THIS SUBROUTINE CHECKS FOR AN IMPROVED FEASIBLE SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	OPTI	100
11:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	OPTI	110
12:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	OPTI	120
13:	1 RLBTY,WEIGHT	OPTI	130
14:	DIMENSION NUNIT1(50)	OPTI	140
15:	C	OPTI	150
16:	MOVE = 0	OPTI	160
17:	1 DO 10 I = 1,NMOD	OPTI	170
18:	NUNIT1(I) = NUNIT(I)	OPTI	180
19:	10 NUNIT(I) = NUNIT(I) + RGRADN(I)	OPTI	190
20:	DO 20 I=1,NMOD	OPTI	200
21:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) GO TO 100	OPTI	210
22:	20 CONTINUE	OPTI	220
23:	CALL CONSTR	OPTI	230
24:	IF(CONST(1).GT. 0.0 .OR. CONST(2).GT. 0.0) GO TO 100	OPTI	240
25:	CALL MODULE	OPTI	250
26:	CALL RELIAB	OPTI	260
27:	IF (RLBTY .LE. OBJFN) GO TO 100	OPTI	270
28:	MOVE = MOVE + 1	OPTI	280
29:	OBJFN = RLBTY	OPTI	290
30:	GO TO 999	OPTI	300
31:	100 DO 30 I=1,NMOD	OPTI	310
32:	30 NUNIT(I) = NUNIT1(I)	OPTI	320
33:	999 RETURN	OPTI	330
34:	END	OPTI	340

1:	SUBROUTINE OPTIMZ	OPTI	10
2:	C THIS SUBROUTINE FINDS A LOCALLY OPTIMAL SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/CONST(10),COSTF(50),G(50),RGRADN(50),RIGRDN(50),	OPTI	100
11:	1 RUNIT(50),WUNIT(50),CMCOST,OBJFN,RLBTY,WEIGHT	OPTI	110
12:	REAL*4 CONST,COSTF,G,RGRADN,RIGRDN,RUNIT,WUNIT,CMCOST,OBJFN,	OPTI	120
13:	1 RLBTY,WEIGHT	OPTI	130
14:	REAL*4 DUMMY(50)	OPTI	140
15:	C	OPTI	150
16:	CALL MODULE	OPTI	160
17:	CALL RELIAB	OPTI	170
18:	OBJFN = RLBTY	OPTI	180
19:	CALL SENSE	OPTI	190
20:	CALL SORT	OPTI	200
21:	CALL CONSTR	OPTI	210
22:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	OPTI	220
23:	WRITE(6,1500) OBJFN	OPTI	230
24:	WRITE(6,1600) CONST(1),CONST(2)	OPTI	240
25:	DO 5 I=1,NMOD	OPTI	250
26:	5 DUMMY(I) = RGRADN(I)	OPTI	260
27:	DO 200 I=1,NONZRO	OPTI	270
28:	DO 110 J=1,NMOD	OPTI	280
29:	110 RGRADN(J) = 0.0	OPTI	290
30:	I1 = NONZRO - I + 1	OPTI	300
31:	K = NSORT(I1)	OPTI	310
32:	RGRADN(K) = DUMMY(K)/ABS(DUMMY(K))	OPTI	320
33:	CALL OPTIM	OPTI	330
34:	IF(MOVE .EQ. 0) GO TO 200	OPTI	340
35:	GO TO 1000	OPTI	350

36:	200 CONTINUE	OPTI 360
37:	1000 RETURN	OPTI 370
38:	1400 FORMAT(1H0,/,10X,'NUNIT(I) =',20I5)	OPTI 380
39:	1500 FORMAT (1H0,/,30X,'SYSTEM RELIABILITY =',G13.6)	OPTI 390
40:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	OPTI 400
41:	1 G13.6)	OPTI 410
42:	END	OPTI 420

SYSTEM AVAILABILITY OPTIMIZATION

1:	C	MAIN PROGRAM	MAIN	10
2:	C	SYSTEM AVAILABILITY OPTIMIZATION	MAIN	20
3:	C	MODIFIED INTEGER GRADIENT METHOD	MAIN	30
4:		IMPLICIT INTEGER*4 (A-Z)	MAIN	40
5:		COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MAIN	50
6:		1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MAIN	60
7:		2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	MAIN	70
8:		INTEGER*2 P,PP,PATH	MAIN	80
9:		COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MAIN	90
10:		1 NMIN(50),FLAG,INDEX,MOVE,ONZRO	MAIN	100
11:		COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	MAIN	110
12:		1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	MAIN	120
13:		2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	MAIN	130
14:		3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MAIN	140
15:		REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	MAIN	150
16:		1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	MAIN	160
17:		2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MAIN	170
18:		REAL*4 YFL,FCOST,MCOST	MAIN	180
19:	C		MAIN	190
20:		PATHL = 300	MAIN	200
21:		PLENTH = 1000	MAIN	210
22:		PPL = 50	MAIN	220
23:	10	CALL DATAIN	MAIN	230
24:		READ(5,1000) INNOD,OUTNOD	MAIN	240
25:		CALL PATHS(INNOD)	MAIN	250
26:		IF(STATUS) 20,10,20	MAIN	260
27:	20	CALL PATHP(OUTNOD)	MAIN	270
28:		WRITE(6,1100)	MAIN	280
29:		DO 30 L=1,NPATH	MAIN	290
30:		WRITE(6,1200) L,(PATH(L,M),M=1,NMOD)	MAIN	300
31:	30	CONTINUE	MAIN	310
32:		CALL PATHC	MAIN	320
33:		WRITE(6,1001) NPATH,MPATH	MAIN	330
34:		WRITE(6,2600)	MAIN	340
35:		DO 35 L=1,MPATH	MAIN	350

36:	WRITE(6,2700) L,PSIGN(L),(PATH(L,M),M=1,NMOD)	MAIN 360
37:	35 CONTINUE	MAIN 370
38:	KOUNT = 1	MAIN 380
39:	ISEED = 984376521	MAIN 390
40:	IX = ISEED	MAIN 400
41:	DO 45 I=1,NMOD	MAIN 410
42:	45 NUNIT(I) = 1	MAIN 420
43:	TM = TMIN + (TMAX-TMIN)/2.	MAIN 430
44:	GO TO 55	MAIN 440
45:	C GENERATING A UNIFORMLY DISTRIBUTED RANDOM STARTING POINT	MAIN 450
46:	40 DO 50 I=1,NMOD	MAIN 460
47:	CALL RANDU(IX,IY,YFL)	MAIN 470
48:	NUNIT(I) = YFL*(NMAX(I)- NMIN(I)+1) + NMIN(I)	MAIN 480
49:	IX = IY	MAIN 490
50:	50 CONTINUE	MAIN 500
51:	CALL RANDU(IX,IY,YFL)	MAIN 510
52:	TM = YFL*(TMAX-TMIN) + TMIN	MAIN 520
53:	IX = IY	MAIN 530
54:	55 WRITE (6,1300) KOUNT	MAIN 540
55:	WRITE(6,2100)	MAIN 550
56:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 560
57:	WRITE(6,1450) TM	MAIN 570
58:	C LOCATING A FEASIBLE POINT	MAIN 580
59:	WRITE(6,2200)	MAIN 590
60:	CALL FESIBL	MAIN 600
61:	WRITE(6,1410) MOVE,INDEX	MAIN 610
62:	IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD1) GO TO 90	MAIN 620
63:	GO TO 60	MAIN 630
64:	60 WRITE(6,2400)	MAIN 640
65:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 650
66:	WRITE(6,1450) TM	MAIN 660
67:	WRITE(6,1600) CONST(1),CONST(2)	MAIN 670
68:	C SEARCHING FOR LOCAL OPTIMUM	MAIN 680
69:	WRITE(6,2300)	MAIN 690
70:	65 CALL OPTIMZ	MAIN 700
71:	IF(MOVE .EQ. 0) GO TO 70	MAIN 710

72:	GO TO 65	MAIN 720
73:	70 WRITE(6,2500)	MAIN 730
74:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 740
75:	WRITE(6,1450) TM	MAIN 750
76:	FLAG = 1	MAIN 760
77:	CALL AVLBTY	MAIN 770
78:	CALL CONSTR	MAIN 780
79:	FCOST = 0.0	MAIN 790
80:	DO 80 I=1,NMOD	MAIN 800
81:	80 FCOST = FCOST + NUNIT(I)*COSTF(I)	MAIN 810
82:	MCOST = CONST(1) + CMCOST - FCOST	MAIN 820
83:	WRITE(6,1500) ASYS,PSYS,RSYS	MAIN 830
84:	WRITE(6,1600) CONST(1),CONST(2)	MAIN 840
85:	WRITE(6,1700) FCOST,MCOST	MAIN 850
86:	WRITE(6,1900) MTBM,MDT	MAIN 860
87:	90 KOUNT = KOUNT + 1	MAIN 870
88:	IF (KOUNT .GT. 25) GO TO 10	MAIN 880
89:	GO TO 40	MAIN 890
90:	1000 FORMAT (2I3)	MAIN 900
91:	1001 FORMAT(1H0,/,30X,'NPATH = ',I3,/,30X,'MPATH = ',I5)	MAIN 910
92:	1100 FORMAT (1H0,/,T20,'PATHS IN BINARY CODE',/)	MAIN 920
93:	1200 FORMAT (1H0,T16,'PATH',I3,' = ',50I2)	MAIN 930
94:	1300 FORMAT(1H1,40X,'KOUNT = ',I3,//)	MAIN 940
95:	1400 FORMAT(1H0,/,10X,'NUNIT(I) = ',20I5)	MAIN 950
96:	1410 FORMAT (1H0,/,30X,'MOVE = ',I4,10X,'INDEX = ',I4)	MAIN 960
97:	1450 FORMAT (1H0,10X,'TM = ',G13.5)	MAIN 970
98:	1500 FORMAT (1H0,/,30X,'SYSTEM AVAILABILITY = ',G13.6,/,30X,	MAIN 980
99:	1 'SYSTEM PROFIT = ',G13.6,/,30X,'SYSTEM RELIABILITY = ',G13.6)	MAIN 990
100:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) = ',G13.6,/,30X,'CONSTRAINT(2) = ',	MAIN1000
101:	1 G13.6)	MAIN1010
102:	1700 FORMAT(1H0,/,30X,'FCOST = ',G13.6,/,30X,'MCOST = ',G13.6)	MAIN1020
103:	1900 FORMAT(1H0,20X,'MTBM = ',G13.6,20X,'MDT = ',G13.6)	MAIN1030
104:	2100 FORMAT(1H0,50X,'THE INITIAL POINT IS')	MAIN1040
105:	2200 FORMAT(1H0,50X,'SEARCHING FOR THE FEASIBLE POINT')	MAIN1050
106:	2300 FORMAT(1H0,50X,'SEARCHING FOR THE LOCAL OPTIMUM POINT')	MAIN1060
107:	2400 FORMAT(1H0,////,50X,'THE FEASIBLE POINT IS')	MAIN1070

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108: 2500 FORMAT(1H0,////,50X,'THE OPTIMAL POINT IS') MAIN1080
109: 2600 FORMAT(1H0,//,T20,'PATHS AND PATH COMBINATIONS IN BINARY CODE',/) MAIN1090
110: 2700 FORMAT(1H0,T16,'PATHC',I3,' = ',I2,4X,50I2) MAIN1100
111:      END MAIN1110
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1:	SUBROUTINE DATAIN	DATA	10
2:	C THIS SUBROUTINE READS AND ECHOCHECKS THE DATA	DATA	20
3:	IMPLICIT INTEGER*4 (A-Z)	DATA	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	DATA	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	DATA	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	DATA	60
7:	INTEGER*2 P,PP,PATH	DATA	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	DATA	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	DATA	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	DATA	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	DATA	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	DATA	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	DATA	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	DATA	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	DATA	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	DATA	160
17:	DIMENSION TITLE(30)	DATA	170
18:	C	DATA	180
19:	READ(5,999,END=100)(TITLE(I),I=1,30)	DATA	190
20:	READ(5,1000) NNODES,NMOD,CMCOST,WEIGHT,INCOM,COSTMF,TLF,TEM,	DATA	200
21:	1 TSC,TMIN,TMAX,TSTEP,INTVL	DATA	210
22:	DO 10 I=1,NMOD	DATA	220
23:	10 READ(5,1007) KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),WUNIT(I),	DATA	230
24:	1 LAMDA(I)	DATA	240
25:	NMOD1 = NMOD + 1	DATA	250
26:	FOR=0	DATA	260
27:	REV=0	DATA	270
28:	DO 40 I=1,1000	DATA	280
29:	READ(5,1001)P(I,1),P(I,2),P(I,4)	DATA	290
30:	IF(P(I,1)-P(I,2)) 20,50,30	DATA	300
31:	20 FOR=FOR + 1	DATA	310
32:	GO TO 40	DATA	320
33:	30 REV=REV + 1	DATA	330
34:	40 CONTINUE	DATA	340
35:	50 NEDGES=FOR + REV	DATA	350

36:	NEDGS1=NEDGES + 1	DATA 360
37:	FREV=FOR + 1	DATA 370
38:	WRITE(6,1002) (TITLE(I),I=1,30)	DATA 380
39:	WRITE(6,1003)	DATA 390
40:	WRITE(6,1004)	DATA 400
41:	WRITE(6,1003)	DATA 410
42:	DO 60 I=1,NEDGES	DATA 420
43:	WRITE(6,1005) I,P(I,1),P(I,2),P(I,4)	DATA 430
44:	WRITE(6,1003)	DATA 440
45:	60 CONTINUE	DATA 450
46:	WRITE(6,1008)	DATA 460
47:	WRITE(6,1011)	DATA 470
48:	WRITE(6,1009)	DATA 480
49:	WRITE(6,1011)	DATA 490
50:	DO 70 I=1,NMOD	DATA 500
51:	WRITE(6,1010) I,KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),	DATA 510
52:	1 WUNIT(I),LAMDA(I)	DATA 520
53:	WRITE(6,1011)	DATA 530
54:	70 CONTINUE	DATA 540
55:	WRITE(6,1006) CMCOST,WEIGHT,INCOM,COSTMF,TLF,TEM,TSC,TMIN,TMAX,	DATA 550
56:	1 TSTEP,INTVL	DATA 560
57:	999 FORMAT(15A4)	DATA 570
58:	1000 FORMAT (2I10,5F10.0/6F10.0)	DATA 580
59:	1001 FORMAT (3I10)	DATA 590
60:	1002 FORMAT(1H1,/,T10,30A4,/,T50,'INVENTORY OF BRANCHES',/)	DATA 600
61:	1003 FORMAT (1H+,T38,44(' _'))	DATA 610
62:	1004 FORMAT (1H ,T38,' ' BRANCH ' ' ORIGIN ' ' END ' ',	DATA 620
63:	1 ' MODULE ' ')	DATA 630
64:	1005 FORMAT (1H ,T38,4(' ',4X,12,4X),' ')	DATA 640
65:	1006 FORMAT (1H0,T10,'CMCOST=',G13.6,/,T10,'WEIGHT=',G13.6,/,T10,	DATA 650
66:	1 'INCOM =',G13.6,/,T10,'COSTMF=',G13.6,/,T10,'TLF =',G13.6,/,	DATA 660
67:	2 T10,'TEM =',G13.6,/,T10,'TSC =',G13.6,/,T10,'TMIN =',	DATA 670
68:	3 G13.6,/,T10,'TMAX =',G13.6,/,T10,'TSTEP =',G13.6,/,T10,	DATA 680
69:	4 'INTVL =',G13.6)	DATA 690
70:	1007 FORMAT(4I10,2F10.0,E20.0)	DATA 700
71:	1008 FORMAT(1H0,T50,'INVENTORY OF MODULES',/)	DATA 710

72:	1009	FORMAT(1H ,T20,' MODULE ', ' KODE ', ' MAX-UNIT ',	DATA 720
73:	1	' MIN-UNIT ', 'M OUT OF N ', 'UNIT COSTF ', 'UNIT WEGHT ',	DATA 730
74:	1	' FAILURE RATE ')	DATA 740
75:	1010	FORMAT(1H ,T20,5(' ',4X,I2,4X),2(' ',F10.3),' ',G14.5,' ')	DATA 750
76:	1011	FORMAT(1H+,T20,92('_'))	DATA 760
77:		RETURN	DATA 770
78:	100	CALL EXIT	DATA 780
79:		END	DATA 790

1:	SUBROUTINE MODULE	MODU 10
2:	C THIS SUBROUTINE COMPUTES THE MODULE RELIABILITY AND	MODU 20
3:	C SENSITIVITY OF MODULE RELIABILITY TO THE MODULE REDUNDANCY.	MODU 30
4:	IMPLICIT INTEGER*4 (A-Z)	MODU 40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MODU 50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MODU 60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	MODU 70
8:	INTEGER*2 P,PP,PATH	MODU 80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MODU 90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	MODU 100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	MODU 110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	MODU 120
13:	2 ASYS,COSTMF,INCUM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	MODU 130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MODU 140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	MODU 150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	MODU 160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MODU 170
18:	REAL*4 DUMMY1,DUMMY2,DUMMY3	MODU 180
19:	C	MODU 190
20:	DO 200 K=1,NMOD	MODU 200
21:	IF(KODE(K) .EQ. 0) GO TO 100	MODU 210
22:	DUMMY1 = EXP(-LAMDA(K)*T*MUNIT(K))	MODU 220
23:	DUMMY2 = 1.0	MODU 230
24:	DUMMY3 = LAMDA(K) * T * MUNIT(K)	MODU 240
25:	G(K) = 1.0	MODU 250
26:	NM = NUNIT(K) - MUNIT(K)	MODU 260
27:	IF(NM .EQ. 0) GO TO 20	MODU 270
28:	DO 10 I=1,NM	MODU 280
29:	DUMMY2 = DUMMY2 * DUMMY3/I	MODU 290
30:	G(K) = G(K) + DUMMY2	MODU 300
31:	10 CONTINUE	MODU 310
32:	20 G(K) = G(K) * DUMMY1	MODU 320
33:	RIGRDN(K) = DUMMY1 * DUMMY2 * DUMMY3/(NM+1)	MODU 330
34:	GO TO 200	MODU 340
35:	100 DUMMY1 = EXP(-LAMDA(K)*T)	MODU 350

36:	DUMMY2 = DUMMY1**NUNIT(K)	MODU 360
37:	G(K) = DUMMY2	MODU 370
38:	NM = NUNIT(K) - MUNIT(K)	MODU 380
39:	IF(NM .EQ. 0) GO TO 120	MODU 390
40:	DO 110 I=1,NM	MODU 400
41:	DUMMY2 = DUMMY2 * ((1.-DUMMY1)/DUMMY1) * (NUNIT(K)-I+1)/I	MODU 410
42:	G(K) = G(K) + DUMMY2	MODU 420
43:	110 CONTINUE	MODU 430
44:	120 RIGRDN(K) = DUMMY2 * (1.-DUMMY1) * MUNIT(K)/(NM+1)	MODU 440
45:	200 CONTINUE	MODU 450
46:	RETURN	MODU 460
47:	END	MODU 470

1:	SUBROUTINE RELIAB	RELI	10
2:	C THIS SUBROUTINE CALCULATES THE SYSTEM RELIABILITY FROM THE	RELI	20
3:	C MODULE RELIABILITIES.	RELI	30
4:	IMPLICIT INTEGER*4 (A-Z)	RELI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	RELI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	RELI	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	RELI	70
8:	INTEGER*2 P,PP,PATH	RELI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	RELI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	RELI	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	RELI	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	RELI	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	RELI	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	RELI	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	RELI	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	RELI	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	RELI	170
18:	REAL*4 GPATH	RELI	180
19:	C	RELI	190
20:	RLBTY = 0.0	RELI	200
21:	DO 10 I=1,MPATH	RELI	210
22:	GPATH = PSIGN(I)	RELI	220
23:	DO 20 J=1,NMOD	RELI	230
24:	IF (PATH(I,J) .EQ. 0) GO TO 20	RELI	240
25:	GPATH = GPATH * G(J)	RELI	250
26:	20 CONTINUE	RELI	260
27:	RLBTY = RLBTY + GPATH	RELI	270
28:	10 CONTINUE	RELI	280
29:	RETURN	RELI	290
30:	END	RELI	300

1:	SUBROUTINE SENSE	SENS	10
2:	C THIS SUBROUTINE DETERMINES THE SENSITIVITY OF THE SYSTEM	SENS	20
3:	C RELIABILITY TO MODULE REDUNDANCIES.	SENS	30
4:	IMPLICIT INTEGER*4 (A-Z)	SENS	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SENS	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SENS	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	SENS	70
8:	INTEGER*2 P,PP,PATH	SENS	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SENS	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SENS	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	SENS	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	SENS	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	SENS	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SENS	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	SENS	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	SENS	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SENS	170
18:	REAL*4 DUMMY(50)	SENS	180
19:	C	SENS	190
20:	DO 10 I=1,NMOD	SENS	200
21:	10 DUMMY(I) = G(I)	SENS	210
22:	DO 100 I=1,NMOD	SENS	220
23:	DO 90 J=1,NMOD	SENS	230
24:	90 G(J) = DUMMY(J)	SENS	240
25:	G(I) = 1.0	SENS	250
26:	CALL RELIAB	SENS	260
27:	RGRADN(I) = RLBTY	SENS	270
28:	G(I) = 0.0	SENS	280
29:	CALL RELIAB	SENS	290
30:	RGRADN(I) = (RGRADN(I)-RLBTY) * RIGRDN(I)	SENS	300
31:	100 CONTINUE	SENS	310
32:	DO 150 I=1,NMOD	SENS	320
33:	150 G(I) = DUMMY(I)	SENS	330
34:	RETURN	SENS	340
35:	END	SENS	350

1:	SUBROUTINE AVLBTY	AVLB	10
2:	C THIS SUBROUTINE DETERMINES SYSTEM AVAILABILITY, SYSTEM PROFIT,	AVLB	20
3:	C AND SENSITIVITY OF SYSTEM AVAILABILITY AND SYSTEM PROFIT TO EACH	AVLB	30
4:	C MODULE'S REDUNDANCY AND TO SYSTEM MAINTENANCE INTERVAL.	AVLB	40
5:	IMPLICIT INTEGER*4 (A-Z)	AVLB	50
6:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	AVLB	60
7:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	AVLB	70
8:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	AVLB	80
9:	INTEGER*2 P,PP,PATH	AVLB	90
10:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	AVLB	100
11:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	AVLB	110
12:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	AVLB	120
13:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	AVLB	130
14:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	AVLB	140
15:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	AVLB	150
16:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	AVLB	160
17:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	AVLB	170
18:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	AVLB	180
19:	REAL*4 DELTA,DUMMY1	AVLB	190
20:	C	AVLB	200
21:	EVEN = IFIX((TM/INTVL) + 1.0)	AVLB	210
22:	EVEN = (EVEN/2) * 2	AVLB	220
23:	IF (EVEN .EQ. 0) EVEN = 2	AVLB	230
24:	DELTA = TM/EVEN	AVLB	240
25:	MTBM = 1.0	AVLB	250
26:	DO 10 J=1,NMOD	AVLB	260
27:	10 MTGRDN(J) = 0.0	AVLB	270
28:	N = EVEN/2	AVLB	280
29:	N1 = N-1	AVLB	290
30:	DO 50 I=1,N	AVLB	300
31:	T = DELTA * (2*I-1)	AVLB	310
32:	CALL MODULE	AVLB	320
33:	CALL RELIAB	AVLB	330
34:	RSYS = RLBTY	AVLB	340
35:	MTBM = MTBM + 4.0 * RSYS	AVLB	350

36:	IF(FLAG .EQ. 1) GO TO 50	AVLB 360
37:	CALL SENSE	AVLB 370
38:	DO 30 J=1,NMOD	AVLB 380
39:	30 MTGRDN(J) = MTGRDN(J) + 4.0 * RGRADN(J)	AVLB 390
40:	50 CONTINUE	AVLB 400
41:	DO 100 I=1,N1	AVLB 410
42:	IF (N1 .EQ. 0) GO TO 100	AVLB 420
43:	T = DELTA * 2 * I	AVLB 430
44:	CALL MODULE	AVLB 440
45:	CALL RELIAB	AVLB 450
46:	RSYS = RLBTY	AVLB 460
47:	MTBM = MTBM + 2. * RSYS	AVLB 470
48:	IF(FLAG .EQ. 1) GO TO 100	AVLB 480
49:	CALL SENSE	AVLB 490
50:	DO 70 J=1,NMOD	AVLB 500
51:	70 MTGRDN(J) = MTGRDN(J) + 2.0 * RGRADN(J)	AVLB 510
52:	100 CONTINUE	AVLB 520
53:	T = TM	AVLB 530
54:	CALL MODULE	AVLB 540
55:	CALL RELIAB	AVLB 550
56:	RSYS = RLBTY	AVLB 560
57:	MTBM = (MTBM+RSYS) * DELTA/3.	AVLB 570
58:	MDT = TEM * (1.-RSYS) + TSC * RSYS	AVLB 580
59:	ASYS = MTBM/(MTBM+MDT)	AVLB 590
60:	DUMMY1 = 0.0	AVLB 600
61:	DO 115 I=1,NMOD	AVLB 610
62:	115 DUMMY1 = DUMMY1 + NUNIT(I) * COSTF(I)	AVLB 620
63:	PSYS = INCOM*TLF*ASYS - DUMMY1*(1.+(1.-ASYS)*TLF*COSTMF/100.)	AVLB 630
64:	IF(FLAG .EQ. 1) GO TO 1000	AVLB 640
65:	CALL SENSE	AVLB 650
66:	DO 120 I=1,NMOD	AVLB 660
67:	120 MTGRDN(I) = (MTGRDN(I)+RGRADN(I)) * DELTA/3.	AVLB 670
68:	RGRADT = 0.0	AVLB 680
69:	DO 130 I=1,NMOD	AVLB 690
70:	IF(KODE(I) .EQ. 0) GO TO 140	AVLB 700
71:	RGRADT = RGRADT - ((NUNIT(I)-MUNIT(I)+1)/TM) * RGRADN(I)	AVLB 710

72:	GO TO 130	AVLB 720
73:	140 RGRADT = RGRADT - (LAMDA(I)*(NUNIT(I)-MUNIT(I)+1)/	AVLB 730
74:	1 (1.-EXP(-LAMDA(I)*TM))) * RGRADN(I)	AVLB 740
75:	130 CONTINUE	AVLB 750
76:	DO 150 I=1,NMOD	AVLB 760
77:	AGRADN(I) = (MDT*MTGRDN(I) + MTBM*(TEM-TSC)*RGRADN(I))/	AVLB 770
78:	1 (MTBM+MDT)**2	AVLB 780
79:	PGRADN(I) = INCOM*TLF*AGRADN(I) + DUMMY1*TLF*AGRADN(I)*COSTMF/100.	AVLB 790
80:	1 - COSTF(I)*(1.+(1.-ASYS)*TLF*COSTMF/100.)	AVLB 800
81:	150 CONTINUE	AVLB 810
82:	AGRADT = (MDT*RSYS + MTBM*(TEM-TSC)*RGRADT)/(MTBM+MDT)**2	AVLB 820
83:	AGRADT = AGRADT * TSTEP	AVLB 830
84:	PGRADT = (INCOM + DUMMY1*COSTMF/100.) *TLF*AGRADT	AVLB 840
85:	1000 RETURN	AVLB 850
86:	END	AVLB 860

1:	SUBROUTINE CONSTR	CONS	10
2:	C THIS SUBROUTINE COMPUTES THE SYSTEM CONSTRAINTS.	CONS	20
3:	IMPLICIT INTEGER*4 (A-Z)	CONS	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	CONS	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	CONS	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	CONS	60
7:	INTEGER*2 P,PP,PATH	CONS	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	CONS	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,ONZRO	CONS	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	CONS	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	CONS	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	CONS	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	CONS	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	CONS	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	CONS	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	CONS	160
17:	C	CONS	170
18:	CONST(1) = 0.0	CONS	180
19:	CONST(2) = 0.0	CONS	190
20:	DO 10 I=1,NMOD	CONS	200
21:	CONST(1) = CONST(1) + NUNIT(I)*COSTF(I)	CONS	210
22:	CONST(2) = CONST(2) + NUNIT(I)*WUNIT(I)	CONS	220
23:	10 CONTINUE	CONS	230
24:	CONST(1) = CONST(1) * (1.+(1.-ASYS)*TLF*COSTMF/100.)	CONS	240
25:	CONST(1) = CONST(1) - CMCOST	CONS	250
26:	CONST(2) = CONST(2) - WEIGHT	CONS	260
27:	RETURN	CONS	270
28:	END	CONS	280

1:	SUBROUTINE FESIBL	FESI	10
2:	C THIS SUBROUTINE FINDS THE FEASIBLE SOLUTION BY THE METHOD	FESI	20
3:	C OF WEIGHTED PERPENDICULARS.	FESI	30
4:	IMPLICIT INTEGER*4 (A-Z)	FESI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	FESI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	FESI	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	FESI	70
8:	INTEGER*2 P,PP,PATH	FESI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	FESI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	FESI	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	FESI	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	FESI	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	FESI	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	FESI	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	FESI	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	FESI	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	FESI	170
18:	REAL*4 Z(50),ZMIN,S,DUMMY1,ZMAX	FESI	180
19:	C	FESI	190
20:	MOVE = 0	FESI	200
21:	INDEX = 0	FESI	210
22:	5 FLAG = 0	FESI	220
23:	CALL AVLBTY	FESI	230
24:	CALL CONSTR	FESI	240
25:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	FESI	250
26:	WRITE(6,1450) TM	FESI	260
27:	WRITE(6,1600) CONST(1),CONST(2)	FESI	270
28:	IF(CONST(1) .LE. 0.0 .AND. CONST(2) .LE. 0.0) GO TO 1000	FESI	280
29:	IF(CONST(1) .LE. 0.0) GO TO 30	FESI	290
30:	S = CONST(1)	FESI	300
31:	DUMMY1 = 0.0	FESI	310
32:	DO 10 I=1,NMOD	FESI	320
33:	10 DUMMY1 = DUMMY1 + NUNIT(I)*COSTF(I)	FESI	330
34:	DO 20 I=1,NMOD	FESI	340
35:	20 Z(I) = -CONST(1) * (COSTF(I)*(1.+(1.-ASYS)*TLF*COSTMF/100.) -	FESI	350

36:	1 DUMMY1*AGRADN(I)*TLF*COSTMF/100.)	FESI 360
37:	Z(NMOD1) = -CONST(1) * (-DUMMY1*AGRADT*TLF*COSTMF/100.)	FESI 370
38:	30 IF(CONST(2) .LE. 0.0) GO TO 50	FESI 380
39:	S = S + CONST(2)	FESI 390
40:	DO 40 I=1,NMOD	FESI 400
41:	40 Z(I) = Z(I) - CONST(2) * WUNIT(I)	FESI 410
42:	Z(NMOD1) = Z(NMOD1) - CONST(2)*0.0	FESI 420
43:	50 DO 60 I=1,NMOD1	FESI 430
44:	60 Z(I) = +Z(I)/S	FESI 440
45:	ZMAX = ABS(Z(1))	FESI 450
46:	DO 70 J=2,NMOD1	FESI 460
47:	IF (ZMAX .LT. ABS(Z(J))) ZMAX = ABS(Z(J))	FESI 470
48:	70 CONTINUE	FESI 480
49:	DO 80 I=1,NMOD	FESI 490
50:	IF(Z(I)) 90,80,100	FESI 500
51:	90 Z(I) = (Z(I)/ZMAX) * 2. - 0.5	FESI 510
52:	Z(I) = AINT(Z(I))	FESI 520
53:	GO TO 80	FESI 530
54:	100 Z(I) = (Z(I)/ZMAX) * 2. + 0.5	FESI 540
55:	Z(I) = AINT(Z(I))	FESI 550
56:	80 CONTINUE	FESI 560
57:	Z(NMOD1) = (Z(NMOD1)/ZMAX) * 2.	FESI 570
58:	INDEX = 0	FESI 580
59:	MOVE = MOVE + 1	FESI 590
60:	DO 120 I=1,NMOD	FESI 600
61:	NUNIT(I) = NUNIT(I) + Z(I)	FESI 610
62:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) INDEX=INDEX+1	FESI 620
63:	IF(NUNIT(I) .LT. NMIN(I)) NUNIT(I) = NMIN(I)	FESI 630
64:	IF(NUNIT(I) .GT. NMAX(I)) NUNIT(I) = NMAX(I)	FESI 640
65:	120 CONTINUE	FESI 650
66:	TM = TM + Z(NMOD1) * TSTEP	FESI 660
67:	IF(TM.LT.TMIN .OR. TM.GT.TMAX) INDEX = INDEX+1	FESI 670
68:	IF(TM .LT. TMIN) TM = TMIN	FESI 680
69:	IF(TM .GT. TMAX) TM = TMAX	FESI 690
70:	IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD1) GO TO 999	FESI 700
71:	GO TO 5	FESI 710

72:	999 WRITE(6,2500)	FESI 720
73:	1000 RETURN	FESI 730
74:	1400 FORMAT(1H0,/,10X,'NUNIT(1) =',20I5)	FESI 740
75:	1450 FORMAT (1H0,10X,'TM =',G13.5)	FESI 750
76:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	FESI 760
77:	1 G13.6)	FESI 770
78:	2500 FORMAT(1H0,////,30X,'FESIBL POINT CAN NOT BE LOCATED',////)	FESI 780
79:	END	FESI 790

1:	SUBROUTINE SORT	SORT 10
2:	C THIS SUBROUTINE ARRANGES THE GRADIENT VECTOR IN ASCENDING ORDER.	SORT 20
3:	IMPLICIT INTEGER*4 (A-Z)	SORT 30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SORT 40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SORT 50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	SORT 60
7:	INTEGER*2 P,PP,PATH	SORT 70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SORT 80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SORT 90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	SORT 100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	SORT 110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	SORT 120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SORT 130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	SORT 140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	SORT 150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SORT 160
17:	REAL*4 GRDMIN,DUMMY(50)	SORT 170
18:	C	SORT 180
19:	DO 5 I=1,NMOD	SORT 190
20:	DUMMY(I) = AGRADN(I)	SORT 200
21:	5 CONTINUE	SORT 210
22:	DUMMY(NMOD1) = AGRADT	SORT 220
23:	NONZRO = 0	SORT 230
24:	DO 10 I=1,NMOD	SORT 240
25:	IF(ABS(DUMMY(I)) .LT. 1.E-15) GO TO 10	SORT 250
26:	NONZRO = NONZRO + 1	SORT 260
27:	10 CONTINUE	SORT 270
28:	DO 20 I=1,NONZRO	SORT 280
29:	GRDMIN = 1.E30	SORT 290
30:	DO 30 J=1,NMOD	SORT 300
31:	IF(ABS(DUMMY(J)) .LT. 1.E-15) GO TO 30	SORT 310
32:	I1 = I - 1	SORT 320
33:	IF(I1 .EQ.0) GO TO 50	SORT 330
34:	DO 40 K=1,I1	SORT 340
35:	IF(J .EQ. NSORT(K)) GO TO 30	SORT 350

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36:      40 CONTINUE
37:      50 IF (GRDMIN .LE. ABS(DUMMY (J))) GO TO 30
38:          GRDMIN = ABS(DUMMY (J))
39:          NSORT(I) = J
40:      30 CONTINUE
41:      20 CONTINUE
42:          RETURN
43:          END
```

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SORT 360
SORT 370
SORT 380
SORT 390
SORT 400
SORT 410
SORT 420
SORT 430
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1:	SUBROUTINE OPTIM	OPTI	10
2:	C THIS SUBROUTINE CHECKS FOR AN IMPROVED FEASIBLE SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	OPTI	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	OPTI	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	OPTI	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	OPTI	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	OPTI	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	160
17:	REAL*4 TM1	OPTI	170
18:	DIMENSION NUNIT1(50)	OPTI	180
19:	C	OPTI	190
20:	MOVE = 0	OPTI	200
21:	1 DO 10 I = 1,NMOD	OPTI	210
22:	NUNIT1(I) = NUNIT(I)	OPTI	220
23:	10 NUNIT(I) = NUNIT(I) + AGRADN(I)	OPTI	230
24:	TM1 = TM	OPTI	240
25:	TM = TM + AGRADT * TSTEP	OPTI	250
26:	DO 20 I=1,NMOD	OPTI	260
27:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) GO TO 100	OPTI	270
28:	20 CONTINUE	OPTI	280
29:	IF(TM.LT.TMIN .OR. TM.GT.TMAX) GO TO 100	OPTI	290
30:	FLAG = 1	OPTI	300
31:	CALL AVLBTY	OPTI	310
32:	IF(ASYS .LE. OBJFN) GO TO 100	OPTI	320
33:	CALL CONSTR	OPTI	330
34:	IF(CONST(1).GT. 0.0 .OR. CONST(2).GT. 0.0) GO TO 100	OPTI	340
35:	MOVE = MOVE + 1	OPTI	350

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36:      OBJFN = ASYS
37:      GO TO 999
38:  100 DO 30 I=1,NMOD
39:      30 NUNIT(I) = NUNIT1(I)
40:      TM = TM1
41:  999 RETURN
42:      END
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OPTI 360
OPTI 370
OPTI 380.
OPTI 390
OPTI 400
OPTI 410
OPTI 420
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1:	SUBROUTINE OPTIMZ	OPTI	10
2:	C THIS SUBROUTINE FINDS A LOCALLY OPTIMAL SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	OPTI	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	OPTI	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	OPTI	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	OPTI	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	OPTI	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	160
17:	REAL*4 DUMMY(50)	OPTI	170
18:	C	OPTI	180
19:	FLAG = 0	OPTI	190
20:	CALL AVLBTY	OPTI	200
21:	CALL SORT	OPTI	210
22:	CALL CONSTR	OPTI	220
23:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	OPTI	230
24:	WRITE(6,1450) TM	OPTI	240
25:	WRITE(6,1500) ASYS,PSYS,RSYS	OPTI	250
26:	WRITE(6,1600) CONST(1),CONST(2)	OPTI	260
27:	OBJFN = ASYS	OPTI	270
28:	DO 5 I=1,NMOD	OPTI	280
29:	5 DUMMY(I) = AGRADN(I)	OPTI	290
30:	DUMMY(NMOD1) = AGRADT	OPTI	300
31:	DO 200 I=1,NONZRO	OPTI	310
32:	DO 110 J=1,NMOD	OPTI	320
33:	110 AGRADN(J) = 0.0	OPTI	330
34:	I1 = NONZRO - I + 1	OPTI	340
35:	K = NSORT(I1)	OPTI	350

36:	AGRADN(K) = DUMMY(K)/ABS(DUMMY(K))	OPTI 360
37:	AGRADT = DUMMY(NMOD1)/ABS(DUMMY(K))	OPTI 370
38:	CALL OPTIM	OPTI 380
39:	IF(MOVE .EQ. 0) GO TO 200	OPTI 390
40:	GO TO 1000	OPTI 400
41:	200 CONTINUE	OPTI 410
42:	DO 300 I=1,NMOD	OPTI 420
43:	300 AGRADN(I) = 0.0	OPTI 430
44:	AGRADT = (100./TSTEP) * DUMMY(NMOD1)/ABS(DUMMY(NMOD1))	OPTI 440
45:	CALL OPTIM	OPTI 450
46:	1000 RETURN	OPTI 460
47:	1400 FORMAT(1H0,/,10X,'NUNIT(I) =',20I5)	OPTI 470
48:	1450 FORMAT (1H0,10X,'TM =',G13.5)	OPTI 480
49:	1500 FORMAT (1H0,/,30X,'SYSTEM AVAILABILITY =',G13.6,/,30X,	OPTI 490
50:	1 'SYSTEM PROFIT =',G13.6,/,30X,'SYSTEM RELIABILITY =',G13.6)	OPTI 500
51:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	OPTI 510
52:	1 G13.6)	OPTI 520
53:	END	OPTI 530

SYSTEM PROFIT OPTIMIZATION

1:	C	MAIN PROGRAM	MAIN	10
2:	C	SYSTEM PROFIT OPTIMIZATION	MAIN	20
3:	C	MODIFIED INTEGER GRADIENT METHOD	MAIN	30
4:		IMPLICIT INTEGER*4 (A-Z)	MAIN	40
5:		COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MAIN	50
6:	1	NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MAIN	60
7:	2	PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	MAIN	70
8:		INTEGER*2 P,PP,PATH	MAIN	80
9:		COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MAIN	90
10:	1	NMIN(50),FLAG,INDEX,MOVE,ONZRO	MAIN	100
11:		COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	MAIN	110
12:	1	MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	MAIN	120
13:	2	ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	MAIN	130
14:	3	RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MAIN	140
15:		REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	MAIN	150
16:	1	WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	MAIN	160
17:	2	CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MAIN	170
18:		REAL*4 YFL,FCOST,MCOST	MAIN	180
19:	C		MAIN	190
20:		PATHL = 300	MAIN	200
21:		PLENTH = 1000	MAIN	210
22:		PPL = 50	MAIN	220
23:	10	CALL DATAIN	MAIN	230
24:		READ(5,1000) INNOD,OUTNOD	MAIN	240
25:		CALL PATHS(INNOD)	MAIN	250
26:		IF(STATUS) 20,10,20	MAIN	260
27:	20	CALL PATHP(OUTNOD)	MAIN	270
28:		WRITE(6,1100)	MAIN	280
29:		DO 30 L=1,NPATH	MAIN	290
30:		WRITE(6,1200) L,(PATH(L,M),M=1,NMOD)	MAIN	300
31:	30	CONTINUE	MAIN	310
32:		CALL PATHC	MAIN	320
33:		WRITE(6,1001) NPATH,MPATH	MAIN	330
34:		WRITE(6,2600)	MAIN	340
35:		DO 35 L=1,MPATH	MAIN	350

36:		WRITE(6,2700) L,PSIGN(L),(PATH(L,M),M=1,NMOD)	MAIN 360
37:	35	CONTINUE	MAIN 370
38:		KOUNT = 1	MAIN 380
39:		ISEED = 984376521	MAIN 390
40:		IX = ISEED	MAIN 400
41:		DO 45 I=1,NMOD	MAIN 410
42:	45	NUNIT(I) = 1	MAIN 420
43:		TM = TMIN + (TMAX-TMIN)/2.	MAIN 430
44:		GO TO 55	MAIN 440
45:	C	GENERATING A UNIFORMLY DISTRIBUTED RANDOM STARTING POINT	MAIN 450
46:	40	DO 50 I=1,NMOD	MAIN 460
47:		CALL RANDU(IX,IY,YFL)	MAIN 470
48:		NUNIT(I) = YFL*(NMAX(I)- NMIN(I)+1) + NMIN(I)	MAIN 480
49:		IX = IY	MAIN 490
50:	50	CONTINUE	MAIN 500
51:		CALL RANDU(IX,IY,YFL)	MAIN 510
52:		TM = YFL*(TMAX-TMIN) + TMIN	MAIN 520
53:		IX = IY	MAIN 530
54:	55	WRITE (6,1300) KOUNT	MAIN 540
55:		WRITE(6,2100)	MAIN 550
56:		WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 560
57:		WRITE(6,1450) TM	MAIN 570
58:	C	LOCATING A FEASIBLE POINT	MAIN 580
59:		WRITE(6,2200)	MAIN 590
60:		CALL FESIBL	MAIN 600
61:		WRITE(6,1410) MOVE,INDEX	MAIN 610
62:		IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD1) GO TO 90	MAIN 620
63:		GO TO 60	MAIN 630
64:	60	WRITE(6,2400)	MAIN 640
65:		WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 650
66:		WRITE(6,1450) TM	MAIN 660
67:		WRITE(6,1600) CONST(1),CONST(2)	MAIN 670
68:	C	SEARCHING FOR LOCAL OPTIMUM	MAIN 680
69:		WRITE(6,2300)	MAIN 690
70:	65	CALL OPTIMZ	MAIN 700
71:		IF(MOVE .EQ. 0) GO TO 70	MAIN 710

72:	GO TO 65	MAIN 720
73:	70 WRITE(6,2500)	MAIN 730
74:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	MAIN 740
75:	WRITE(6,1450) TM	MAIN 750
76:	FLAG = 1	MAIN 760
77:	CALL AVLBTY	MAIN 770
78:	CALL CONSTR	MAIN 780
79:	FCOST = CONST(1) + CMCOST	MAIN 790
80:	MCOST = FCOST * (1.-ASYS) * TLF * COSTMF/100.	MAIN 800
81:	WRITE(6,1500) ASYS,PSYS,RSYS	MAIN 810
82:	WRITE(6,1600) CONST(1),CONST(2)	MAIN 820
83:	WRITE(6,1700) FCOST,MCOST	MAIN 830
84:	WRITE(6,1900) MTBM,MDT	MAIN 840
85:	90 KOUNT = KOUNT + 1	MAIN 850
86:	IF (KOUNT .GT. 25) GO TO 10	MAIN 860
87:	GO TO 40	MAIN 870
88:	1000 FORMAT (2I3)	MAIN 880
89:	1001 FORMAT(1H0,/,30X,'NPATH = ',I3,/,30X,'MPATH = ',I5)	MAIN 890
90:	1100 FORMAT (1H0,/,T20,'PATHS IN BINARY CODE',/)	MAIN 900
91:	1200 FORMAT (1H0,T16,'PATH',I3,' = ',50I2)	MAIN 910
92:	1300 FORMAT(1H1,40X,'KOUNT = ',I3,/,)	MAIN 920
93:	1400 FORMAT(1H0,/,10X,'NUNIT(I) = ',20I5)	MAIN 930
94:	1410 FORMAT (1H0,/,30X,'MOVE = ',I4,10X,'INDEX = ',I4)	MAIN 940
95:	1450 FORMAT (1H0,10X,'TM = ',G13.5)	MAIN 950
96:	1500 FORMAT (1H0,/,30X,'SYSTEM AVAILABILITY = ',G13.6,/,30X,	MAIN 960
97:	1 'SYSTEM PROFIT = ',G13.6,/,30X,'SYSTEM RELIABILITY = ',G13.6)	MAIN 970
98:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) = ',G13.6,/,30X,'CONSTRAINT(2) = ',	MAIN 980
99:	1 G13.6)	MAIN 990
100:	1700 FORMAT(1H0,/,30X,'FCOST = ',G13.6,/,30X,'MCOST = ',G13.6)	MAIN1000
101:	1900 FORMAT(1H0,20X,'MTBM = ',G13.6,20X,'MDT = ',G13.6)	MAIN1010
102:	2100 FORMAT(1H0,50X,'THE INITIAL POINT IS')	MAIN1020
103:	2200 FORMAT(1H0,50X,'SEARCHING FOR THE FEASIBLE POINT')	MAIN1030
104:	2300 FORMAT(1H0,50X,'SEARCHING FOR THE LOCAL OPTIMUM POINT')	MAIN1040
105:	2400 FORMAT(1H0,////,50X,'THE FEASIBLE POINT IS')	MAIN1050
106:	2500 FORMAT(1H0,////,50X,'THE OPTIMAL POINT IS')	MAIN1060
107:	2600 FORMAT(1H0,/,T20,'PATHS AND PATH COMBINATIONS IN BINARY CODE',/)	MAIN1070

108: 2700 FORMAT(1H0,T16,'PATHC',I3,' = ',I2,4X,50I2)
109: END

MAIN1080
MAIN1090

1:	SUBROUTINE DATAIN	DATA	10
2:	C THIS SUBROUTINE READS AND ECHOCHECKS THE DATA	DATA	20
3:	IMPLICIT INTEGER*4 (A-Z)	DATA	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	DATA	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	DATA	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	DATA	60
7:	INTEGER*2 P,PP,PATH	DATA	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	DATA	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	DATA	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	DATA	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	DATA	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	DATA	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	DATA	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	DATA	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	DATA	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	DATA	160
17:	DIMENSION TITLE(30)	DATA	170
18:	C	DATA	180
19:	READ(5,999,END=100)(TITLE(I),I=1,30)	DATA	190
20:	READ(5,1000) NNODES,NMOD,CMCOST,WEIGHT,INCOM,COSTMF,TLF,TEM,	DATA	200
21:	1 TSC,TMIN,TMAX,TSTEP,INTVL	DATA	210
22:	DO 10 I=1,NMOD	DATA	220
23:	10 READ(5,1007) KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),WUNIT(I),	DATA	230
24:	1 LAMDA(I)	DATA	240
25:	NMOD1 = NMOD + 1	DATA	250
26:	FOR=0	DATA	260
27:	REV=0	DATA	270
28:	DO 40 I=1,1000	DATA	280
29:	READ(5,1001)P(I,1),P(I,2),P(I,4)	DATA	290
30:	IF(P(I,1)-P(I,2)) 20,50,30	DATA	300
31:	20 FOR=FOR + 1	DATA	310
32:	GO TO 40	DATA	320
33:	30 REV=REV + 1	DATA	330
34:	40 CONTINUE	DATA	340
35:	50 NEDGES=FOR + REV	DATA	350

36:	NEDGSI=NEDGES + 1	DATA 360
37:	FREV=FOR + 1	DATA 370
38:	WRITE(6,1002) (TITLE(I),I=1,30)	DATA 380
39:	WRITE(6,1003)	DATA 390
40:	WRITE(6,1004)	DATA 400
41:	WRITE(6,1003)	DATA 410
42:	DO 60 I=1,NEDGES	DATA 420
43:	WRITE(6,1005) I,P(I,1),P(I,2),P(I,4)	DATA 430
44:	WRITE(6,1003)	DATA 440
45:	60 CONTINUE	DATA 450
46:	WRITE(6,1008)	DATA 460
47:	WRITE(6,1011)	DATA 470
48:	WRITE(6,1009)	DATA 480
49:	WRITE(6,1011)	DATA 490
50:	DO 70 I=1,NMOD	DATA 500
51:	WRITE(6,1010) I,KODE(I),NMAX(I),NMIN(I),MUNIT(I),COSTF(I),	DATA 510
52:	1 WUNIT(I),LAMDA(I)	DATA 520
53:	WRITE(6,1011)	DATA 530
54:	70 CONTINUE	DATA 540
55:	WRITE(6,1006) CMCOST,WEIGHT,INCOM,COSTMF,TLF,TEM,TSC,TMIN,TMAX,	DATA 550
56:	1 TSTEP,INTVL	DATA 560
57:	999 FORMAT(15A4)	DATA 570
58:	1000 FORMAT (2I10,5F10.0/6F10.0)	DATA 580
59:	1001 FORMAT (3I10)	DATA 590
60:	1002 FORMAT(1H1,/,T10,30A4,/,T50,'INVENTORY OF BRANCHES',/)	DATA 600
61:	1003 FORMAT (1H+,T38,44(' _'))	DATA 610
62:	1004 FORMAT (1H ,T38,' BRANCH ',' ORIGIN ',' END ',	DATA 620
63:	1 ' MODULE ')	DATA 630
64:	1005 FORMAT (1H ,T38,4(' ',4X,12,4X),' ')	DATA 640
65:	1006 FORMAT (1H0,T10,'CMCOST=',G13.6,/,T10,'WEIGHT=',G13.6,/,T10,	DATA 650
66:	1 'INCOM =',G13.6,/,T10,'COSTMF=',G13.6,/,T10,'TLF =',G13.6,/,	DATA 660
67:	2 T10,'TEM =',G13.6,/,T10,'TSC =',G13.6,/,T10,'TMIN =',	DATA 670
68:	3 G13.6,/,T10,'TMAX =',G13.6,/,T10,'TSTEP =',G13.6,/,T10,	DATA 680
69:	4 'INTVL =',G13.6)	DATA 690
70:	1007 FORMAT(4I10,2F10.0,E20.0)	DATA 700
71:	1008 FORMAT(1H0,T50,'INVENTORY OF MODULES',/)	DATA 710

72:	1009	FORMAT(1H ,T20,' MODULE ',' KODE ',' MAX-UNIT ',	DATA 720
73:	1	' MIN-UNIT ','M OUT OF N ','UNIT COSTF ','UNIT WEGHT ',	DATA 730
74:	1	' FAILURE RATE ')	DATA 740
75:	1010	FORMAT(1H ,T20,5(' ','4X,I2,4X),2(' ','F10.3'),' ','G14.5',' ')	DATA 750
76:	1011	FORMAT(1H+,T20,92(' _'))	DATA 760
77:		RETURN	DATA 770
78:	100	CALL EXIT	DATA 780
79:		END	DATA 790

1:	SUBROUTINE MODULE	MODU 10
2:	C THIS SUBROUTINE COMPUTES THE MODULE RELIABILITY AND	MODU 20
3:	C SENSITIVITY OF MODULE RELIABILITY TO THE MODULE REDUNDANCY.	MODU 30
4:	IMPLICIT INTEGER*4 (A-Z)	MODU 40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MODU 50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MODU 60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	MODU 70
8:	INTEGER*2 P,PP,PATH	MODU 80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	MODU 90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	MODU 100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	MODU 110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	MODU 120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	MODU 130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MODU 140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	MODU 150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	MODU 160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	MODU 170
18:	REAL*4 DUMMY1,DUMMY2,DUMMY3	MODU 180
19:	C	MODU 190
20:	DO 200 K=1,NMOD	MODU 200
21:	IF(KODE(K) .EQ. 0) GO TO 100	MODU 210
22:	DUMMY1 = EXP(-LAMDA(K)*T*MUNIT(K))	MODU 220
23:	DUMMY2 = 1.0	MODU 230
24:	DUMMY3 = LAMDA(K) * T * MUNIT(K)	MODU 240
25:	G(K) = 1.0	MODU 250
26:	NM = NUNIT(K) - MUNIT(K)	MODU 260
27:	IF(NM .EQ. 0) GO TO 20	MODU 270
28:	DO 10 I=1,NM	MODU 280
29:	DUMMY2 = DUMMY2 * DUMMY3/I	MODU 290
30:	G(K) = G(K) + DUMMY2	MODU 300
31:	10 CONTINUE	MODU 310
32:	20 G(K) = G(K) * DUMMY1	MODU 320
33:	RIGRDN(K) = DUMMY1 * DUMMY2 * DUMMY3/(NM+1)	MODU 330
34:	GO TO 200	MODU 340
35:	100 DUMMY1 = EXP(-LAMDA(K)*T)	MODU 350

36:	DUMMY2 = DUMMY1**NUNIT(K)	MODU 360
37:	G(K) = DUMMY2	MODU 370
38:	NM = NUNIT(K) - MUNIT(K)	MODU 380
39:	IF(NM .EQ. 0) GO TO 120	MODU 390
40:	DO 110 I=1,NM	MODU 400
41:	DUMMY2 = DUMMY2 * ((1.-DUMMY1)/DUMMY1) * (NUNIT(K)-I+1)/I	MODU 410
42:	G(K) = G(K) + DUMMY2	MODU 420
43:	110 CONTINUE	MODU 430
44:	120 RIGRON(K) = DUMMY2 * (1.-DUMMY1) * MUNIT(K)/(NM+1)	MODU 440
45:	200 CONTINUE	MODU 450
46:	RETURN	MODU 460
47:	END	MODU 470

1:	SUBROUTINE RELIAB	RELI	10
2:	C THIS SUBROUTINE CALCULATES THE SYSTEM RELIABILITY FROM THE	RELI	20
3:	C MODULE RELIABILITIES.	RELI	30
4:	IMPLICIT INTEGER*4 (A-Z)	RELI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	RELI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	RELI	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	RELI	70
8:	INTEGER*2 P,PP,PATH	RELI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	RELI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	RELI	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	RELI	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	RELI	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	RELI	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	RELI	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	RELI	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	RELI	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	RELI	170
18:	REAL*4 GPATH	RELI	180
19:	C	RELI	190
20:	RLBTY = 0.0	RELI	200
21:	DO 10 I=1,MPATH	RELI	210
22:	GPATH = PSIGN(I)	RELI	220
23:	DO 20 J=1,NMOD	RELI	230
24:	IF (PATH(I,J) .EQ. 0) GO TO 20	RELI	240
25:	GPATH = GPATH * G(J)	RELI	250
26:	20 CONTINUE	RELI	260
27:	RLBTY = RLBTY + GPATH	RELI	270
28:	10 CONTINUE	RELI	280
29:	RETURN	RELI	290
30:	END	RELI	300

1:	SUBROUTINE SENSE	SENS	10
2:	C THIS SUBROUTINE DETERMINES THE SENSITIVITY OF THE SYSTEM	SENS	20
3:	C RELIABILITY TO MODULE REDUNDANCIES.	SENS	30
4:	IMPLICIT INTEGER*4 (A-Z)	SENS	40
5:	COMMON/MASON/FUR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SENS	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SENS	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	SENS	70
8:	INTEGER*2 P,PP,PATH	SENS	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SENS	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SENS	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	SENS	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	SENS	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	SENS	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SENS	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	SENS	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	SENS	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SENS	170
18:	REAL*4 DUMMY(50)	SENS	180
19:	C	SENS	190
20:	DO 10 I=1,NMOD	SENS	200
21:	10 DUMMY(I) = G(I)	SENS	210
22:	DO 100 I=1,NMOD	SENS	220
23:	DO 90 J=1,NMOD	SENS	230
24:	90 G(J) = DUMMY(J)	SENS	240
25:	G(I) = 1.0	SENS	250
26:	CALL RELIAB	SENS	260
27:	RGRADN(I) = RLBTY	SENS	270
28:	G(I) = 0.0	SENS	280
29:	CALL RELIAB	SENS	290
30:	RGRADN(I) = (RGRADN(I)-RLBTY) * RIGRDN(I)	SENS	300
31:	100 CONTINUE	SENS	310
32:	DO 150 I=1,NMOD	SENS	320
33:	150 G(I) = DUMMY(I)	SENS	330
34:	RETURN	SENS	340
35:	END	SENS	350

1:	SUBROUTINE AVLBTY	AVLB	10
2:	C THIS SUBROUTINE DETERMINES SYSTEM AVAILABILITY, SYSTEM PROFIT,	AVLB	20
3:	C AND SENSITIVITY OF SYSTEM AVAILABILITY AND SYSTEM PROFIT TO EACH	AVLB	30
4:	C MODULE'S REDUNDANCY AND TO SYSTEM MAINTENANCE INTERVAL.	AVLB	40
5:	IMPLICIT INTEGER*4 (A-Z)	AVLB	50
6:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	AVLB	60
7:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	AVLB	70
8:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	AVLB	80
9:	INTEGER*2 P,PP,PATH	AVLB	90
10:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	AVLB	100
11:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	AVLB	110
12:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	AVLB	120
13:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	AVLB	130
14:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	AVLB	140
15:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	AVLB	150
16:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	AVLB	160
17:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	AVLB	170
18:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	AVLB	180
19:	REAL*4 DELTA,DUMMY1	AVLB	190
20:	C	AVLB	200
21:	EVEN = IFIX((TM/INTVL) + 1.0)	AVLB	210
22:	EVEN = (EVEN/2) * 2	AVLB	220
23:	IF (EVEN .EQ. 0) EVEN = 2	AVLB	230
24:	DELTA = TM/EVEN	AVLB	240
25:	MTBM = 1.0	AVLB	250
26:	DO 10 J=1,NMOD	AVLB	260
27:	10 MTGRDN(J) = 0.0	AVLB	270
28:	N = EVEN/2	AVLB	280
29:	N1 = N-1	AVLB	290
30:	DO 50 I=1,N	AVLB	300
31:	T = DELTA * (2*I-1)	AVLB	310
32:	CALL MODULE	AVLB	320
33:	CALL RELIAB	AVLB	330
34:	RSYS = RLBTY	AVLB	340
35:	MTBM = MTBM + 4.0 * RSYS	AVLB	350

36:	IF(FLAG .EQ. 1) GO TO 50	AVLB 360
37:	CALL SENSE	AVLB 370
38:	DO 30 J=1,NMOD	AVLB 380
39:	30 MTGRDN(J) = MTGRDN(J) + 4.0 * RGRADN(J)	AVLB 390
40:	50 CONTINUE	AVLB 400
41:	DO 100 I=1,N1	AVLB 410
42:	IF (N1 .EQ. 0) GO TO 100	AVLB 420
43:	T = DELTA * 2 * I	AVLB 430
44:	CALL MODULE	AVLB 440
45:	CALL RELIAB	AVLB 450
46:	RSYS = RLBTY	AVLB 460
47:	MTBM = MTBM + 2. * RSYS	AVLB 470
48:	IF(FLAG .EQ. 1) GO TO 100	AVLB 480
49:	CALL SENSE	AVLB 490
50:	DO 70 J=1,NMOD	AVLB 500
51:	70 MTGRDN(J) = MTGRDN(J) + 2.0 * RGRADN(J)	AVLB 510
52:	100 CONTINUE	AVLB 520
53:	T = TM	AVLB 530
54:	CALL MODULE	AVLB 540
55:	CALL RELIAB	AVLB 550
56:	RSYS = RLBTY	AVLB 560
57:	MTBM = (MTBM+RSYS) * DELTA/3.	AVLB 570
58:	MDT = TEM * (1.-RSYS) + TSC * RSYS	AVLB 580
59:	ASYS = MTBM/(MTBM+MDT)	AVLB 590
60:	DUMMY1 = 0.0	AVLB 600
61:	DO 115 I=1,NMOD	AVLB 610
62:	115 DUMMY1 = DUMMY1 + NUNIT(I) * COSTF(I)	AVLB 620
63:	PSYS = INCOM*TLF*ASYS - DUMMY1*(1.+(1.-ASYS)*TLF*COSTMF/100.)	AVLB 630
64:	IF(FLAG .EQ. 1) GO TO 1000	AVLB 640
65:	CALL SENSE	AVLB 650
66:	DO 120 I=1,NMOD	AVLB 660
67:	120 MTGRDN(I) = (MTGRDN(I)+RGRADN(I)) * DELTA/3.	AVLB 670
68:	RGRADT = 0.0	AVLB 680
69:	DO 130 I=1,NMOD	AVLB 690
70:	IF(KODE(I) .EQ. 0) GO TO 140	AVLB 700
71:	RGRADT = RGRADT - ((NUNIT(I)-MUNIT(I)+1)/TM) * RGRADN(I)	AVLB 710

72:	GO TO 130	AVLB 720
73:	140 RGRADT = RGRADT - (LAMDA(I)*(NUNIT(I)-MUNIT(I)+1)/	AVLB 730
74:	1 (1.-EXP(-LAMDA(I)*TM))) * RGRADN(I)	AVLB 740
75:	130 CONTINUE	AVLB 750
76:	DO 150 I=1,NMOD	AVLB 760
77:	AGRADN(I) = (MDT*MTGRDN(I) + MTBM*(TEM-TSC)*RGRADN(I))/	AVLB 770
78:	1 (MTBM+MDT)**2	AVLB 780
79:	PGRADN(I) = INCOM*TLF*AGRADN(I) + DUMMY1*TLF*AGRADN(I)*COSTMF/100.	AVLB 790
80:	1 - COSTF(I)*(1.+(1.-ASYS)*TLF*COSTMF/100.)	AVLB 800
81:	150 CONTINUE	AVLB 810
82:	AGRADT = (MDT*RSYS + MTBM*(TEM-TSC)*RGRADT)/(MTBM+MDT)**2	AVLB 820
83:	AGRADT = AGRADT * TSTEP	AVLB 830
84:	PGRADT = (INCOM + DUMMY1*COSTMF/100.) *TLF*AGRADT	AVLB 840
85:	1000 RETURN	AVLB 850
86:	END	AVLB 860

1:	SUBROUTINE CONSTR	CONS	10
2:	C THIS SUBROUTINE COMPUTES THE SYSTEM CONSTRAINTS.	CONS	20
3:	IMPLICIT INTEGER*4 (A-Z)	CONS	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	CONS	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	CONS	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	CONS	60
7:	INTEGER*2 P,PP,PATH	CONS	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	CONS	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	CONS	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	CONS	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	CONS	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	CONS	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	CONS	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	CONS	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	CONS	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	CONS	160
17:	C	CONS	170
18:	CONST(1) = 0.0	CONS	180
19:	CONST(2) = 0.0	CONS	190
20:	DO 10 I=1,NMOD	CONS	200
21:	CONST(1) = CONST(1) + NUNIT(I)*COSTF(I)	CONS	210
22:	CONST(2) = CONST(2) + NUNIT(I)*WUNIT(I)	CONS	220
23:	10 CONTINUE	CONS	230
24:	CONST(1) = CONST(1) - CMCOST	CONS	240
25:	CONST(2) = CONST(2) - WEIGHT	CONS	250
26:	RETURN	CONS	260
27:	END	CONS	270

1:	SUBROUTINE FESIBL	FESI	10
2:	C THIS SUBROUTINE FINDS THE FEASIBLE SOLUTION BY THE METHOD	FESI	20
3:	C OF WEIGHTED PERPENDICULARS.	FESI	30
4:	IMPLICIT INTEGER*4 (A-Z)	FESI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	FESI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	FESI	60
7:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	FESI	70
8:	INTEGER*2 P,PP,PATH	FESI	80
9:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	FESI	90
10:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	FESI	100
11:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	FESI	110
12:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	FESI	120
13:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	FESI	130
14:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	FESI	140
15:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	FESI	150
16:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	FESI	160
17:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	FESI	170
18:	REAL*4 Z(50),ZMIN,S,DUMMY1,ZMAX	FESI	180
19:	C	FESI	190
20:	MOVE = 0	FESI	200
21:	INDEX = 0	FESI	210
22:	5 CALL CONSTR	FESI	220
23:	WRITE(6,1400) (NUNIT(I),I=1,NEDGES)	FESI	230
24:	WRITE(6,1450) TM	FESI	240
25:	WRITE(6,1600) CONST(1),CONST(2)	FESI	250
26:	IF(CONST(1) .LE. 0.0 .AND. CONST(2) .LE. 0.0) GO TO 1000	FESI	260
27:	IF(CONST(1) .LE. 0.0) GO TO 30	FESI	270
28:	S = CONST(1)	FESI	280
29:	DO 20 I=1,NMOD	FESI	290
30:	20 Z(I) = -CONST(1) * COSTF(I)	FESI	300
31:	Z(NMOD1) = -CONST(1) * 0.0	FESI	310
32:	30 IF(CONST(2) .LE. 0.0) GO TO 50	FESI	320
33:	S = S + CONST(2)	FESI	330
34:	DO 40 I=1,NMOD	FESI	340
35:	40 Z(I) = Z(I) - CONST(2) * WUNIT(I)	FESI	350

36:	Z(NMOD1) = Z(NMOD1) - CONST(2)*0.0	FESI 360
37:	50 DO 60 I=1,NMOD1	FESI 370
38:	60 Z(I) = +Z(I)/S	FESI 380
39:	ZMAX = ABS(Z(1))	FESI 390
40:	DO 70 J=2,NMOD1	FESI 400
41:	IF (ZMAX .LT. ABS(Z(J))) ZMAX = ABS(Z(J))	FESI 410
42:	70 CONTINUE	FESI 420
43:	DO 80 I=1,NMOD	FESI 430
44:	IF(Z(I)) 90,80,100	FESI 440
45:	90 Z(I) = (Z(I)/ZMAX) * 2. - 0.5	FESI 450
46:	Z(I) = AINT(Z(I))	FESI 460
47:	GO TO 80	FESI 470
48:	100 Z(I) = (Z(I)/ZMAX) * 2. + 0.5	FESI 480
49:	Z(I) = AINT(Z(I))	FESI 490
50:	80 CONTINUE	FESI 500
51:	Z(NMOD1) = (Z(NMOD1)/ZMAX) * 2.	FESI 510
52:	INDEX = 0	FESI 520
53:	MOVE = MOVE + 1	FESI 530
54:	DO 120 I=1,NMOD	FESI 540
55:	NUNIT(I) = NUNIT(I) + Z(I)	FESI 550
56:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) INDEX=INDEX+1	FESI 560
57:	IF(NUNIT(I) .LT. NMIN(I)) NUNIT(I) = NMIN(I)	FESI 570
58:	IF(NUNIT(I) .GT. NMAX(I)) NUNIT(I) = NMAX(I)	FESI 580
59:	120 CONTINUE	FESI 590
60:	TM = TM + Z(NMOD1) * TSTEP	FESI 600
61:	IF(TM.LT.TMIN .OR. TM.GT.TMAX) INDEX = INDEX+1	FESI 610
62:	IF(TM .LT. TMIN) TM = TMIN	FESI 620
63:	IF(TM .GT. TMAX) TM = TMAX	FESI 630
64:	IF (MOVE .GT. 20 .OR. INDEX .EQ. NMOD1) GO TO 999	FESI 640
65:	GO TO 5	FESI 650
66:	999 WRITE(6,2500)	FESI 660
67:	1000 RETURN	FESI 670
68:	1400 FORMAT(1H0,/,10X,'NUNIT(I) =',20I5)	FESI 680
69:	1450 FORMAT (1H0,10X,'TM =',G13.5)	FESI 690
70:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	FESI 700
71:	1 G13.6)	FESI 710

72: 2500 FORMAT(1H0,////,30X,'FESIBL POINT CAN NOT BE LOCATED',////)
73: END

FESI 720
FESI 730

1:	SUBROUTINE SORT	SORT 10
2:	C THIS SUBROUTINE ARRANGES THE GRADIENT VECTOR IN ASCENDING ORDER.	SORT 20
3:	IMPLICIT INTEGER*4 (A-Z)	SORT 30
4:	COMMON/MASON/FUR,FREV,LEND,NEDGES,NEDGS1,NNODES,	SORT 40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	SORT 50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	SORT 60
7:	INTEGER*2 P,PP,PATH	SORT 70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	SORT 80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	SORT 90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	SORT 100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	SORT 110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	SORT 120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SORT 130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	SORT 140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	SORT 150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	SORT 160
17:	REAL*4 GRDMIN,DUMMY(50)	SORT 170
18:	C	SORT 180
19:	DO 5 I=1,NMOD	SORT 190
20:	DUMMY(I) = PGRADN(I)	SORT 200
21:	5 CONTINUE	SORT 210
22:	DUMMY(NMOD1) = PGRADT	SORT 220
23:	NONZRO = 0	SORT 230
24:	DO 10 I=1,NMOD	SORT 240
25:	IF(ABS(DUMMY(I))) .LT. 1.E-15) GO TO 10	SORT 250
26:	NONZRO = NONZRO + 1	SORT 260
27:	10 CONTINUE	SORT 270
28:	DO 20 I=1,NONZRO	SORT 280
29:	GRDMIN = 1.E30	SORT 290
30:	DO 30 J=1,NEDGES	SORT 300
31:	IF(ABS(DUMMY(J))) .LT. 1.E-15) GO TO 30	SORT 310
32:	I1 = I - 1	SORT 320
33:	IF(I1 .EQ.0) GO TO 50	SORT 330
34:	DO 40 K=1,I1	SORT 340
35:	IF(J .EQ. NSORT(K)) GO TO 30	SORT 350

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36:      40 CONTINUE
37:      50 IF (GRDMIN .LE. ABS(DUMMY (J))) GO TO 30
38:          GRDMIN = ABS(DUMMY (J))
39:          NSORT(I) = J
40:      30 CONTINUE
41:      20 CONTINUE
42:          RETURN
43:          END
```

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SORT 360
SORT 370
SORT 380
SORT 390
SORT 400
SORT 410
SORT 420
SORT 430
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1:	SUBROUTINE OPTIM	OPTI	10
2:	C THIS SUBROUTINE CHECKS FOR AN IMPROVED FEASIBLE SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	OPTI	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	OPTI	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	OPTI	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	OPTI	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	OPTI	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	160
17:	REAL*4 TM1	OPTI	170
18:	DIMENSION NUNIT1(50)	OPTI	180
19:	C	OPTI	190
20:	MOVE = 0	OPTI	200
21:	1 DO 10 I = 1,NMOD	OPTI	210
22:	NUNIT1(I) = NUNIT(I)	OPTI	220
23:	10 NUNIT(I) = NUNIT(I) + PGRADN(I)	OPTI	230
24:	TM1 = TM	OPTI	240
25:	TM = TM + PGRADT * TSTEP	OPTI	250
26:	DO 20 I=1,NMOD	OPTI	260
27:	IF(NUNIT(I).LT. NMIN(I) .OR. NUNIT(I).GT.NMAX(I)) GO TO 100	OPTI	270
28:	20 CONTINUE	OPTI	280
29:	IF(TM.LT.TMIN .OR. TM.GT.TMAX) GO TO 100	OPTI	290
30:	FLAG = 1	OPTI	300
31:	CALL AVLBTY	OPTI	310
32:	IF(PSYS .LE. OBJFN) GO TO 100	OPTI	320
33:	CALL CONSTR	OPTI	330
34:	IF(CONST(1).GT. 0.0 .OR. CONST(2).GT. 0.0) GO TO 100	OPTI	340
35:	MOVE = MOVE + 1	OPTI	350

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36:      OBJFN = PSYS
37:      GO TO 999
38:  100 DO 30 I=1,NMOD
39:      30 NUNIT(I) = NUNIT1(I)
40:      TM = TM1
41:  999 RETURN
42:      END
```

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OPTI 360
OPTI 370
OPTI 380
OPTI 390
OPTI 400
OPTI 410
OPTI 420
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1:	SUBROUTINE OPTIMZ	OPTI	10
2:	C THIS SUBROUTINE FINDS A LOCALLY OPTIMAL SOLUTION.	OPTI	20
3:	IMPLICIT INTEGER*4 (A-Z)	OPTI	30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	OPTI	40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	OPTI	50
6:	2 PATH(300,50),NMOD,NMOD1,MPATH,PSIGN(300)	OPTI	60
7:	INTEGER*2 P,PP,PATH	OPTI	70
8:	COMMON/OPTM1/NUNIT(50),NMAX(50),MUNIT(50),KODE(50),NSORT(50),	OPTI	80
9:	1 NMIN(50),FLAG,INDEX,MOVE,NONZRO	OPTI	90
10:	COMMON/OPTM2/AGRADN(50),CONST(10),COSTF(50),G(50),LAMDA(50),	OPTI	100
11:	1 MTGRDN(50),PGRADN(50),RGRADN(50),RIGRDN(50),WUNIT(50),AGRADT,	OPTI	110
12:	2 ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,CMCOST,	OPTI	120
13:	3 RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	130
14:	REAL*4 AGRADN,CONST,COSTF,G,LAMDA,MTGRDN,PGRADN,RGRADN,RIGRDN,	OPTI	140
15:	1 WUNIT,AGRADT,ASYS,COSTMF,INCOM,INTVL,MTBM,MDT,OBJFN,PGRADT,PSYS,	OPTI	150
16:	2 CMCOST,RLBTY,RGRADT,RSYS,T,TEM,TLF,TM,TMIN,TMAX,TSC,TSTEP,WEIGHT	OPTI	160
17:	REAL*4 DUMMY(50)	OPTI	170
18:	C	OPTI	180
19:	FLAG = 0	OPTI	190
20:	CALL AVLBTY	OPTI	200
21:	CALL SORT	OPTI	210
22:	CALL CONSTR	OPTI	220
23:	WRITE(6,1400) (NUNIT(I),I=1,NMOD)	OPTI	230
24:	WRITE(6,1450) TM	OPTI	240
25:	WRITE(6,1500) ASYS,PSYS,RSYS	OPTI	250
26:	WRITE(6,1600) CONST(1),CONST(2)	OPTI	260
27:	OBJFN = PSYS	OPTI	270
28:	DO 5 I=1,NMOD	OPTI	280
29:	5 DUMMY(I) = PGRADN(I)	OPTI	290
30:	DUMMY(NMOD1) = PGRADT	OPTI	300
31:	DO 200 I=1,NONZRO	OPTI	310
32:	DO 110 J=1,NMOD	OPTI	320
33:	110 PGRADN(J) = 0.0	OPTI	330
34:	I1 = NONZRO - I + 1	OPTI	340
35:	K = NSORT(I1)	OPTI	350

36:	PGRADN(K) = DUMMY(K)/ABS(DUMMY(K))	OPTI 360
37:	PGRADT = DUMMY(NMOD1)/ABS(DUMMY(K))	OPTI 370
38:	CALL OPTIM	OPTI 380
39:	IF(MOVE .EQ. 0) GO TO 200	OPTI 390
40:	GO TO 1000	OPTI 400
41:	200 CONTINUE	OPTI 410
42:	DO 300 I=1,NMOD	OPTI 420
43:	300 PGRADN(I) = 0.0	OPTI 430
44:	PGRADT = (100./TSTEP) * DUMMY(NMOD1)/ABS(DUMMY(NMOD1))	OPTI 440
45:	CALL OPTIM	OPTI 450
46:	1000 RETURN	OPTI 460
47:	1400 FORMAT(1H0,/,10X,'NUNIT(I) =',20I5)	OPTI 470
48:	1450 FORMAT (1H0,10X,'TM =',G13.5)	OPTI 480
49:	1500 FORMAT (1H0,/,30X,'SYSTEM AVAILABILITY =',G13.6,/,30X,	OPTI 490
50:	1 'SYSTEM PROFIT =',G13.6,/,30X,'SYSTEM RELIABILITY =',G13.6)	OPTI 500
51:	1600 FORMAT(1H0,/,30X,'CONSTRAINT(1) =',G13.6,/,30X,'CONSTRAINT(2) =',	OPTI 510
52:	1 G13.6)	OPTI 520
53:	END	OPTI 530

COMPUTER LISTING FOR COMPUTATION SCHEME II

SYSTEM RELIABILITY OPTIMIZATION

1:	C	MAIN PROGRAM	MAIN	10
2:	C	SYSTEM RELIABILITY OPTIMIZATION	MAIN	20
3:	C	PARTIAL ENUMERATION METHOD	MAIN	30
4:		IMPLICIT INTEGER*4 (A-Z)	MAIN	40
5:		INTEGER*2 XTEST(40),XTEST1(40),XONE(40),XSTAR(40),XSTR1(40),	MAIN	50
6:		1 NSTAR1(40),MSTAR1(40),XHAT(40),NHAT(40),MHAT(40),NXTEST(40,5),	MAIN	60
7:		2 MXTEST(40,5),NXSTR1(40,5),MXSTR1(40,5)	MAIN	70
8:		DIMENSION COUNT(4)	MAIN	80
9:		INTEGER*2 P,PP,PATH	MAIN	90
10:		COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MAIN	100
11:		1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MAIN	110
12:		2 PATH(300,50),NMOD,MPATH,PSIGN(300)	MAIN	120
13:		COMMON/OPTM1/RKODE(40),NKODE(40),MKODE(40),NTEST(40),MTEST(40)	MAIN	130
14:		COMMON/OPTM2/R(40),G(40),GHAT(40),KINDEX(40),SINDEX(40),	MAIN	140
15:		1 CONSTR(40),COST,COSTHT,RLBTY,OBJFN,OBJHAT	MAIN	150
16:		REAL*4 R,G,GHAT,KINDEX,SINDEX,CONSTR,COST,COSTHT,RLBTY,OBJFN,	MAIN	160
17:		1 OBJHAT	MAIN	170
18:	C		MAIN	180
19:		PLENTH = 1000	MAIN	190
20:		PPL = 50	MAIN	200
21:		PATHL = 300	MAIN	210
22:	10	CALL DATAIN	MAIN	220
23:		READ(5,1000) INNOD,OUTNOD	MAIN	230
24:		CALL PATHS(INNOD)	MAIN	240
25:		IF(STATUS) 20,10,20	MAIN	250
26:	20	CALL PATHP(OUTNOD)	MAIN	260
27:		WRITE(6,1002)	MAIN	270
28:		DO 30 L= 1,NPATH	MAIN	280
29:		WRITE(6,1003) L,(PATH(L,M),M=1,NMOD)	MAIN	290
30:	30	CONTINUE	MAIN	300
31:		CALL PATHC	MAIN	310
32:		WRITE(6,1001) NPATH,MPATH	MAIN	320
33:		WRITE(6,1004)	MAIN	330
34:		DO 35 L=1,MPATH	MAIN	340
35:		WRITE(6,1005) L,PSIGN(L),(PATH(L,M),M=1,NMOD)	MAIN	350

36:	35	CONTINUE	MAIN 360
37:		OBJHAT = 1.0	MAIN 370
38:		DO 1 I = 1,4	MAIN 380
39:	1	COUNT(I) = 0	MAIN 390
40:	C	GENERATES TOTAL NUMBER OF VARIABLES IN THE PROBLEM	MAIN 400
41:		NVAR = 0	MAIN 410
42:		DO 40 I = 1,NMOD	MAIN 420
43:		NVAR = NVAR + NKODE(I) + MKODE(I)	MAIN 430
44:	40	CONTINUE	MAIN 440
45:	C	GENERATES INITIAL VALUES OF XTEST(I)	MAIN 450
46:		DO 50 I = 1,NVAR	MAIN 460
47:		XTEST(I) = 0	MAIN 470
48:		XONE(I) = 0	MAIN 480
49:	50	CONTINUE	MAIN 490
50:		XONE(I) = 1	MAIN 500
51:		XTEST(I) = 1	MAIN 510
52:	C	GENERATE XSTAR	MAIN 520
53:	100	CARRY = 0	MAIN 530
54:		DO 140 I=1,NVAR	MAIN 540
55:		XTEST1(I) = XTEST(I) + 1 + CARRY	MAIN 550
56:		IF (XTEST1(I).LE.1) GO TO 150	MAIN 560
57:		CARRY = 1	MAIN 570
58:		XTEST1(I) = XTEST1(I) - 2	MAIN 580
59:		GO TO 140	MAIN 590
60:	150	CARRY = 0	MAIN 600
61:	140	CONTINUE	MAIN 610
62:		TEST1 = 0	MAIN 620
63:		DO 200 I=1,NVAR	MAIN 630
64:		IF((XTEST(I).EQ.1) .OR. (XTEST1(I).EQ.1)) GO TO 210	MAIN 640
65:		XSTR1(I) = 0	MAIN 650
66:		GO TO 200	MAIN 660
67:	210	XSTR1(I) = 1	MAIN 670
68:		TEST1 = TEST1 + XSTR1(I)	MAIN 680
69:	200	CONTINUE	MAIN 690
70:		CARRY = 0	MAIN 700
71:		DO 220 I=1,NVAR	MAIN 710

72:	XSTAR(I) = XSTR1(I) + XONE(I) + CARRY	MAIN 720
73:	IF(XSTAR(I).LE.1) GO TO 230	MAIN 730
74:	CARRY = 1	MAIN 740
75:	XSTAR(I) = XSTAR(I) - 2	MAIN 750
76:	GO TO 220	MAIN 760
77:	230 CARRY = 0	MAIN 770
78:	220 CONTINUE	MAIN 780
79:	C GENERATE VALUES FOR NTEST(I),MTEST(I),NSTAR1(I),AND MSTAR1(I)	MAIN 790
80:	K = 1	MAIN 800
81:	DO 60 J = 1,5	MAIN 810
82:	DO 70 I=1,NMOD	MAIN 820
83:	IF (J.GT.NKODE(I)) GO TO 70	MAIN 830
84:	NXTEST(I,J) = XTEST(K)	MAIN 840
85:	NXSTR1(I,J) = XSTR1(K)	MAIN 850
86:	K = K + 1	MAIN 860
87:	IF (J.GT.MKODE(I)) GO TO 70	MAIN 870
88:	MXTEST(I,J) = XTEST(K)	MAIN 880
89:	MXSTR1(I,J) = XSTR1(K)	MAIN 890
90:	K = K + 1	MAIN 900
91:	70 CONTINUE	MAIN 910
92:	60 CONTINUE	MAIN 920
93:	DO 80 I=1,NMOD	MAIN 930
94:	NTEST(I) = 1	MAIN 940
95:	MTEST(I) = 1	MAIN 950
96:	NSTAR1(I) = 1	MAIN 960
97:	MSTAR1(I) = 1	MAIN 970
98:	L = NKODE(I)	MAIN 980
99:	IF (L .EQ. 0) GO TO 95	MAIN 990
100:	DO 90 J=1,L	MAIN1000
101:	NTEST(I) = NTEST(I) + (2**(J-1))*(1-NXTEST(I,J))	MAIN1010
102:	NSTAR1(I) = NSTAR1(I) + (2**(J-1))*(1-NXSTR1(I,J))	MAIN1020
103:	90 CONTINUE	MAIN1030
104:	95 L = MKODE(I)	MAIN1040
105:	IF (L .EQ. 0) GO TO 80	MAIN1050
106:	DO 110 J=1,L	MAIN1060
107:	MTEST(I) = MTEST(I) + (2**(J-1))*MXTEST(I,J)	MAIN1070

108:	MSTAR1(I) = MSTAR1(I) + (2** (J-1)) * MXSTR1(I,J)	MAIN1080
109:	110 CONTINUE	MAIN1090
110:	80 CONTINUE	MAIN1100
111:	C TEST FOR GI1(X*-1) - GI2(X)	MAIN1110
112:	COUNT(1) = COUNT(1) + 1	MAIN1120
113:	DO 310 I=1,NMOD	MAIN1130
114:	CONSTR(I) = NTEST(I) - MTEST(I)	MAIN1140
115:	IF(CONSTR(I).GE.0.0) GO TO 310	MAIN1150
116:	GO TO 400	MAIN1160
117:	310 CONTINUE	MAIN1170
118:	M = NMOD + 1	MAIN1180
119:	CONSTR(M) = 0.0	MAIN1190
120:	DO 320 I=1,NMOD	MAIN1200
121:	320 CONSTR(M) = CONSTR(M) + NSTAR1(I)*KINDEX(I)*(1./MSTAR1(I))**	MAIN1210
122:	1SINDEX(I)	MAIN1220
123:	CONSTR(M) = COST - CONSTR(M)	MAIN1230
124:	IF(CONSTR(M).GE.0.0) GO TO 330	MAIN1240
125:	GO TO 400	MAIN1250
126:	C TEST FOR GO(X) - GO(XHAT).	MAIN1260
127:	330 COUNT(2) = COUNT(2) + 1	MAIN1270
128:	CALL MODULE	MAIN1280
129:	CALL RELIAB	MAIN1290
130:	OBJFN = 1. - RLBTY	MAIN1300
131:	IF(OBJFN.LT.OBJHAT) GO TO 250	MAIN1310
132:	400 IF(TEST1.EQ.NVAR) GO TO 500	MAIN1320
133:	DO 260 I=1,NVAR	MAIN1330
134:	XTEST(I) = XSTAR(I)	MAIN1340
135:	260 CONTINUE	MAIN1350
136:	GO TO 100	MAIN1360
137:	C TEST FOR GI1(X)-GI2(X)	MAIN1370
138:	250 COUNT(3) = COUNT(3) + 1	MAIN1380
139:	M = NMOD + 1	MAIN1390
140:	CONSTR(M) = 0.0	MAIN1400
141:	DO 340 I=1,NMOD	MAIN1410
142:	340 CONSTR(M) = CONSTR(M) + NTEST (I)*KINDEX(I)*(1./MTEST (I))**	MAIN1420
143:	1SINDEX(I)	MAIN1430

144:	CONSTR(M) = COST - CONSTR(M)	MAIN1440
145:	IF(CONSTR(M).GE.0.0) GO TO 350	MAIN1450
146:	CARRY = 0	MAIN1460
147:	TEST2 = 0	MAIN1470
148:	DO 360 I=1,NVAR	MAIN1480
149:	TEST2 = TEST2 + XTEST(I)	MAIN1490
150:	XTEST(I) = XTEST(I) + XONE(I) + CARRY	MAIN1500
151:	IF(XTEST(I).LE.1) GO TO 370	MAIN1510
152:	CARRY = 1	MAIN1520
153:	XTEST(I) = XTEST(I) - 2	MAIN1530
154:	GO TO 360	MAIN1540
155:	370 CARRY = 0	MAIN1550
156:	360 CONTINUE	MAIN1560
157:	IF(TEST2.EQ.NVAR) GO TO 500	MAIN1570
158:	GO TO 100	MAIN1580
159:	350 COUNT(4) = COUNT(4) + 1	MAIN1590
160:	DO 380 I=1,NVAR	MAIN1600
161:	380 XHAT(I) = XTEST(I)	MAIN1610
162:	OBJHAT = OBJFN	MAIN1620
163:	DO 390 I=1,NMOD	MAIN1630
164:	NHAT(I) = NTEST(I)	MAIN1640
165:	MHAT(I) = MTEST(I)	MAIN1650
166:	GHAT(I) = G(I)	MAIN1660
167:	390 CONTINUE	MAIN1670
168:	M = NMOD + 1	MAIN1680
169:	COSTHT = COST - CONSTR(M)	MAIN1690
170:	GO TO 400	MAIN1700
171:	500 RLBTY = 1. - OBJHAT	MAIN1710
172:	WRITE(6,1006)	MAIN1720
173:	WRITE(6,1007)	MAIN1730
174:	WRITE(6,1008)	MAIN1740
175:	WRITE(6,1007)	MAIN1750
176:	DO 410 I=1,NMOD	MAIN1760
177:	WRITE(6,1009) I,NHAT(I),MHAT(I),GHAT(I)	MAIN1770
178:	WRITE(6,1007)	MAIN1780
179:	410 CONTINUE	MAIN1790

180:	WRITE(6,1010) RLBTY,COSTHT,(J,COUNT(J),J=1,4)	MAIN1800
181:	GO TO 10	MAIN1810
182:	1000 FORMAT(2I3)	MAIN1820
183:	1001 FORMAT(1H0,/,30X,'NPATH = ',I3,/,30X,'MPATH = ',I5)	MAIN1830
184:	1002 FORMAT(1H0,/,T20,'PATHS IN BINARY CODE',/)	MAIN1840
185:	1003 FORMAT(1H0,T16,'PATH',I3,' = ',50I2)	MAIN1850
186:	1004 FORMAT(1H0,/,T20,'PATHS AND PATH COMBINATIONS IN BINARY CODE',/)	MAIN1860
187:	1005 FORMAT(1H0,T16,'PATHC',I3,' = ',I2,4X,50I2)	MAIN1870
188:	1006 FORMAT(1H1,/,T40,'THE OPTIMAL POINT IS',/)	MAIN1880
189:	1007 FORMAT(1H+,T30,55('_'))	MAIN1890
190:	1008 FORMAT(1H ,T30,' MODULE ', ' N-OPTM ', ' M-OPTM ',	MAIN1900
191:	1 ' MODULE RELIABILITY ')	MAIN1910
192:	1009 FORMAT(1H ,T30,3(' ',4X,I2,4X), ' ',D20.10,' ')	MAIN1920
193:	1010 FORMAT(1H0,/,30X,'OPTIMIZED SYSTEM RELIABILITY = ',D20.10,/,30X,	MAIN1930
194:	1 'SYSTEM COST = ',G13.5,/, (30X,'COUNT',I3,' = ',I8,/))	MAIN1940
195:	END	MAIN1950

1:	SUBROUTINE DATAIN	DATA 10
2:	C THIS SUBROUTINE READS AND ECHOCHECKS THE DATA	DATA 20
3:	IMPLICIT INTEGER*4 (A-Z)	DATA 30
4:	INTEGER*2 P,PP,PATH	DATA 40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	DATA 50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	DATA 60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	DATA 70
8:	COMMON/OPTM1/RKODE(40),NKODE(40),MKODE(40),NTEST(40),MTEST(40)	DATA 80
9:	COMMON/OPTM2/R(40),G(40),GHAT(40),KINDEX(40),SINDEX(40),	DATA 90
10:	1 CONSTR(40),COST,COSTHT,RLBTY,OBJFN,OBJHAT	DATA 100
11:	REAL*4 R,G,GHAT,KINDEX,SINDEX,CONSTR,COST,COSTHT,RLBTY,OBJFN,	DATA 110
12:	1 OBJHAT	DATA 120
13:	DIMENSION TITLE(30)	DATA 130
14:	C	DATA 140
15:	READ(5,999,END=70) (TITLE(I),I=1,30)	DATA 150
16:	READ(5,1000)NNODES,NMOD,COST	DATA 160
17:	DO 10 I=1,NMOD	DATA 170
18:	10 READ(5,1006) RKODE(I),NKODE(I),MKODE(I),SINDEX(I),KINDEX(I),R(I)	DATA 180
19:	FOR=0	DATA 190
20:	REV=0	DATA 200
21:	DO 40 I=1,1000	DATA 210
22:	READ(5,1001) P(I,1),P(I,2),P(I,4)	DATA 220
23:	IF(P(I,1)-P(I,2)) 20,50,30	DATA 230
24:	20 FOR=FOR + 1	DATA 240
25:	GO TO 40	DATA 250
26:	30 REV=REV + 1	DATA 260
27:	40 CONTINUE	DATA 270
28:	50 NEDGES=FOR + REV	DATA 280
29:	NEDGS1=NEDGES + 1	DATA 290
30:	FREV=FOR + 1	DATA 300
31:	WRITE(6,1002) (TITLE(I),I=1,30)	DATA 310
32:	WRITE(6,1003)	DATA 320
33:	WRITE(6,1004)	DATA 330
34:	WRITE(6,1003)	DATA 340
35:	DO 60 I=1,NEDGES	DATA 350

36:	WRITE(6,1005) I,P(I,1),P(I,2),P(I,4)	DATA 360
37:	WRITE(6,1003)	DATA 370
38:	60 CONTINUE	DATA 380
39:	WRITE(6,1007) COST	DATA 390
40:	WRITE(6,1008)	DATA 400
41:	WRITE(6,1009)	DATA 410
42:	WRITE(6,1010)	DATA 420
43:	WRITE(6,1009)	DATA 430
44:	DO 80 I=1,NMOD	DATA 440
45:	WRITE(6,1011) I,RKODE(I),NKODE(I),MKODE(I),SINDEX(I),KINDEX(I),R(I)	DATA 450
46:	WRITE(6,1009)	DATA 460
47:	80 CONTINUE	DATA 470
48:	RETURN	DATA 480
49:	70 CALL EXIT	DATA 490
50:	999 FORMAT(15A4)	DATA 500
51:	1000 FORMAT(2I10,F10.0)	DATA 510
52:	1001 FORMAT(3I10)	DATA 520
53:	1002 FORMAT(1H1,/,T10,30A4,/,T50,'INVENTORY OF BRANCHES',/)	DATA 530
54:	1003 FORMAT(1H+,T38,44(' _'))	DATA 540
55:	1004 FORMAT (1H ,T38,' ' BRANCH ' ', ' ORIGIN ' ', ' END ' ',	DATA 550
56:	1 ' MODULE ' ')	DATA 560
57:	1005 FORMAT (1H ,T38,4(' ' ',4X,I2,4X),' ')	DATA 570
58:	1006 FORMAT(3I10,3F10.0)	DATA 580
59:	1007 FORMAT(1H0,/,T40,'COST CONSTRAINT =',F15.3,/,/)	DATA 590
60:	1008 FORMAT(1H0,T50,'INVENTORY OF MODULES',/)	DATA 600
61:	1009 FORMAT(1H+,T20,83(' _'))	DATA 610
62:	1010 FORMAT(1H ,T20,' ' MODULE ' ', ' R-KODE ' ', ' N-KODE ' ',	DATA 620
63:	1 ' M-KODE ' ', ' COST-COEFF ' ', ' COST/UNIT ' ', ' UNIT-REL ' ')	DATA 630
64:	1011 FORMAT(1H ,T20,4(' ' ',4X,I2,4X),2(' ' ',F10.3),3X,' ' ',G13.6,' ' ')	DATA 640
65:	END	DATA 650

1:	SUBROUTINE MODULE	MODU 10
2:	C THIS SUBROUTINE COMPUTES THE MODULE RELIABILITY.	MODU 20
3:	IMPLICIT INTEGER*4 (A-Z)	MODU 30
4:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	MODU 40
5:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	MODU 50
6:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	MODU 60
7:	INTEGER*2 P,PP,PATH	MODU 70
8:	COMMON/OPTM1/RKODE(40),NKODE(40),MKODE(40),NTEST(40),MTEST(40)	MODU 80
9:	COMMON/OPTM2/R(40),G(40),GHAT(40),KINDEX(40),SINDEX(40),	MODU 90
10:	1 CONSTR(40),COST,COSTHT,RLBTY,OBJFN,OBJHAT	MODU 100
11:	REAL*4 R,G,GHAT,KINDEX,SINDEX,CONSTR,COST,COSTHT,RLBTY,OBJFN,	MODU 110
12:	1 OBJHAT	MODU 120
13:	REAL*4 DUMMY1,DUMMY2,DUMMY3	MODU 130
14:	C	MODU 140
15:	DO 200 K=1,NMOD	MODU 150
16:	IF(RKODE(K) .EQ. 0) GO TO 100	MODU 160
17:	DUMMY1 = R(K) ** MTEST(K)	MODU 170
18:	DUMMY3 = - MTEST(K) * ALOG(R(K))	MODU 180
19:	G(K) = 1.0	MODU 190
20:	DUMMY2 = 1.0	MODU 200
21:	NM = NTEST(K) - MTEST(K)	MODU 210
22:	IF(NM .EQ. 0) GO TO 20	MODU 220
23:	DO 10 I=1,NM	MODU 230
24:	DUMMY2 = DUMMY2 * DUMMY3/I	MODU 240
25:	G(K) = G(K) + DUMMY2	MODU 250
26:	10 CONTINUE	MODU 260
27:	20 G(K) = G(K) * DUMMY1	MODU 270
28:	GO TO 200	MODU 280
29:	100 DUMMY1 = R(K)	MODU 290
30:	DUMMY2 = DUMMY1**NTEST(K)	MODU 300
31:	G(K) = DUMMY2	MODU 310
32:	NM = NTEST(K) - MTEST(K)	MODU 320
33:	IF(NM .EQ. 0) GO TO 200	MODU 330
34:	DO 110 I=1,NM	MODU 340
35:	DUMMY2 = DUMMY2 * ((1.-DUMMY1)/DUMMY1) * (NTEST(K)-I+1)/I	MODU 350

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36:      G(K) = G(K) + DUMMY2
37:    110 CONTINUE
38:    200 CONTINUE
39:      RETURN
40:      END
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MODU 360
MODU 370
MODU 380
MODU 390
MODU 400
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1:	SUBROUTINE RELIAB	RELI	10
2:	C THIS SUBROUTINE CALCULATES THE SYSTEM RELIABILITY FROM THE	RELI	20
3:	C MODULE RELIABILITIES.	RELI	30
4:	IMPLICIT INTEGER*4 (A-Z)	RELI	40
5:	COMMON/MASON/FOR,FREV,LEND,NEDGES,NEDGS1,NNODES,	RELI	50
6:	1 NPATH,PATHL,PLENTH,PPL,REV,STATUS,P(1000,4),PP(50),	RELI	60
7:	2 PATH(300,50),NMOD,MPATH,PSIGN(300)	RELI	70
8:	INTEGER*2 P,PP,PATH	RELI	80
9:	COMMON/OPTM1/RKODE(40),NKODE(40),MKODE(40),NTEST(40),MTEST(40)	RELI	90
10:	COMMON/OPTM2/R(40),G(40),GHAT(40),KINDEX(40),SINDEX(40),	RELI	100
11:	1 CONSTR(40),COST,COSTHT,RLBTY,OBJFN,OBJHAT	RELI	110
12:	REAL*4 R,G,GHAT,KINDEX,SINDEX,CONSTR,COST,COSTHT,RLBTY,OBJFN,	RELI	120
13:	1 OBJHAT	RELI	130
14:	REAL*4 GPATH	RELI	140
15:	C	RELI	150
16:	RLBTY = 0.0	RELI	160
17:	DO 10 I=1,MPATH	RELI	170
18:	GPATH = PSIGN(I)	RELI	180
19:	DO 20 J=1,NMOD	RELI	190
20:	IF (PATH(I,J) .EQ. 0) GO TO 20	RELI	200
21:	GPATH = GPATH * G(J)	RELI	210
22:	20 CONTINUE	RELI	220
23:	RLBTY = RLBTY + GPATH	RELI	230
24:	10 CONTINUE	RELI	240
25:	RETURN	RELI	250
26:	END	RELI	260